

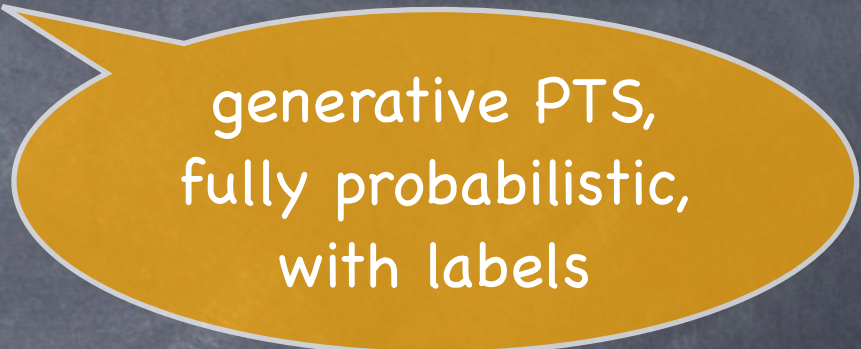
Sound and Complete Axiomatization of Trace Semantics for Probabilistic Transition Systems

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CWI and University of Salzburg

MFPS 2011, Pittsburgh, 27.5.2011

We will discuss

- history
- probabilistic transition systems
- (finite) trace semantics
- the sound and complete axiomatization
- in a coalgebraic setting



generative PTS,
fully probabilistic,
with labels

History

- For LTS Milner '84, JCSS
- expressions for LTS
- Kleene style theorem
- axiomatization
- sound and complete for bisimilarity

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$$P + Q \equiv Q + P, P + 0 \equiv P, \mu x.P \equiv P[\mu x.P/x], \dots$$

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$$P \equiv Q \iff P \sim Q$$

Some years later

- Milner's result was extended by Rabinovich '93 MFPS for trace semantics
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Milner's axioms plus $a.(P + Q) \equiv a.P + a.Q$

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Milner's axioms plus $a.(P + Q) \equiv a.P + a.Q$

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$$P \equiv Q \iff \text{tr}(P) = \text{tr}(Q)$$

Now we do it for PTS

- Expressions/axioms for PTS come in many flavors (mainly for bisimilarity)
we build on Silva, Bonchi, Bonsangue, Rutten '09/'10
- Trace semantics for PTS also exists in variants
we build on Hasuo, Jacobs, S. '06/'07

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expressions, Kleene style theorem, sound and complete
axiomatization for bisimilarity
(generic approach applicable to a large class of systems)

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same expressions, one more axiom,
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Probabilistic transition systems

PTS here are generative, labelled, with explicit termination

$$X \rightarrow \mathcal{D}_\omega(1 + A \times X)$$

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$$\mathcal{D}_\omega X = \{\mu: X \rightarrow [0, 1] \mid \sum_{x \in X} \mu(x) \leq 1, \text{supp}(X) \text{ is finite}\}$$

Probabilistic transition systems

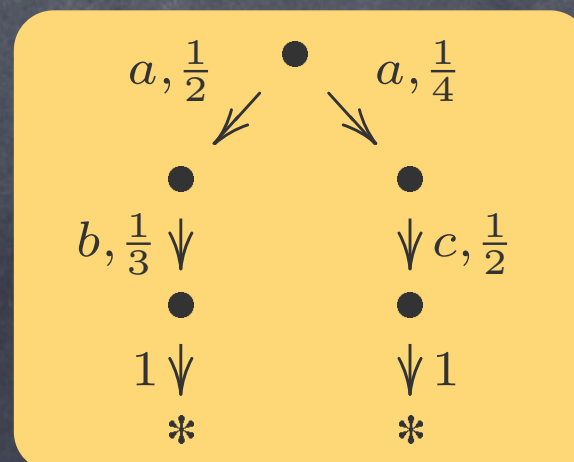
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Example:



Coalgebra basics

Category $\mathbf{C}$, functor F , category of coalgebras:

Coalg _{F}

Objects:

$$X \xrightarrow{c} FX$$

Arrows:

$$\begin{array}{ccc} X & \xrightarrow{h} & Y \\ c \downarrow & & \downarrow d \\ FX & \xrightarrow{Fh} & FY \end{array}$$

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Final coalgebra semantics:

$$\begin{array}{ccc} X & \xrightarrow{\exists! \text{beh}} & \Omega_F \\ c \downarrow & & \downarrow \cong \\ FX & \xrightarrow{F \text{beh}} & F\Omega_F \end{array}$$

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bisimilarity in Sets

(for wpp functors)

trace semantics in $\mathcal{Kl}(T)$

(for TF -coalgebras)

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Coalgebraic trace semantics [HJS'06/'07] applies to PTS

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$\text{tr} : X \rightarrow \mathcal{D}(A^*)$ in **Sets**

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It instantiates to finite trace distribution:

$$\begin{aligned} \text{tr}(x)(\varepsilon) &= c(x)(*) \\ \text{tr}(x)(aw) &= \sum_{x' \in X} c(x)(a, x') \cdot \text{tr}(x')(w) \end{aligned}$$

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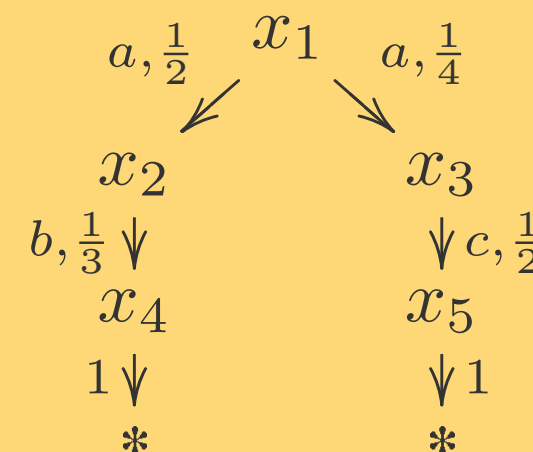
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$$\text{tr}(x_1)(ab) = \frac{1}{6}$$

$$\text{tr}(x_1)(ac) = \frac{1}{8}$$

Expressions for PTS

Are the same ones as for bisimilarity

Silva, Bonchi, Bonsanque, Rutten '09/'10

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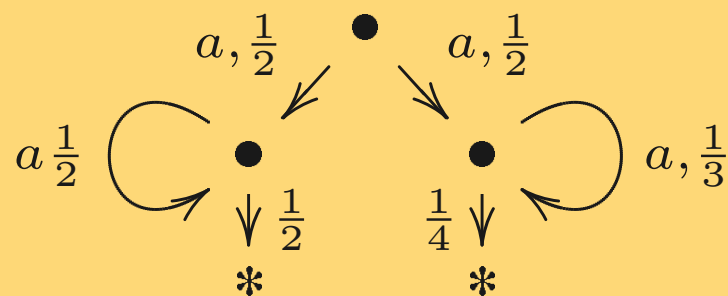
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$$\begin{aligned}
 &\frac{1}{2} \cdot a \cdot \mu x. \left(\frac{1}{2} \cdot a \cdot x \oplus \frac{1}{2} \cdot * \right) \\
 &\oplus \frac{1}{4} \cdot a \cdot \mu y. \left(\frac{1}{3} \cdot a \cdot y \oplus \frac{1}{2} \cdot * \right)
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Expressions for PTS

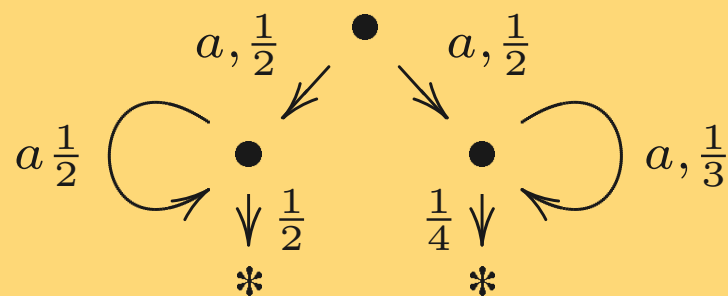
expressions behave! (Kleene-style theorem)

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Axioms

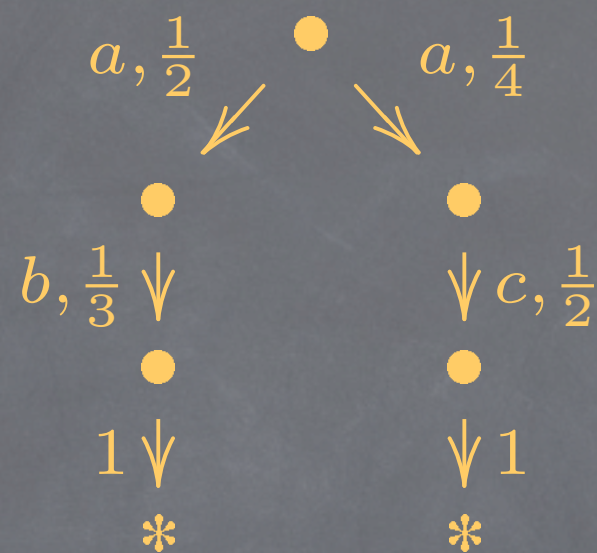
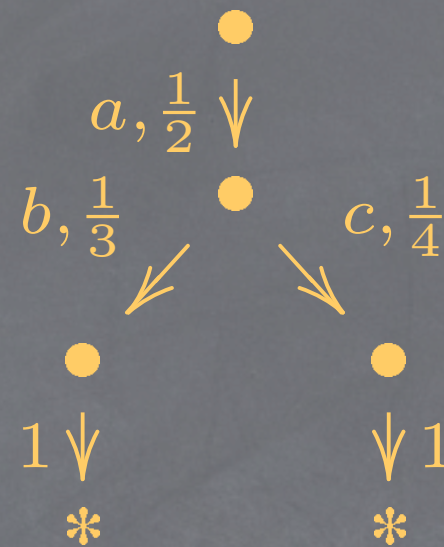
bisimilarity

