

The structure of natural numbers

is helpful for proving
properties

$$\forall n[n \in \mathbb{N} : P(n)]$$

The structure of natural numbers

On natural numbers we can define a notion of a **successor**, a mapping

$$s: \mathbb{N} \rightarrow \mathbb{N}$$

by $s(n) = n+1$

The successor mapping imposes a structure on the set that enables us to **count**:

- 1) there is a **starting** natural number 0
- 2) for every natural number n , there is a **next** natural number $s(n) = n+1$.

(Some) Peano Axioms

Important properties

(I) Different natural numbers have different successors:

$$\forall n, m [n, m \in \mathbb{N} : s(m) = s(n) \Rightarrow m = n]$$

(Some) Peano Axioms

Important properties

(I) Different natural numbers have different successors:

$$\forall n, m [n, m \in \mathbb{N} : s(m) = s(n) \Rightarrow m = n]$$



stated positively

(Some) Peano Axioms

Important properties

(I) Different natural numbers have different successors:

$$\forall n, m [n, m \in \mathbb{N} : s(m) = s(n) \Rightarrow m = n]$$

stated positively

s is injective!

(Some) Peano Axioms

Important properties

(1) Different natural numbers have different successors:

$$\forall n, m [n, m \in \mathbb{N} : s(m) = s(n) \Rightarrow m = n]$$

stated positively

s is injective!

(2) 0 is not a successor: $\forall n [n \in \mathbb{N} : \neg (s(n) = 0)]$

(3) All natural numbers except 0 are successors:

$$\forall n [n \in \mathbb{N} \wedge \neg(n = 0) : \exists m [m \in \mathbb{N} : n = s(m)]]$$

There is more to it - induction

Imagine an infinite sequence of dominos



There is more to it - induction

Imagine an infinite sequence of dominos



If we know that

1. D_0 falls
2. The dominos are close enough together so that if D_i falls, then D_{i+1} falls (for all $i \in \mathbb{N}$)

Then we can conclude that every domino D_n ($n \in \mathbb{N}$) falls!

There is more to it - induction

Imagine an infinite sequence of dominos



If we know that

1. D_0 falls
2. The dominos are close enough together so that if D_i falls, then D_{i+1} falls (for all $i \in \mathbb{N}$)

Then we can conclude that every domino D_n ($n \in \mathbb{N}$) falls!



induction

Induction

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

Induction

P - unary predicate
over \mathbb{N}

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

Induction

P - unary predicate
over \mathbb{N}

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

$$P(0)$$
$$P(0) \Rightarrow P(1)$$

$$P(1)$$
$$P(1) \Rightarrow P(2)$$

$$P(2)$$
$$P(2) \Rightarrow P(3)$$

...

Induction

P - unary predicate
over \mathbb{N}

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim
with 0

$$P(0)$$
$$P(0) \Rightarrow P(1)$$

$$P(1)$$
$$P(1) \Rightarrow P(2)$$

$$P(2)$$
$$P(2) \Rightarrow P(3)$$

...

Induction

P - unary predicate
over \mathbb{N}

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim
with 0

\Rightarrow elim

$$P(0)$$
$$P(0) \Rightarrow P(1)$$

$$P(1)$$
$$P(1) \Rightarrow P(2)$$

$$P(2)$$
$$P(2) \Rightarrow P(3)$$

...

Induction

P - unary predicate
over \mathbb{N}

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim
with 0

\Rightarrow elim



$P(0)$
 $P(0) \Rightarrow P(1)$

$P(1)$
 $P(1) \Rightarrow P(2)$

$P(2)$
 $P(2) \Rightarrow P(3)$

...

Induction

P - unary predicate
over \mathbb{N}

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim
with 0

\Rightarrow elim



$$P(0) \\ P(0) \Rightarrow P(1)$$

$$P(1) \\ P(1) \Rightarrow P(2)$$

$$P(2) \\ P(2) \Rightarrow P(3)$$

...

Variant of the Peano Axiom:

Let $K \subseteq \mathbb{N}$ have the property that

(a) $0 \in K$ and

(b) for all $n \in \mathbb{N}$, $n \in K \Rightarrow (n+1) \in K$.

Then $K = \mathbb{N}$.

Induction

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

P - unary predicate
over \mathbb{N}

Induction

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

P - unary predicate
over \mathbb{N}

...

(m) P(0)
{Assume}

(k) **var** i; i \in \mathbb{N}

(k+1) P(i)
...