$E_1 \oplus (E_2 \oplus E_3)$	$\equiv$	$(E_1 \oplus E_2) \oplus E_3$	(A)
$E_1 \oplus E_2$	$\equiv$	$E_2 \oplus E_1$	(C)
$E \oplus \underline{\emptyset}$	$\equiv$	E	(E)
$\mu x.E$	$\equiv$	$E[\mu x.E/x]$	(FP)
$\gamma[E/x] \equiv E$	$\Rightarrow$	$\mu x.\gamma \equiv E$	(UFP)
$\mu x.E$	$\equiv$	$\mu y.E[y/x]$ if y is not free in E	$(\alpha - equiv)$
$E_1 \equiv E_2$	$\Rightarrow$	$E[E_1/x] \equiv E[E_2/x]$	$(Cong)$
$0 \cdot E$	$\equiv$	$\underline{\emptyset}$	(Z)
$p \cdot E \oplus p' \cdot E$	$\equiv$	$\underline{(p + p') \cdot E}$	(S)

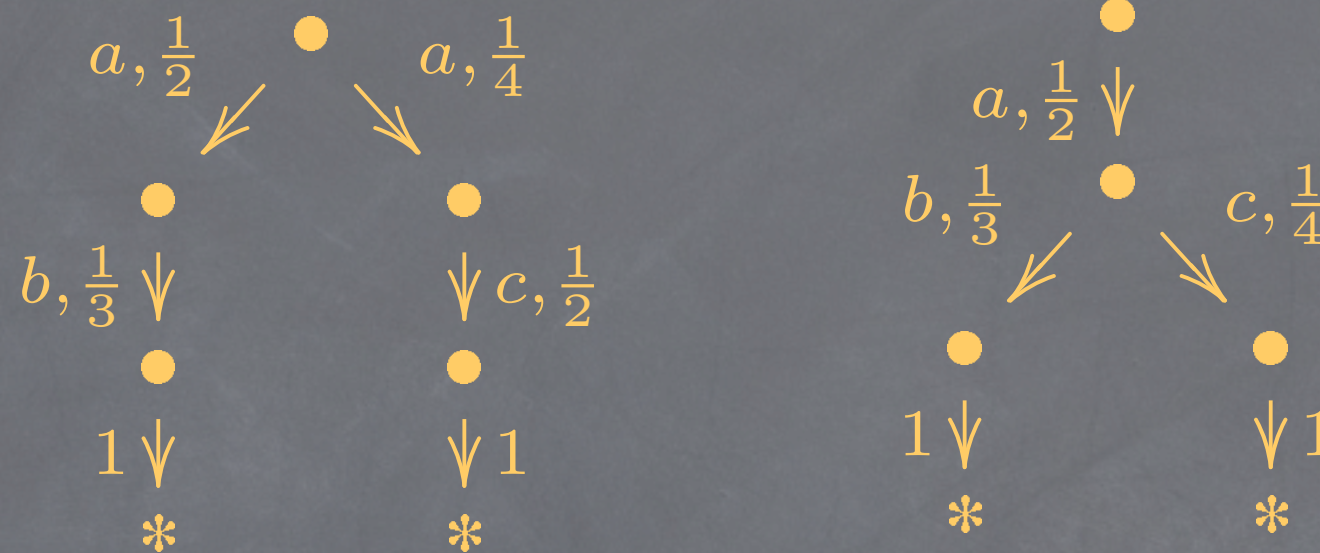
$$p \cdot a \cdot (p_1 E_1 \oplus p_2 E_2) \equiv p_1 \cdot a \cdot p E_1 \oplus p_2 \cdot a \cdot p E_2 \quad (D)$$

trace

Example

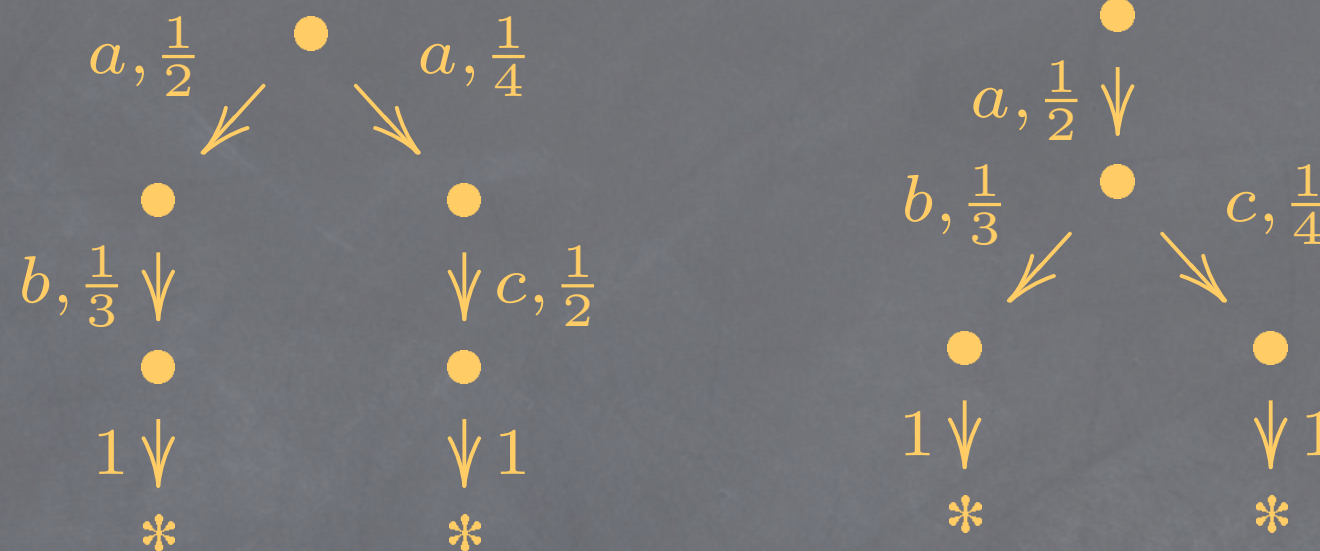


Example



$$\begin{aligned} \left(\frac{1}{2} \cdot a \cdot \frac{1}{3} \cdot b \cdot 1 \cdot * \right) \oplus \left(\frac{1}{4} \cdot a \cdot \frac{1}{2} \cdot c \cdot 1 \cdot * \right) &\stackrel{(D)}{=} \frac{1}{2} \cdot a \cdot \left(\frac{1}{2} \left(\frac{2}{3} \cdot b \cdot 1 \cdot * \right) \oplus \frac{1}{4} (1 \cdot c \cdot 1 \cdot *) \right) \\ &= \frac{1}{2} \cdot a \cdot \left(\frac{1}{3} \cdot b \cdot 1 \cdot * \oplus \frac{1}{4} \cdot c \cdot 1 \cdot * \right) \end{aligned}$$

Example



$$\frac{1}{3} \cdot b \cdot 1 \cdot * = \frac{1}{2} \left(\frac{2}{3} \cdot b \cdot 1 \cdot * \right), \quad \frac{1}{2} \cdot c \cdot 1 \cdot * = \frac{1}{2} (1 \cdot c \cdot 1 \cdot *)$$

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Soundness and Completeness

Find an **injective** map $out_{\equiv}$ with $tr = out_{\equiv} \circ [-]$

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canonical map to
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Soundness

$$\begin{aligned} & E_1 \equiv E_2 \\ \Leftrightarrow & [E_1] = [E_2] \\ \stackrel{(*)}{\Rightarrow} & out_{\equiv}([E_1]) = out_{\equiv}([E_2]) \\ \stackrel{(\Delta)}{\Leftrightarrow} & tr(E_1) = tr(E_2) \end{aligned}$$

Completeness

$$\begin{aligned} & tr(E_1) = tr(E_2) \\ \stackrel{(\Delta)}{\Leftrightarrow} & out_{\equiv}([E_1]) = out_{\equiv}([E_2]) \\ \stackrel{(\heartsuit)}{\Rightarrow} & [E_1] = [E_2] \\ \Leftrightarrow & E_1 \equiv E_2 \end{aligned}$$