(l-1) P(i+1)
{ \Rightarrow -intro on (k+1) and (l-1)}

(l) P(i) \Rightarrow P(i+1)
{ \forall -intro on (k) and (l)}

(l+1) $\forall i[i \in \mathbb{N} : P(i) \Rightarrow P(i+1)]$
{induction on (m) and (l+1)}

(l+2) $\forall n[n \in \mathbb{N} : P(n)]$

Induction

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

P - unary predicate
over \mathbb{N}

...

(m) P(0)
{Assume}

(k) **var** i; i $\in \mathbb{N}$

(k+1) P(i)
...

(l-1) P(i+1)
{ \Rightarrow -intro on (k+1) and (l-1)}

(l) P(i) \Rightarrow P(i+1)
{ \forall -intro on (k) and (l)}

(l+1) $\forall i[i \in \mathbb{N} : P(i) \Rightarrow P(i+1)]$
{induction on (m) and (l+1)}

(l+2) $\forall n[n \in \mathbb{N} : P(n)]$

Basis

Induction

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

P - unary predicate
over \mathbb{N}

...

(m) $P(0)$
{Assume}

(k) **var** $i; i \in \mathbb{N}$

(k+1) $P(i)$
...

(l-1) $P(i+1)$
{ \Rightarrow -intro on (k+1) and (l-1)}

(l) $P(i) \Rightarrow P(i+1)$
{ \forall -intro on (k) and (l)}

(l+1) $\forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)]$
{induction on (m) and (l+1)}

(l+2) $\forall n [n \in \mathbb{N} : P(n)]$

Basis

Induction step

Induction

$$P(0) \wedge \forall i [i \in \mathbb{N} : P(i) \Rightarrow P(i+1)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

P - unary predicate
over \mathbb{N}

...

(m) P(0)
{Assume}

(k) **var** i; i \in \mathbb{N}

(k+1) P(i)
...

(l-1) P(i+1)
{ \Rightarrow -intro on (k+1) and (l-1)}

(l) P(i) \Rightarrow P(i+1)
{ \forall -intro on (k) and (l)}

(l+1) $\forall i[i \in \mathbb{N} : P(i) \Rightarrow P(i+1)]$
{induction on (m) and (l+1)}

(l+2) $\forall n[n \in \mathbb{N} : P(n)]$

Basis

induction
hypothesis

Induction step

Inductive definitions

Inductive proof: truth is passed on

Inductive definition: construction is passed on

Inductive definitions

Inductive proof: truth is passed on

Inductive definition: construction is passed on



well defined by induction

Inductive definitions

Inductive proof: truth is passed on

Inductive definition: construction is passed on

well defined by induction

Example

The sequence of real numbers $(a_i \mid i \in \mathbb{N})$ is defined inductively by

$$a_0 = 2$$

$$a_{i+1} = 2a_i - 1$$

Inductive definitions

Inductive proof: truth is passed on

Inductive definition: construction is passed on

well defined by induction

Example

The sequence of real numbers $(a_i \mid i \in \mathbb{N})$ is defined inductively by

$$a_0 = 2$$

$$a_{i+1} = 2a_i - 1$$

a_0	a_1	a_2	a_3	a_4	...
2	3	5	9	17	...

Inductive definitions

Inductive proof: truth is passed on

Inductive definition: construction is passed on

well defined by induction

Example

The sequence of real numbers $(a_i \mid i \in \mathbb{N})$ is defined inductively by

$$a_0 = 2$$

$$a_{i+1} = 2a_i - 1$$

a_0	a_1	a_2	a_3	a_4	...
2	3	5	9	17	...

Conjecture

For all $n \in \mathbb{N}$ it holds that

$$a_n = 2^{n+1}$$

Inductive definitions

Inductive proof: truth is passed on

Inductive definition: construction is passed on

well defined by induction

Example

The sequence of real numbers $(a_i \mid i \in \mathbb{N})$ is defined inductively by

$$a_0 = 2$$

$$a_{i+1} = 2a_i - 1$$

a_0	a_1	a_2	a_3	a_4	...
2	3	5	9	17	...

proof by induction

Conjecture

For all $n \in \mathbb{N}$ it holds that

$$a_n = 2^{n+1}$$

Strong induction

Strong induction

P - unary predicate
over \mathbb{N}

$$\forall k [k \in \mathbb{N} : \forall j [j \in \mathbb{N} \wedge j < k : P(j)] \Rightarrow P(k)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