Soundness and Completeness

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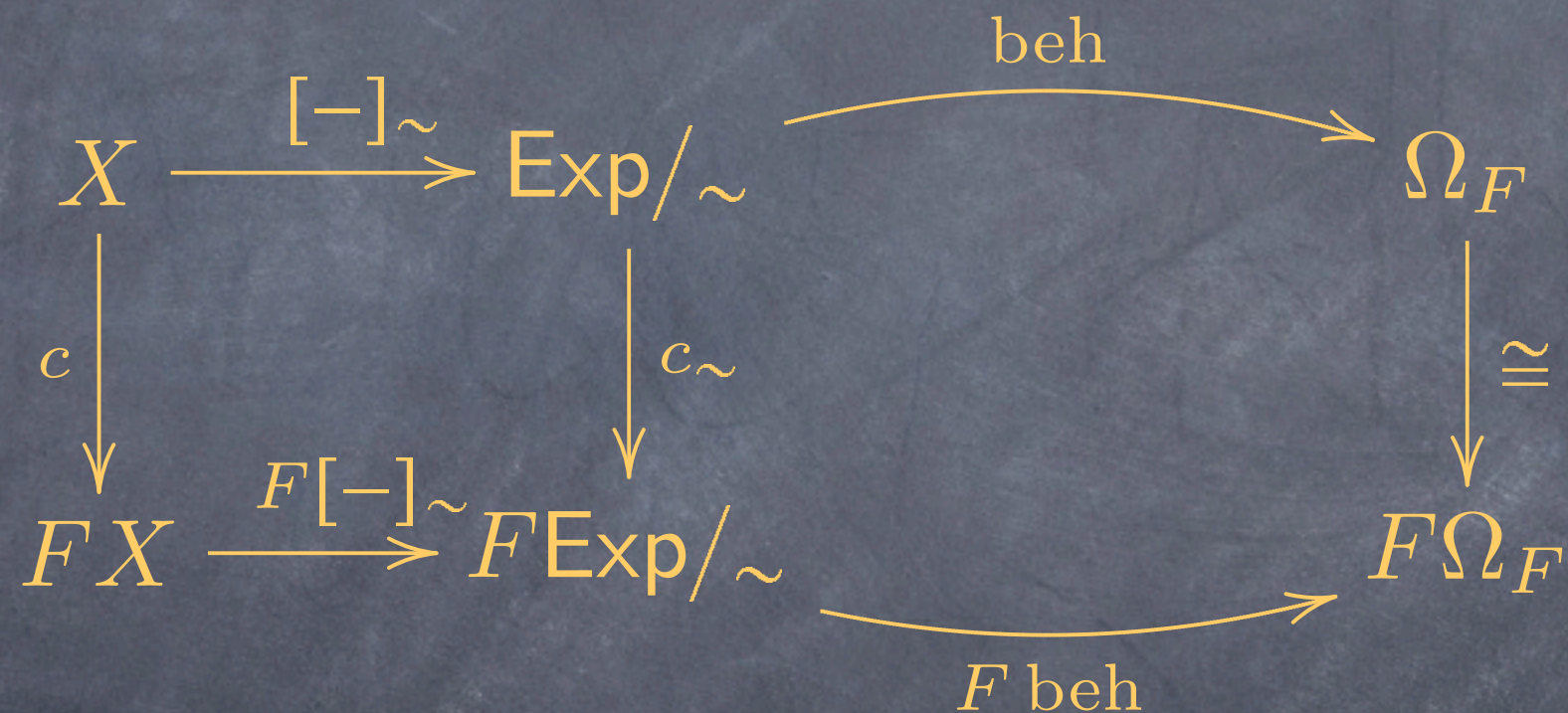
$(*)$ - existence of $out_{\equiv}$ $(\triangle)$ - $tr = out_{\equiv} \circ [-]$ $(\heartsuit)$ - injectivity

How to get out?

Bisimilarity case, F-coalgebras

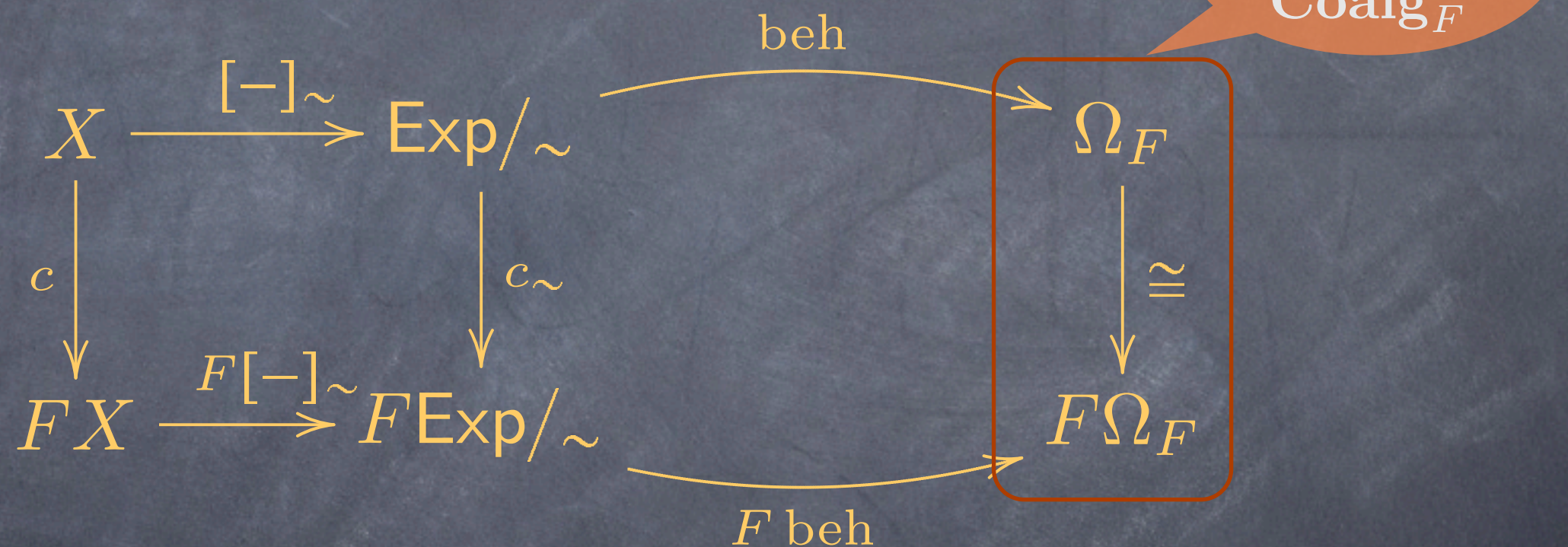
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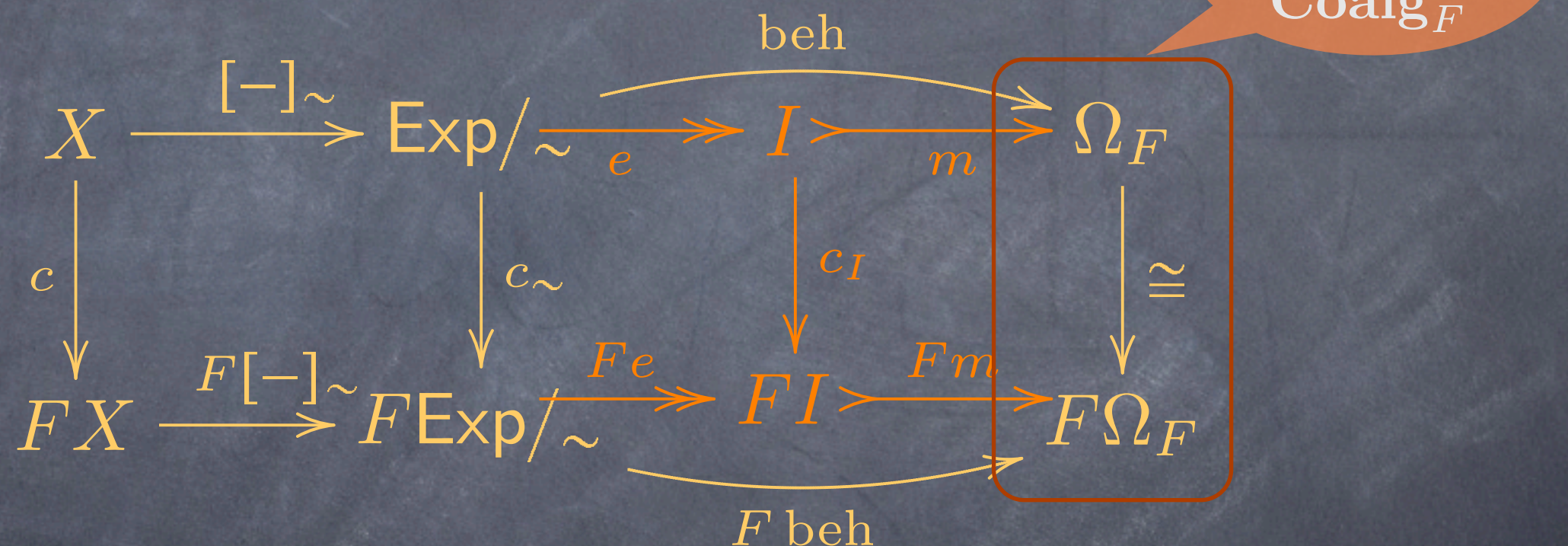
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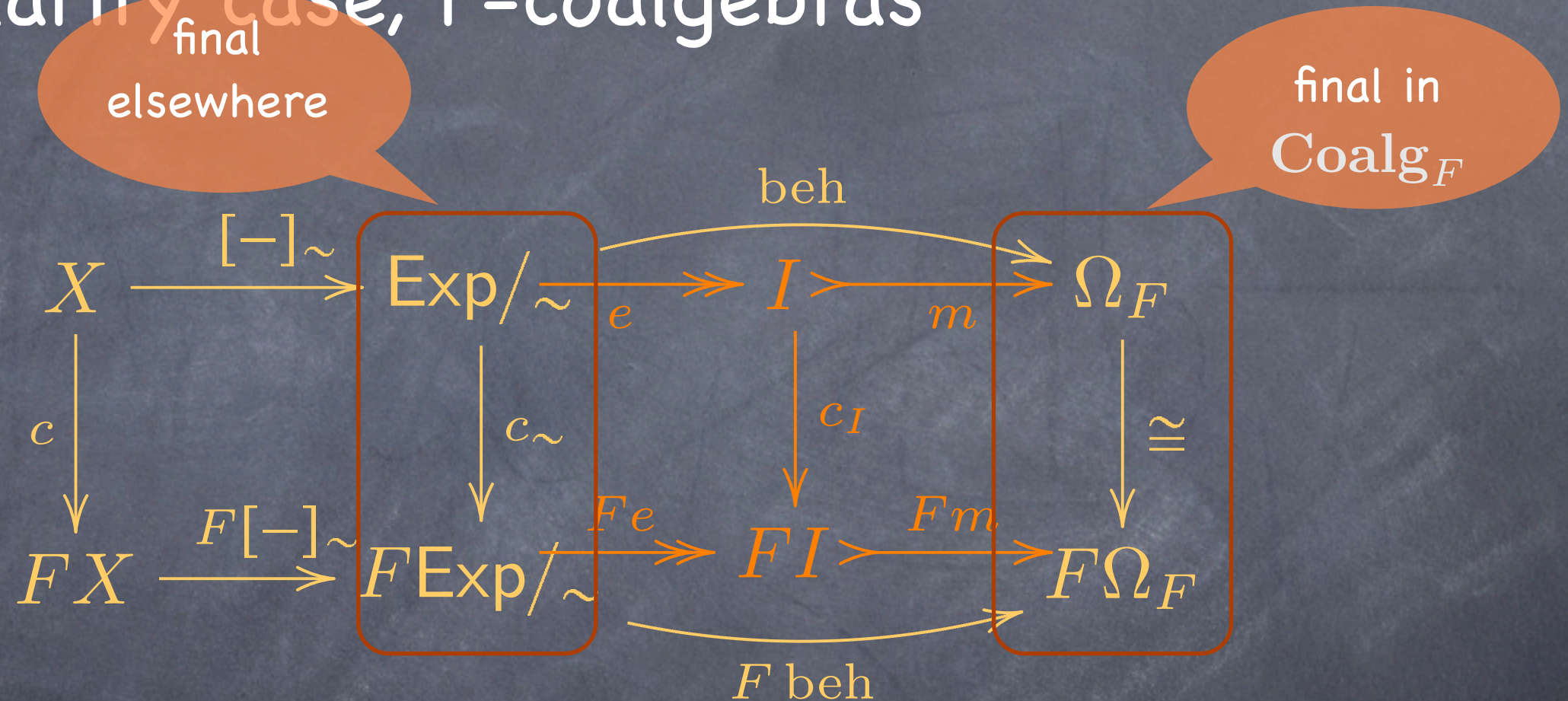
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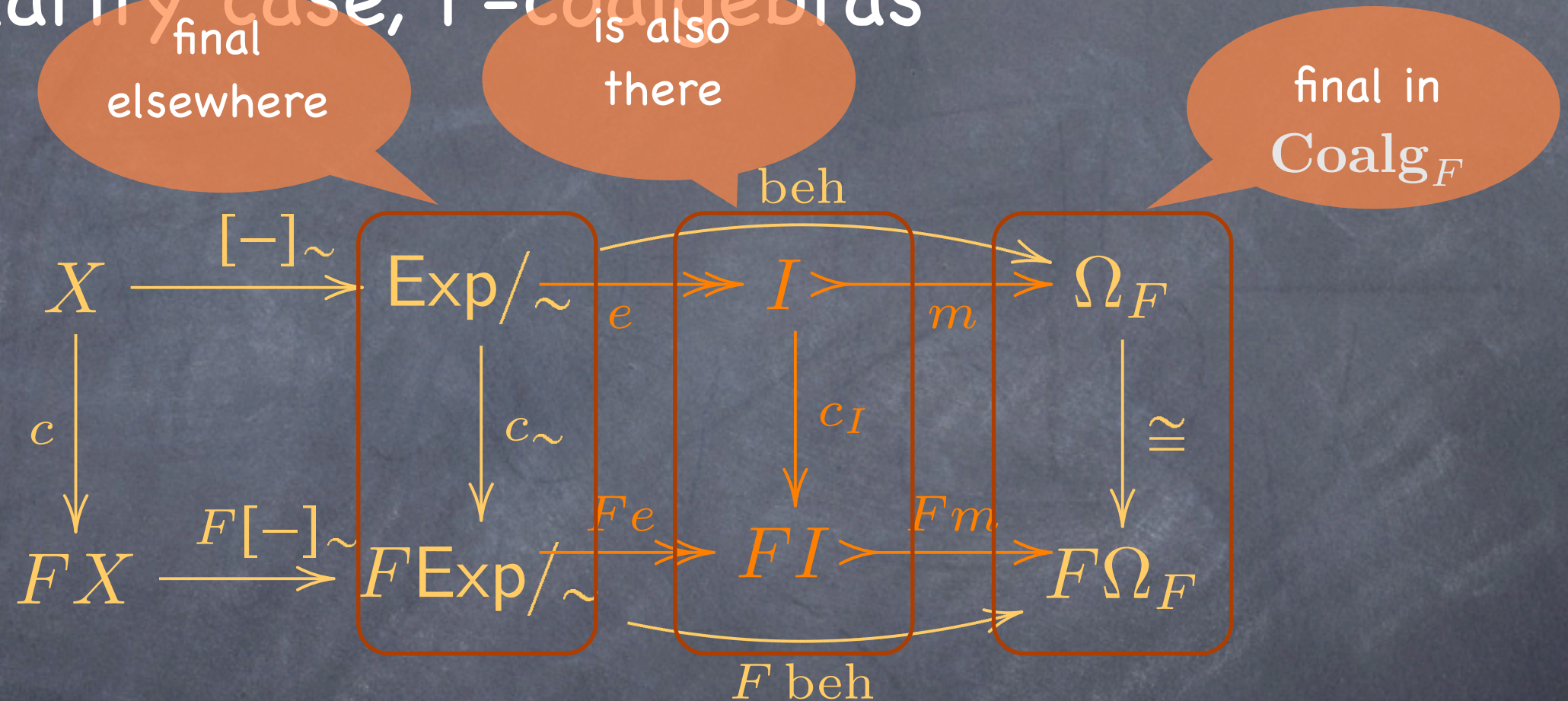
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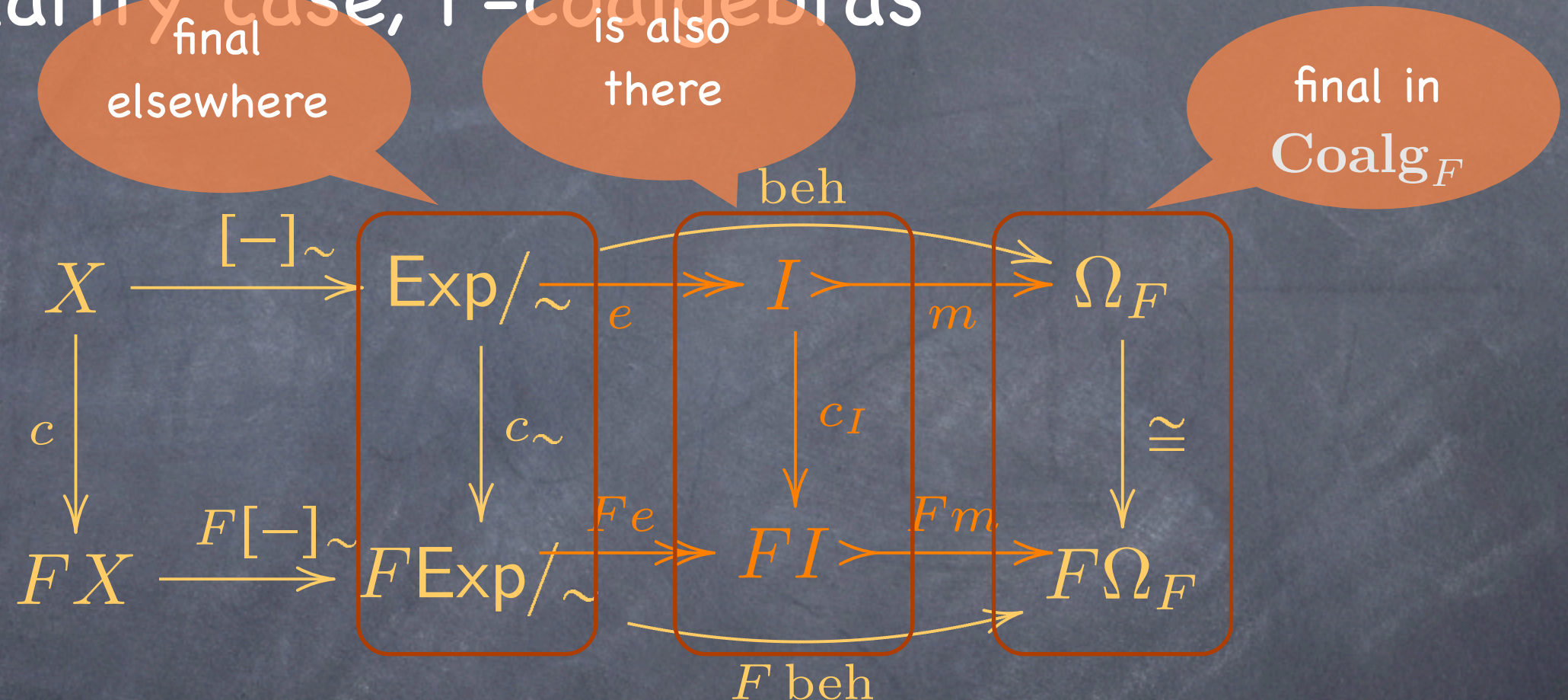
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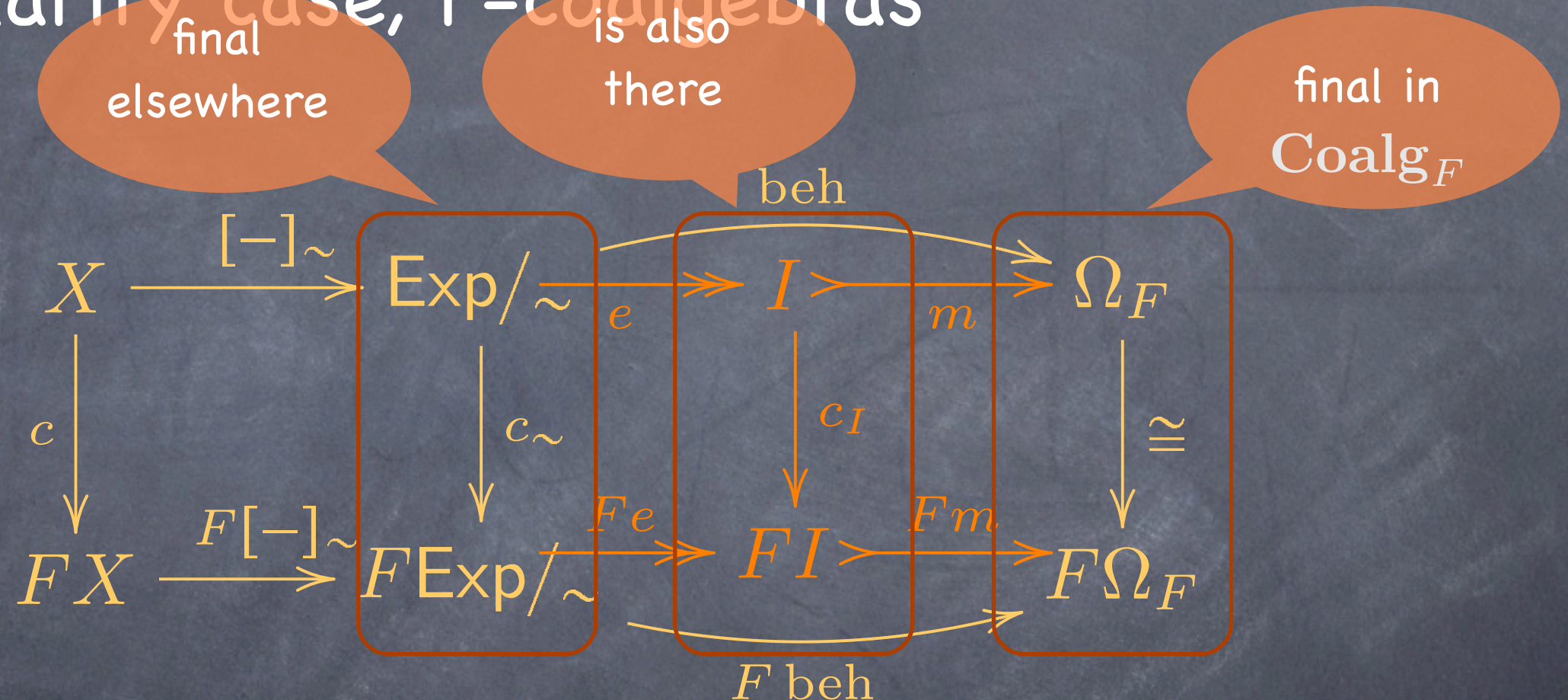
Bisimilarity case, F -coalgebras