Strong induction

P - unary predicate
over \mathbb{N}

$$\forall k [k \in \mathbb{N} : \forall j [j \in \mathbb{N} \wedge j < k : P(j)] \Rightarrow P(k)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

$$\begin{aligned} &P(0) \\ &P(0) \Rightarrow P(1) \\ &P(0) \wedge P(1) \\ &P(0) \wedge P(1) \Rightarrow P(2) \\ &P(0) \wedge P(1) \wedge P(2) \\ &P(0) \wedge P(1) \wedge P(2) \Rightarrow P(3) \\ &\dots \end{aligned}$$

Strong induction

P - unary predicate
over \mathbb{N}

$$\forall k [k \in \mathbb{N} : \forall j [j \in \mathbb{N} \wedge j < k : P(j)] \Rightarrow P(k)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim with $k=1$

$$\begin{aligned} & P(0) \\ & P(0) \Rightarrow P(1) \\ & P(0) \wedge P(1) \\ & P(0) \wedge P(1) \Rightarrow P(2) \\ & P(0) \wedge P(1) \wedge P(2) \\ & P(0) \wedge P(1) \wedge P(2) \Rightarrow P(3) \\ & \dots \end{aligned}$$

Strong induction

P - unary predicate
over \mathbb{N}

$$\forall k [k \in \mathbb{N} : \forall j [j \in \mathbb{N} \wedge j < k : P(j)] \Rightarrow P(k)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim with $k=1$

\Rightarrow elim,
 \wedge intro

$$\begin{aligned} &P(0) \\ &P(0) \Rightarrow P(1) \\ &P(0) \wedge P(1) \\ &P(0) \wedge P(1) \Rightarrow P(2) \\ &P(0) \wedge P(1) \wedge P(2) \\ &P(0) \wedge P(1) \wedge P(2) \Rightarrow P(3) \\ &\dots \end{aligned}$$

Strong induction

P - unary predicate
over \mathbb{N}

$$\forall k [k \in \mathbb{N} : \forall j [j \in \mathbb{N} \wedge j < k : P(j)] \Rightarrow P(k)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim with $k=1$

\Rightarrow elim,
 \wedge intro



$$\begin{aligned} &P(0) \\ &P(0) \Rightarrow P(1) \\ &P(0) \wedge P(1) \\ &P(0) \wedge P(1) \Rightarrow P(2) \\ &P(0) \wedge P(1) \wedge P(2) \\ &P(0) \wedge P(1) \wedge P(2) \Rightarrow P(3) \end{aligned}$$

...

Strong induction

P - unary predicate
over \mathbb{N}

$$\forall k [k \in \mathbb{N} : \forall j [j \in \mathbb{N} \wedge j < k : P(j)] \Rightarrow P(k)] \Rightarrow \forall n [n \in \mathbb{N} : P(n)]$$

\forall elim with $k=1$

\Rightarrow elim,
 \wedge intro



$$\begin{aligned} &P(0) \\ &P(0) \Rightarrow P(1) \\ &P(0) \wedge P(1) \\ &P(0) \wedge P(1) \Rightarrow P(2) \\ &P(0) \wedge P(1) \wedge P(2) \\ &P(0) \wedge P(1) \wedge P(2) \Rightarrow P(3) \\ &\dots \end{aligned}$$

Definition of
 $(a_i \mid i \in \mathbb{N})$
with strong
induction

a_n is defined via
 a_0, \dots, a_{n-1}

Cardinality

Cardinals

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A\rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A\rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A\rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

cardinal
numbers are
 \sim equivalence
classes

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A\rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A\rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

Def.

A set A has at most as large cardinality as a set B if there is an injection $f:A\rightarrow B$.
Notation $|A| \leq |B|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A\rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

Def.

A set A has at most as large cardinality as a set B if there is an injection $f:A\rightarrow B$.
Notation $|A| \leq |B|$.

Def.

A set A has at least as large cardinality as a set B if B is empty or there is a surjection $f:A\rightarrow B$.
Notation $|A| \geq |B|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A \rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

Def.

A set A has at most as large cardinality as a set B if there is an injection $f:A \rightarrow B$.
Notation $|A| \leq |B|$.

Def.

A set A has at least as large cardinality as a set B if B is empty or there is a surjection $f:A \rightarrow B$.
Notation $|A| \geq |B|$.

Def.

A set A has smaller cardinality than a set B if there is an injection $f:A \rightarrow B$ and there is no surjection $f:A \rightarrow B$. Notation $|A| < |B|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Cardinals

Def.

Two sets A and B have the same cardinality (are equinumerous) if there is a bijection $f:A \rightarrow B$.
Notation $A \sim B$, or $|A| = |B|$.