Hence, e is iso, and $\text{out} = \text{beh}$ is injective

How to get out?

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Hence, e is iso, and $\text{out} = \text{beh}$ is injective

Silva et al. '08/'09/'10, Jacobs'06

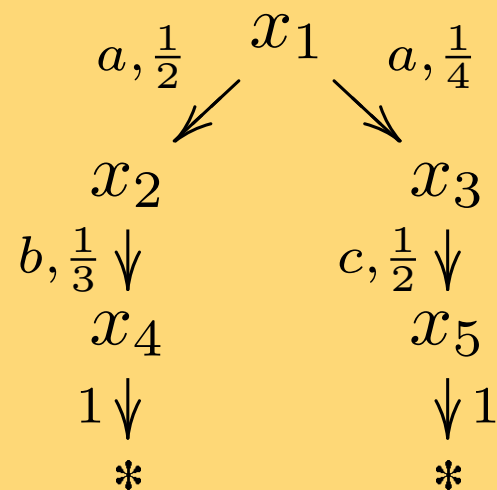
there are also algebras around

A way out for traces?

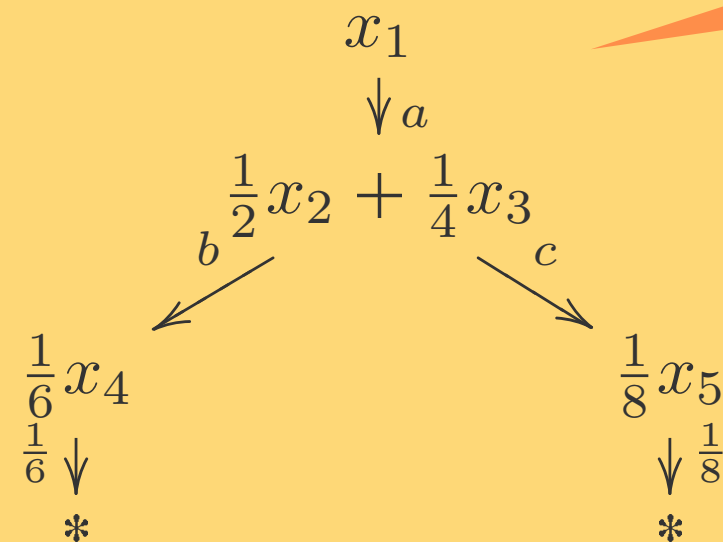
- It is tough to work in Kleisli categories
- Factorization ?
- So we find a way to stay in **Sets**
or rather in **Sets** ^{$\mathcal{D}_\omega$}
- A way out – determinization

Determinization of PTS

PTS example



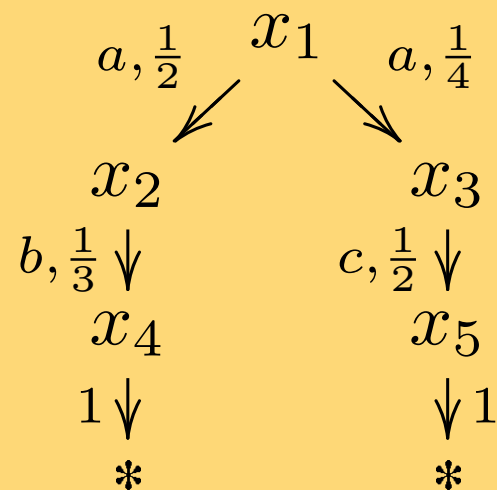
Its determinization



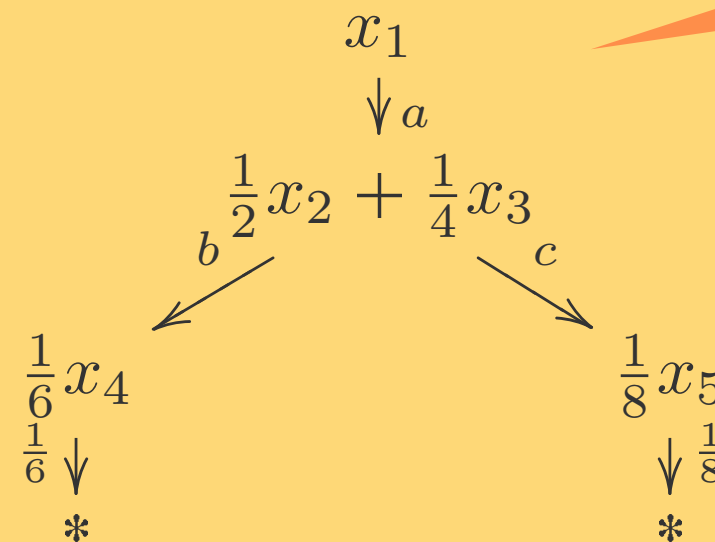
a G-coalgebra

Determinization of PTS

PTS example



Its determinization



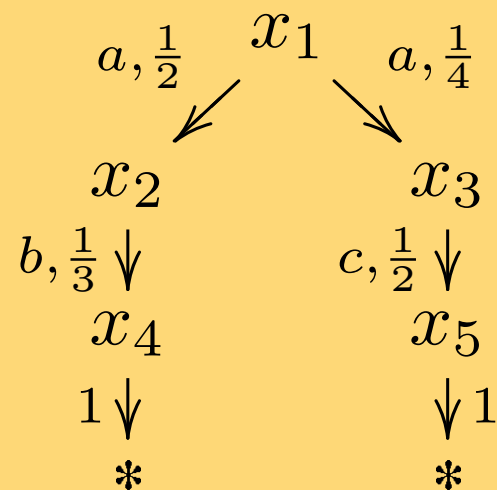
a G-coalgebra

$$GX = [0, 1] \times X^A$$

$$\begin{array}{ccc}
 X & \xrightarrow{\text{out}} & [0, 1]^{A^*} \\
 \downarrow c & & \downarrow \cong \\
 GX & \xrightarrow{G\text{out}} & G([0, 1]^{A^*})
 \end{array}$$

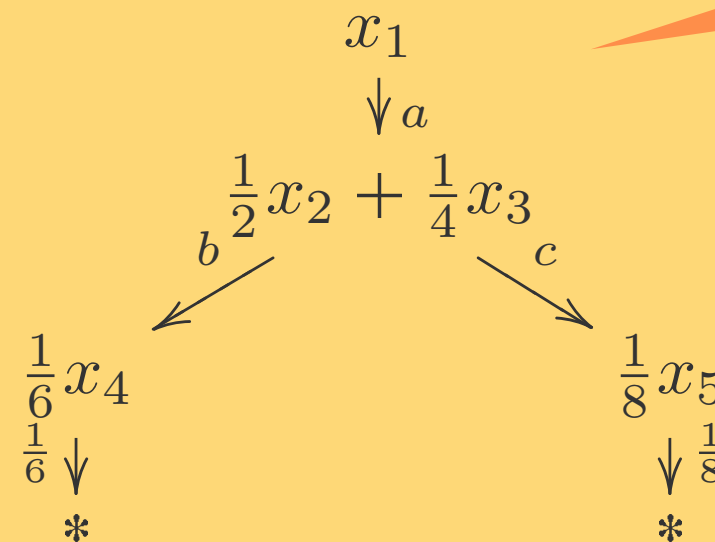
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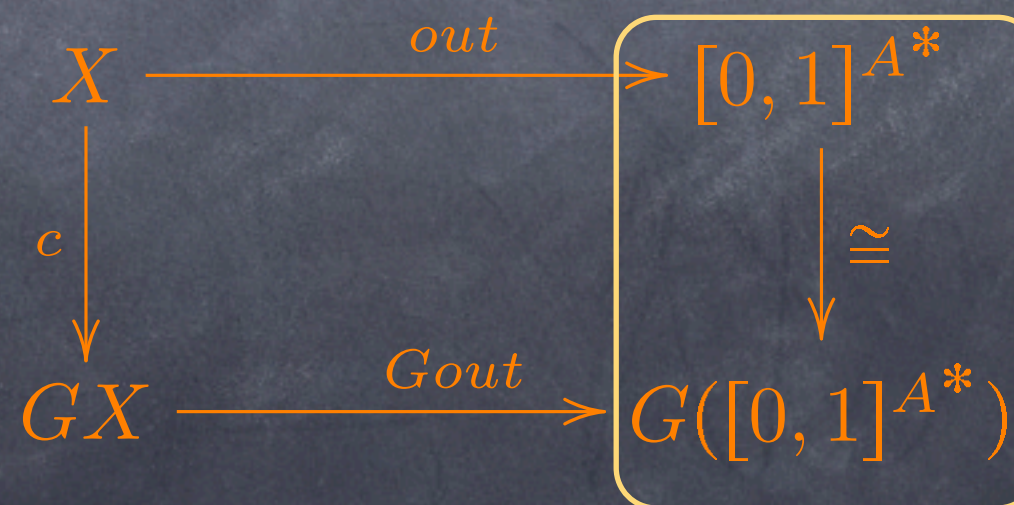
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Its determinization



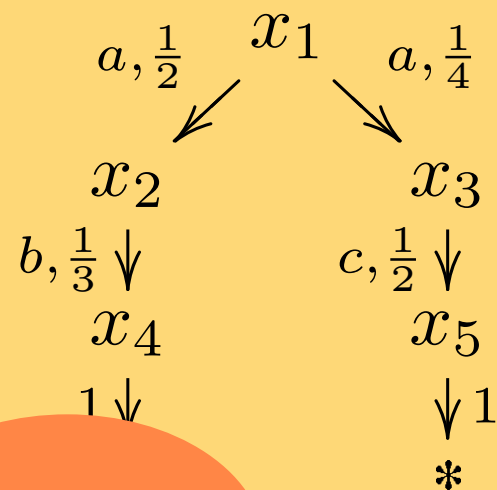
a G-coalgebra

final
G-coalgebra



Determinization of PTS

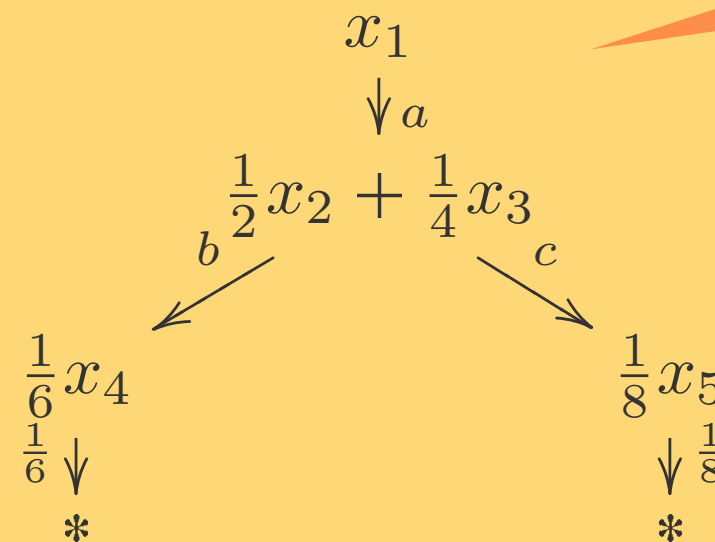
PTS example



Sets $\mathcal{D}_\omega$

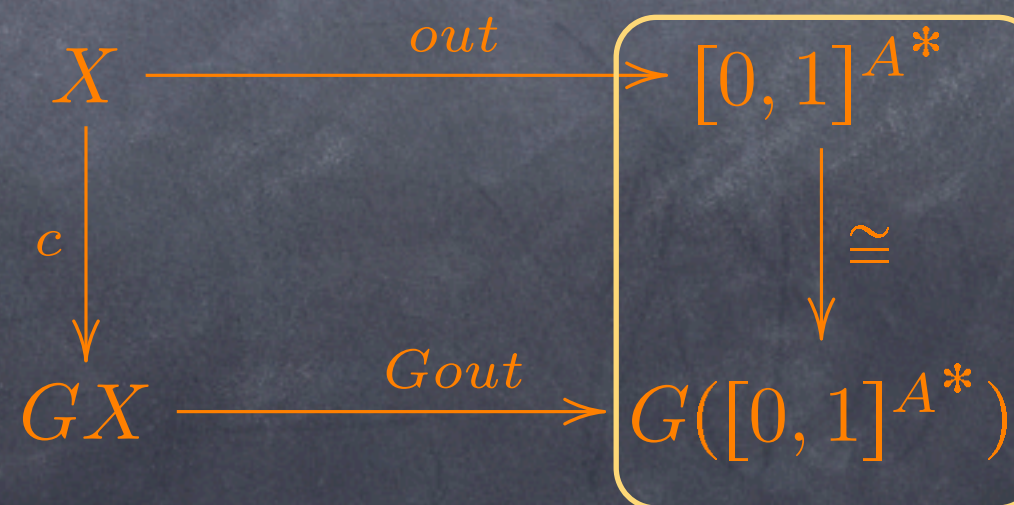
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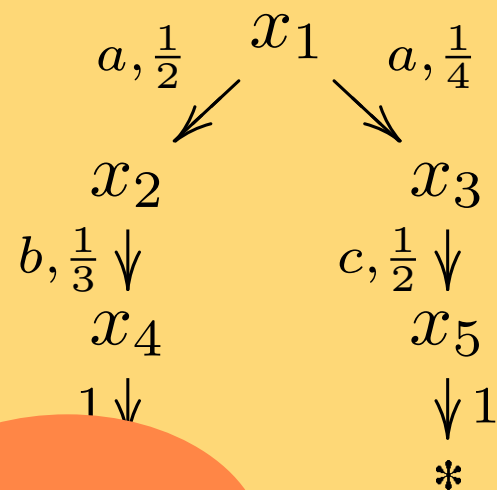
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Determinization of PTS

PTS example



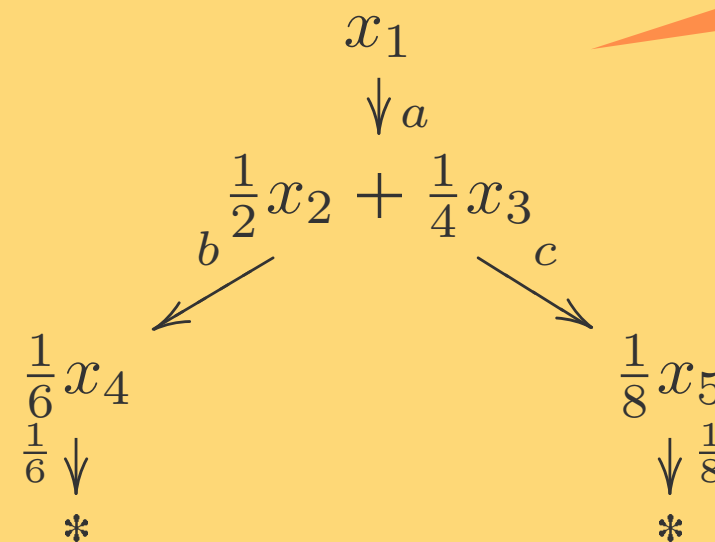
Sets $\mathcal{D}_\omega$

$$GX = [0, 1] \times X^A$$

$$\begin{array}{ccc} X & \xrightarrow{\eta} & \mathcal{D}_\omega(X) \\ \downarrow c & & \downarrow (\delta \circ c)^\# \\ \mathcal{D}_\omega(1 + A \times X) & \xrightarrow{\delta} & G\mathcal{D}_\omega(X) \end{array}$$