Prop.

The relation \sim is an equivalence relation on sets.

Def.

A set A has at most as large cardinality as a set B if there is an injection $f:A \rightarrow B$.
Notation $|A| \leq |B|$.

Def.

A set A has at least as large cardinality as a set B if B is empty or there is a surjection $f:A \rightarrow B$.
Notation $|A| \geq |B|$.

Def.

A set A has smaller cardinality than a set B if there is an injection $f:A \rightarrow B$ and there is no surjection $f:A \rightarrow B$. Notation $|A| < |B|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Theorem (Cantor)

If $|A| \leq |B|$
and
 $|B| \leq |A|$,
then
 $|A| = |B|$.

Operations on cardinals

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Operations on cardinals

Def.

Let A and B be two disjoint sets. Then
 $|A| + |B| = |A \cup B|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Operations on cardinals

Def.

Let A and B be two disjoint sets. Then
 $|A| + |B| = |A \cup B|$.

Def.

Let A and B be two sets. Then
 $|A| \cdot |B| = |A \times B|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Operations on cardinals

Def.

Let A and B be two disjoint sets. Then $|A| + |B| = |A \cup B|$.

Def.

Let A and B be two sets. Then $|A| \cdot |B| = |A \times B|$.

Def.

Let A and B be two sets. Then $|A|^{|B|} = |A^B|$ where A^B is the set of all functions from B to A , i.e. $A^B = \{f \mid f: B \rightarrow A\}$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Operations on cardinals

Def.

Let A and B be two disjoint sets. Then $|A| + |B| = |A \cup B|$.

Def.

Let A and B be two sets. Then $|A| \cdot |B| = |A \times B|$.

Def.

Let A and B be two sets. Then $|A|^{|B|} = |A^B|$ where A^B is the set of all functions from B to A , i.e. $A^B = \{f \mid f: B \rightarrow A\}$.

Prop.

Let A be a set. Then $|\mathcal{P}(A)| = 2^{|A|}$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Operations on cardinals

Def.

Let A and B be two disjoint sets. Then $|A| + |B| = |A \cup B|$.

Def.

Let A and B be two sets. Then $|A| \cdot |B| = |A \times B|$.

Def.

Let A and B be two sets. Then $|A|^{|B|} = |A^B|$ where A^B is the set of all functions from B to A , i.e. $A^B = \{f \mid f: B \rightarrow A\}$.

Prop.

Let A be a set. Then $|\mathcal{P}(A)| = 2^{|A|}$.

$$|A| = [A]_{\sim}$$

cardinal numbers are \sim equivalence classes

Note: $2 = |\{0, 1\}|$

Finite sets, finite cardinals

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Finite sets, finite cardinals

We write \mathbb{N}_k for the set $\{0, 1, \dots, k-1\}$. Then $\mathbb{N}_0 = \emptyset$.

We will also write k for $|\mathbb{N}_k|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Finite sets, finite cardinals

We write \mathbb{N}_k for the set $\{0, 1, \dots, k-1\}$. Then $\mathbb{N}_0 = \emptyset$.

We will also write k for $|\mathbb{N}_k|$.

Def.

A set A is finite if and only if $|A| = |\mathbb{N}_k|$, for some $k \in \mathbb{N}$. We write then $|A| = k$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Finite sets, finite cardinals

We write \mathbb{N}_k for the set $\{0, 1, \dots, k-1\}$. Then $\mathbb{N}_0 = \emptyset$.

We will also write k for $|\mathbb{N}_k|$.

Def.

A set A is finite if and only if $|A| = |\mathbb{N}_k|$, for some $k \in \mathbb{N}$. We write then $|A| = k$.

Hence

A set A is finite if and only if there is a natural number $k \in \mathbb{N}$ and a bijection $f: A \rightarrow \mathbb{N}_k$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Finite sets, finite cardinals

We write \mathbb{N}_k for the set $\{0, 1, \dots, k-1\}$. Then $\mathbb{N}_0 = \emptyset$.

We will also write k for $|\mathbb{N}_k|$.

Def.

A set A is finite if and only if $|A| = |\mathbb{N}_k|$, for some $k \in \mathbb{N}$. We write then $|A| = k$.