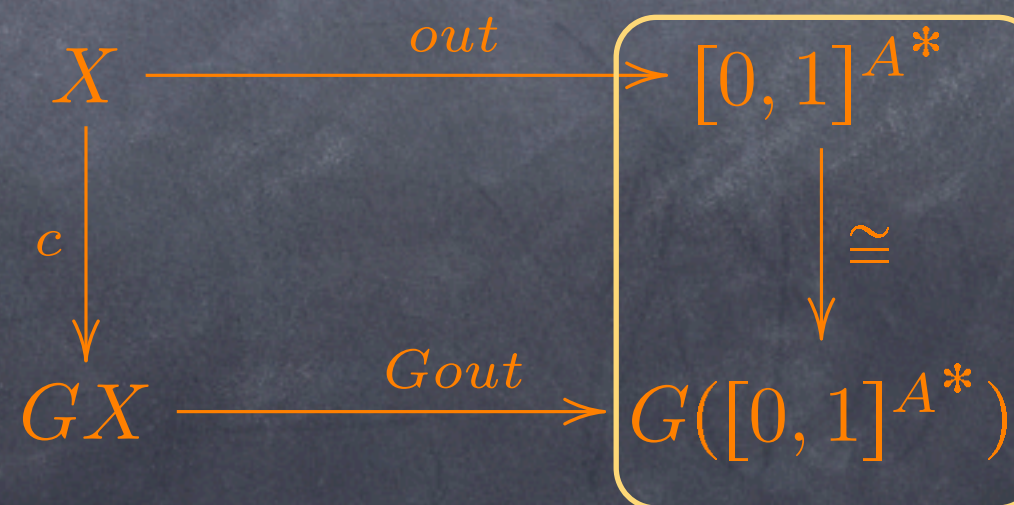
$$\mathcal{D}_\omega(1 + A \times X) \xrightarrow{\delta} G\mathcal{D}_\omega(X)$$

Its determinization



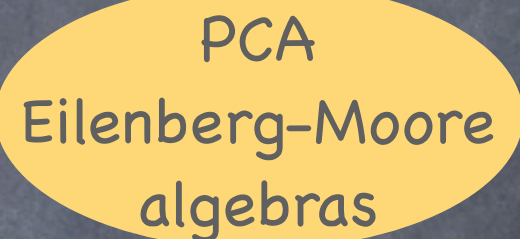
a G-coalgebra

final
G-coalgebra



A way out for traces

Trace case, almost G-coalgebras on $\mathbf{Sets}^{\mathcal{D}_\omega}$



PCA
Eilenberg-Moore
algebras

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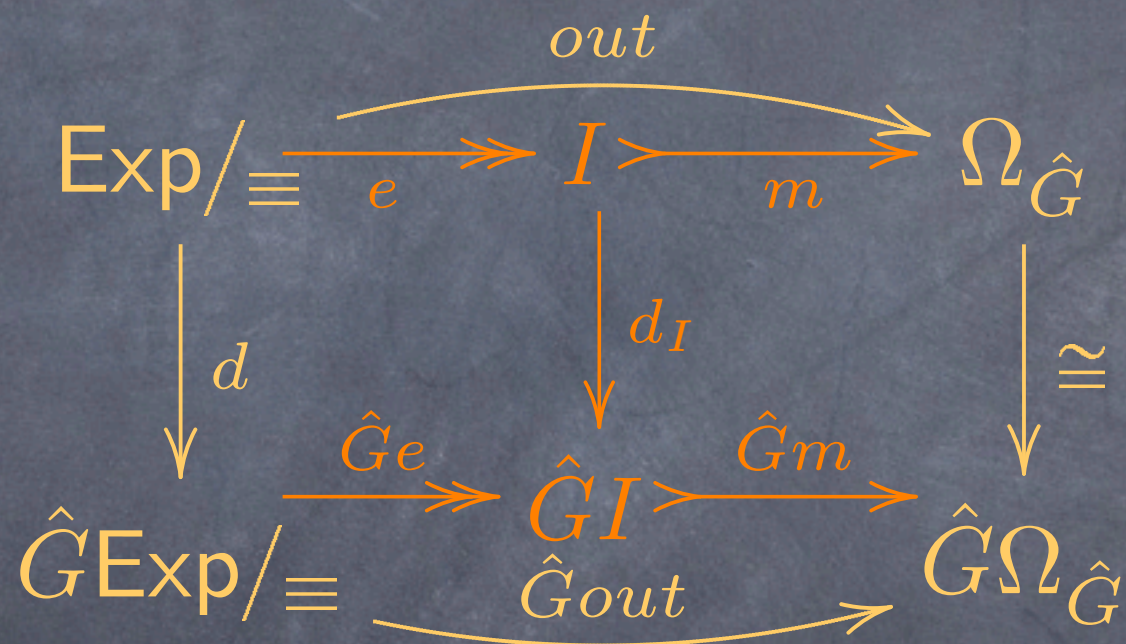
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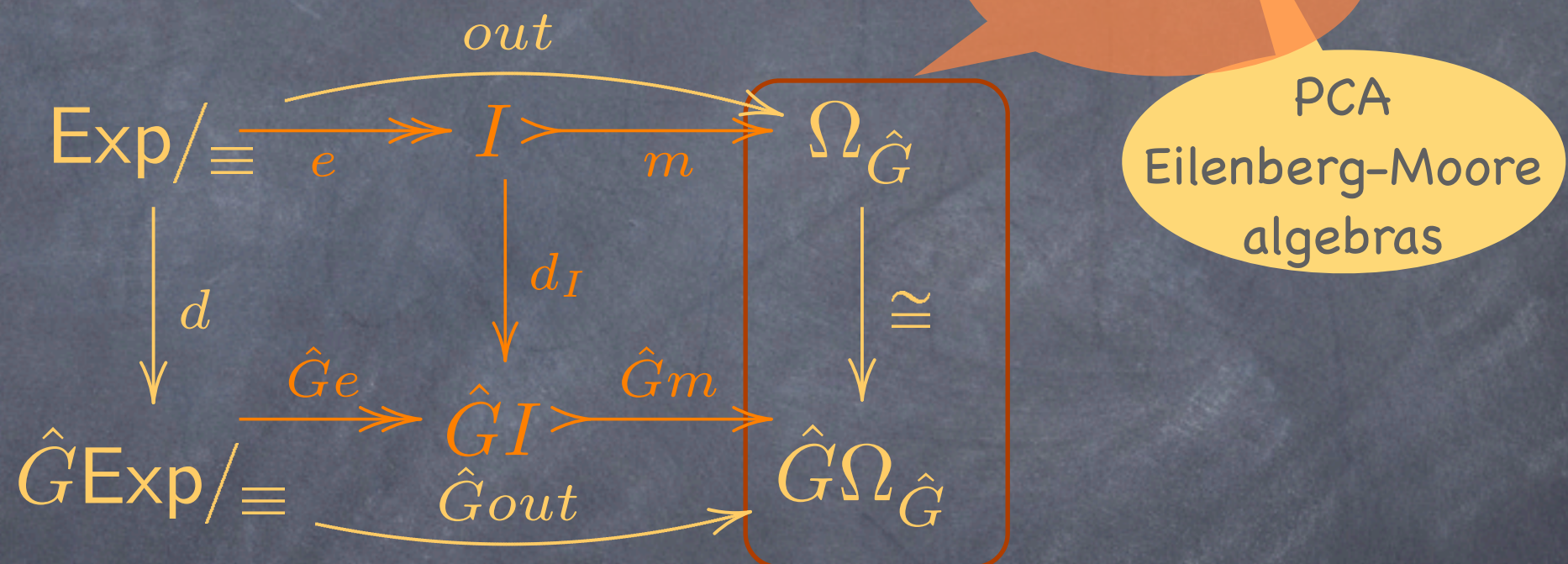
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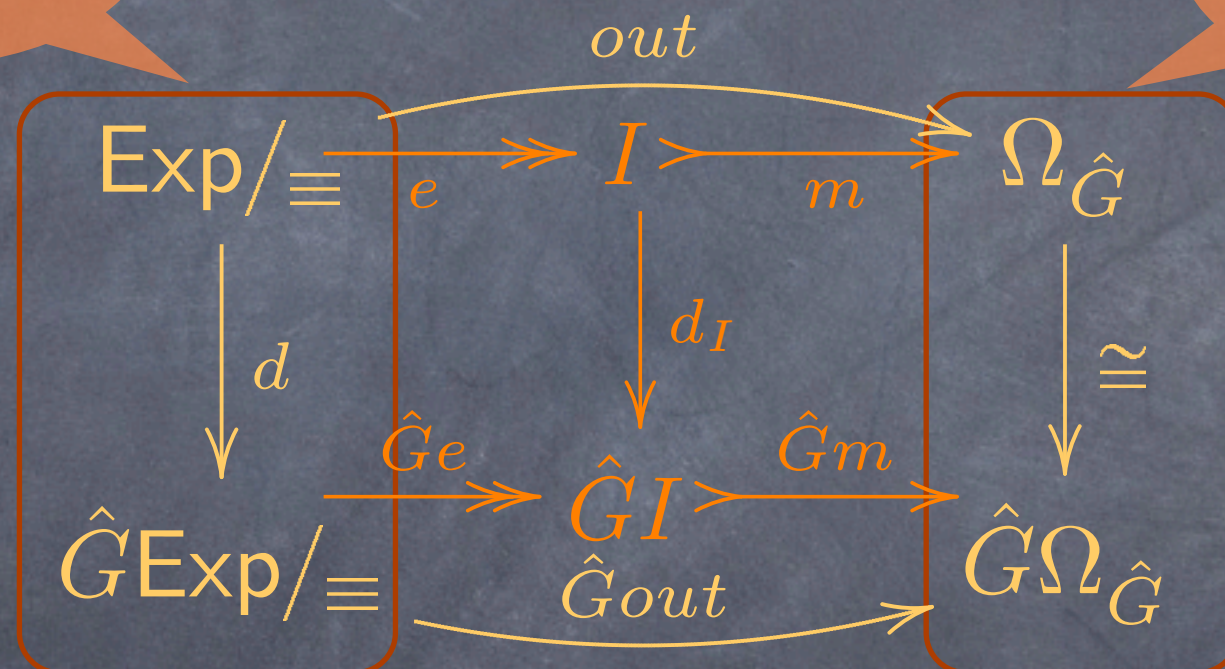
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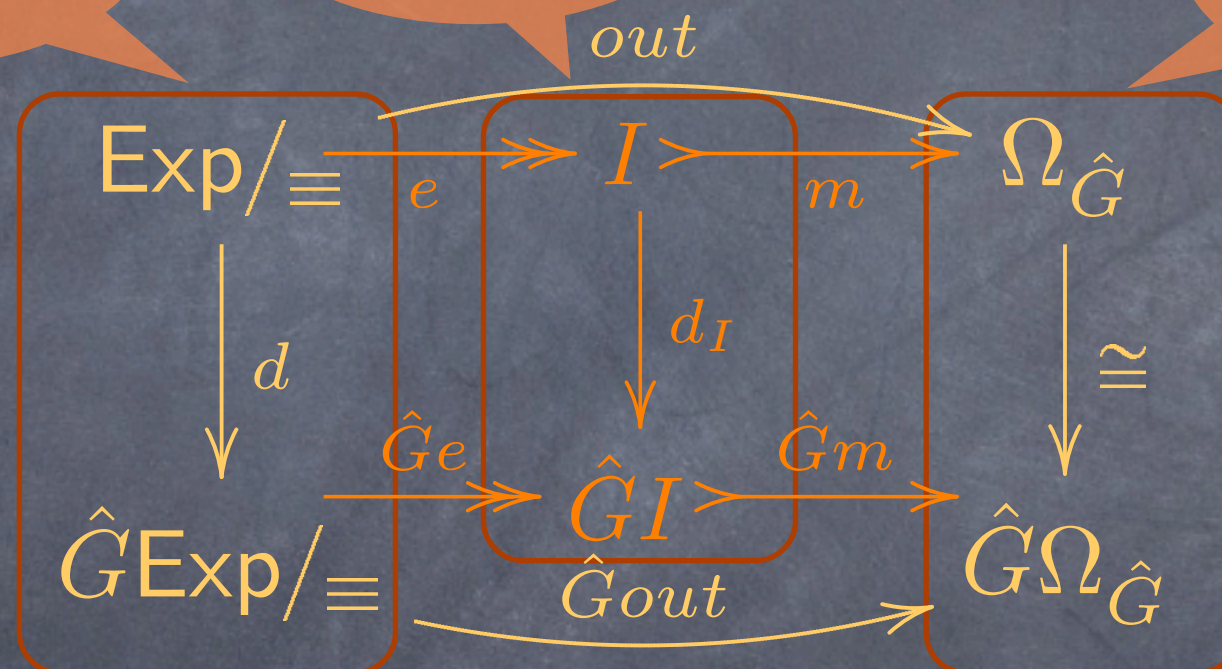
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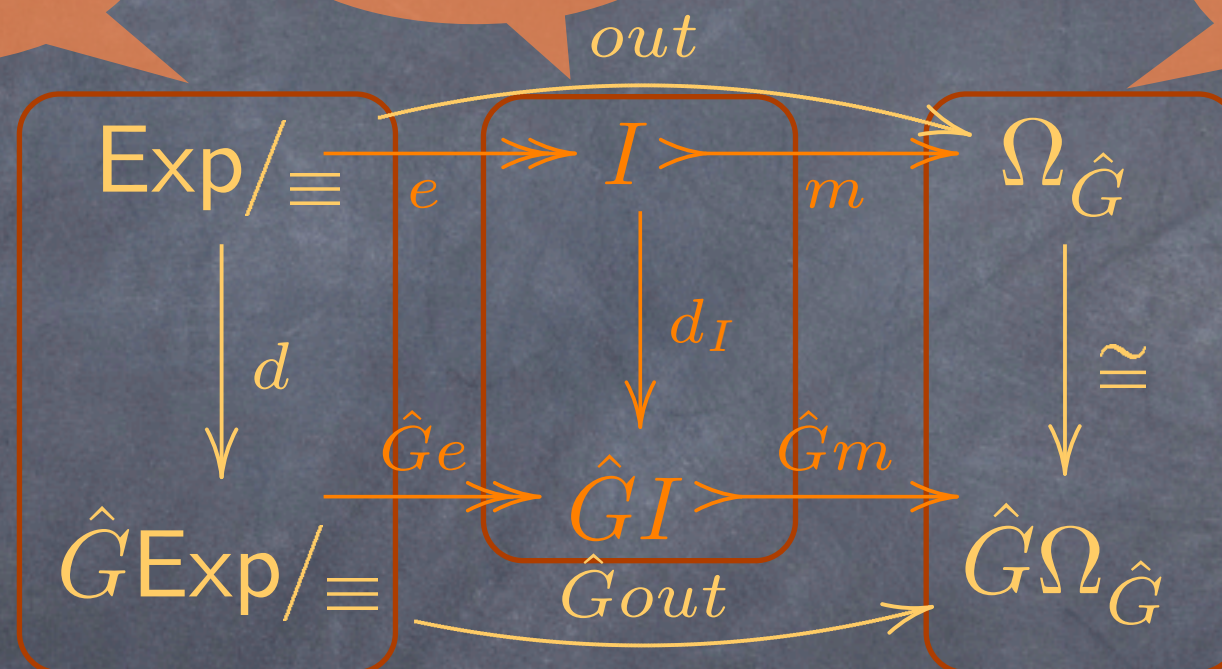
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Hence, e is iso, and out is injective