Hence

A set A is finite if and only if there is a natural number $k \in \mathbb{N}$ and a bijection $f: A \rightarrow \mathbb{N}_k$.

if and only if A has k elements, for some $k \in \mathbb{N}$

$$|A| = [A]_{\sim}$$

cardinal numbers are \sim equivalence classes

Finite sets, finite cardinals

We write \mathbb{N}_k for the set $\{0, 1, \dots, k-1\}$. Then $\mathbb{N}_0 = \emptyset$.

We will also write k for $|\mathbb{N}_k|$.

Def.

A set A is finite if and only if $|A| = |\mathbb{N}_k|$, for some $k \in \mathbb{N}$. We write then $|A| = k$.

Hence

A set A is finite if and only if there is a natural number $k \in \mathbb{N}$ and a bijection $f: A \rightarrow \mathbb{N}_k$.

if and only if A has k elements, for some $k \in \mathbb{N}$

$$|A| = [A]_{\sim}$$

cardinal numbers are \sim equivalence classes

The operations on cardinals when restricted to finite cardinals coincide with the operations on natural numbers!

This justifies the notation.

Finite sets, finite cardinals

We write \mathbb{N}_k for the set $\{0, 1, \dots, k-1\}$. Then $\mathbb{N}_0 = \emptyset$.

We will also write k for $|\mathbb{N}_k|$.

Def.

A set A is finite if and only if $|A| = |\mathbb{N}_k|$, for some $k \in \mathbb{N}$. We write then $|A| = k$.

Hence

A set A is finite if and only if there is a natural number $k \in \mathbb{N}$ and a bijection $f: A \rightarrow \mathbb{N}_k$.

$$|A| = [A]_{\sim}$$

cardinal numbers are \sim equivalence classes

if and only if A has k elements, for some $k \in \mathbb{N}$

E.g. If $|A| = k$ and $|B| = m$ for some $k, m \in \mathbb{N}$ then $|A \times B| = k \cdot m$

The operations on cardinals when restricted to finite cardinals coincide with the operations on natural numbers!
This justifies the notation.

Infinite, countable and uncountable sets

Time for a video!

Infinite, countable and uncountable sets

Time for a video!

Hilbert's
infinite hotel :-)

Infinite, countable and uncountable sets

Time for a video!

Hilbert's
infinite hotel :-)

Infinite, countable and uncountable sets

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

Prop.

\mathbb{N} is countable.
 \mathbb{Z} is countable.
 \mathbb{Q} is countable.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

Prop.

\mathbb{N} is countable.

\mathbb{Z} is countable.

\mathbb{Q} is countable.

Def.

A set is infinite iff $|A| \geq \aleph_0$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

Prop.

\mathbb{N} is countable.

\mathbb{Z} is countable.

\mathbb{Q} is countable.

Def.

A set is infinite iff $|A| \geq \aleph_0$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Hence, every countable set
is infinite

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

Prop.

\mathbb{N} is countable.
 \mathbb{Z} is countable.
 \mathbb{Q} is countable.

Def.

A set is infinite iff $|A| \geq \aleph_0$.

Def.

A set is uncountable iff $|A| > \aleph_0$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Hence, every countable set
is infinite

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

Prop.

\mathbb{N} is countable.
 \mathbb{Z} is countable.
 \mathbb{Q} is countable.

Def.

A set is infinite iff $|A| \geq \aleph_0$.

Def.

A set is uncountable iff $|A| > \aleph_0$.

Prop.

\mathbb{R} is uncountable.

$$|A| = [A]_{\sim}$$

cardinal numbers are \sim equivalence classes

Hence, every countable set is infinite

Infinite, countable and uncountable sets

We write \aleph_0 for the cardinality of natural numbers.
Hence $\aleph_0 = |\mathbb{N}|$.

Def.

A set A is countable iff $|A| = \aleph_0$.

Prop.

\mathbb{N} is countable.
 \mathbb{Z} is countable.
 \mathbb{Q} is countable.

Def.

A set is infinite iff $|A| \geq \aleph_0$.

Def.

A set is uncountable iff $|A| > \aleph_0$.

Prop.

\mathbb{R} is uncountable.

$$|A| = [A]_{\sim}$$

cardinal numbers are \sim equivalence classes

Hence, every countable set is infinite

We write c for $|\mathbb{R}|$

Cardinals are unbounded

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Cardinals are unbounded

Theorem (Cantor)

For every set A we have $|A| < |\mathcal{P}(A)|$.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes

Cardinals are unbounded

Theorem (Cantor)

For every set A we have $|A| < |\mathcal{P}(A)|$.

Hence, for every cardinal there is a larger one.

$$|A| = [A]_{\sim}$$

cardinal
numbers are
 \sim equivalence
classes