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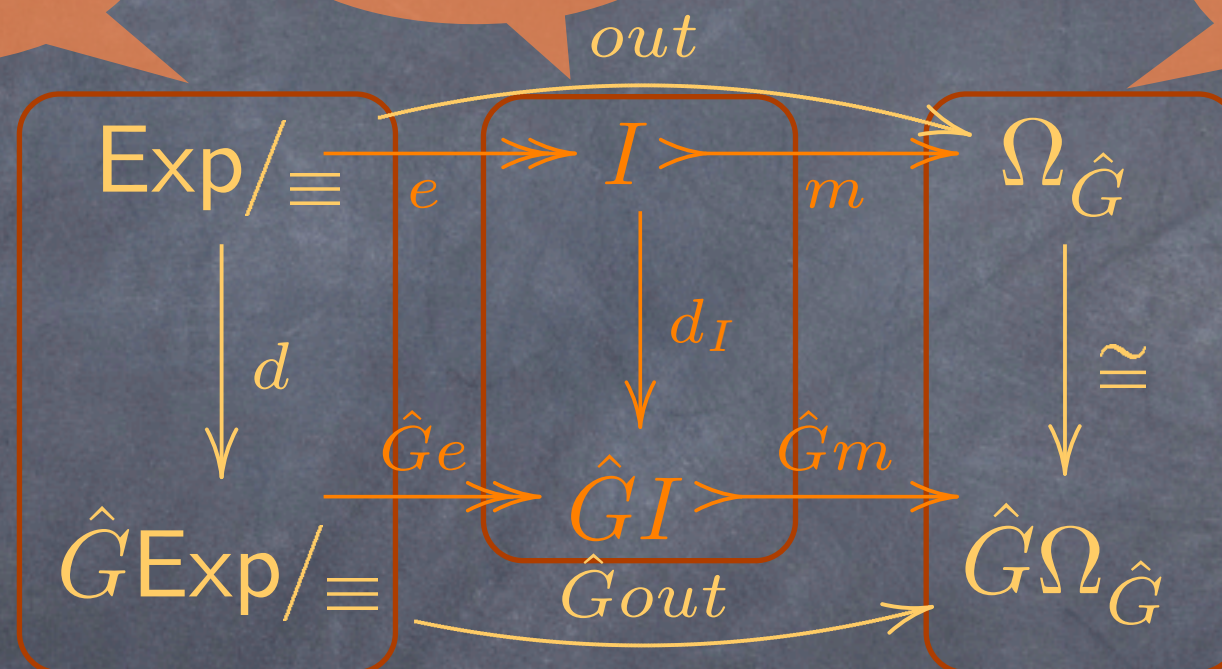
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Moreover: $tr = out \circ [-]$

Conclusions

- we present a solution to a concrete problem

sound and complete axiomatization
of traces for PTS

bisimilarity expressions and axioms plus one new axiom

- in a coalgebraic setting
- it opens many generalization questions...
- all about algebra and coalgebra

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Thank you !