The Power of Isolation



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Safety-critical Real-time Systems





Safety-critical Real-time Systems



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When things go wrong





When things go right...



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Important aspects for isolation: time space





Important aspects for isolation: **time space** Temporal isolation through CBS [Abeni04], VBS [Craciunas 12]

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Important aspects for isolation: **time space** Temporal isolation through CBS [Abeni04], VBS [Craciunas12] Spatial isolation through memory management/hardware



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Power isolation?



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Power isolation?

Is power consumption compositional?



Important aspects for isolation: **time space** Temporal isolation through CBS [Abeni04], VBS [Craciunas12] Spatial isolation through memory management/hardware

Power isolation?

Is power consumption compositional?

Problem: non-linear relationship of power consumption and processor frequency





Adapt system performance to system load





Adapt system performance to system load Dynamic Voltage and Frequency Scaling [Pillai01]



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Example: Intel X-Scale $p(f) = 1520f^3 + 80mWatt$ [Chen08]



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Scaling the frequency results in modified execution time
Deadlines remain the same

Minimize power while maintaining the real-time properties



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Minimize power while maintaining the real-time properties

 $f > \overline{U f_{max}}$ [Pillai01]



Approaches



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Measuring the power consumption [Pathak11]



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Measuring the power consumption [Pathak11] Controlling the power consumption [Cao08]



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Measuring is fine, but we want the information before system start



Approaches

Measuring the power consumption [Pathak11]

Controlling the power consumption [Cao08]

Measuring is fine, but we want the information before system start

Controlling is fine, but we do not want to interfere in the schedule

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Power model Ui System model





Power model Ui System model

Lower and upper bounds on the power consumption of a task as functions of task utilization, frequency scaling, and power model.





Study the compositionality of power consumption (power isolation)




Our approach

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Isolate power consumption through over-provisioning



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Relationship between the power consumption and the contribution of a single task to this power consumption + the trade-off between quality and cost of power isolation.



Our approach

Study the compositionality of power consumption (power isolation)

Isolate power consumption through over-provisioning

Relationship between the power consumption and the contribution of a single task to this power consumption + the trade-off between quality and cost of power isolation.

We discuss the variance between lower and upper bounds (quality) and the power consumption overhead (cost) of power isolation.





CPU energy consumption of a EDF system with utilization $U = \sum_{i=1}^{n} U_i$ running at frequency $\kappa f_{max}, U \leq \kappa \leq 1$ in the interval $[t_0, t_1)$



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$$t_{idle}c_0 + t_{running}(c_0 + c_1(\kappa f_{max})^{\omega})$$



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$$E(\kappa, U) = (t_1 - t_0)c_1 \frac{U}{\kappa} (\kappa f_{max})^{\omega}$$











































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Two frequency levels

$E(1,U) = (t_1 - t_0)c_1 U f_{max}^{\omega}$

 f_{max}



 f_{max}



Jmax

 $bE_{i}^{u} = bE_{i}^{l} = E(1, U_{i}) = (t_{1} - t_{0})c_{1}U_{i}f_{max}^{\omega}$



Jmax

 $bE_{i}^{u} = bE_{i}^{l} = E(1, U_{i}) = (t_{1} - t_{0})c_{1}U_{i}f_{max}^{\omega}$

$$E\left(1,\sum_{i=1}^{n}U_{i}\right) = \sum_{i=1}^{n}E(1,U_{i})$$



Jmax

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Continuous frequency levels

 f_{max}


































task i is the only running task





Continuous frequency levels



Continuous frequency levels $E(U, U) = (t_1 - t_0)c_1(f_{max}U)^{\omega}.$



Continuous frequency levels $E(U,U) = (t_1 - t_0)c_1(f_{max}U)^{\omega}.$ $bE_i^l = E(U_i,U_i) = (t_1 - t_0)c_1(f_{max}U_i)^{\omega}$







Continuous frequency levels $E(U, U) = (t_1 - t_0)c_1(f_{max}U)^{\omega}.$ $bE_{i}^{l} = E(U_{i}, U_{i}) = (t_{1} - t_{0})c_{1}(f_{max}U_{i})^{\omega}$ $bE_i^u = E(1,1) - E(1 - U_i, 1 - U_i)$ $= (t_1 - t_0)c_1 f_{max}^{\omega} (1 - (1 - U_i)^{\omega})$

 $t_0 = 0s, t_1 = 10s, c_1 = 1520$ mWatt



































Minimum contribution of task i to the power consumption





Minimum contribution of task i to the power consumption Maximum contribution of task i to the power consumption









Task i is the only task in the system





Task i is the only task in the system $bE_{i}^{l} = (t_{1} - t_{0})c_{1}\frac{U_{i}}{\kappa_{l}}(\kappa_{l}f_{max})^{\omega}$





Task i is the only task in th $\kappa_{l-1} < U_i \le \kappa_l$ $bE_i^l = (t_1 - t_0)c_1 \frac{U_i}{\kappa_l} (\kappa_l f_{max})^{\omega}$









 $E(1,\min(\kappa+U_i,1)) - E(\kappa,\min(\kappa_{u-1},\kappa_{m-2})))$





 $E(1,\min(\kappa+U_i,1)) - E(\kappa,\min(\kappa_{u-1},\kappa_{m-2})))$

$$\kappa = \min(\kappa_u, \kappa_{m-1})$$





 $E(1,\min(\kappa+U_i,1)) - E(\kappa,\min(\kappa_{u-1},\kappa_{m-2})))$

$$\kappa = \min(\kappa_u, \kappa_{m-1})$$

$$-1 < 1 - U_i \le \kappa_u$$

 κ_u



Discrete Frequencies - Bounds

Jmax



$$bE_i^l = (t_1 - t_0)c_1 \frac{U_i}{\kappa_l} (\kappa_l f_{max})^{\omega}$$



 $\left(\right)$

3 discrete frequencies - bounds



fmax





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3 discrete frequencies - jitter

 f_{max}





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Discrete Frequencies - Jitter



 f_{max}







Difference between the optimal and actual power consumption



Difference between the optimal and actual power consumption Optimal is achieved with infinite frequency levels



Difference between the optimal and actual power consumption

Optimal is achieved with infinite frequency levels

The cost of power isolation depends on
the sum of task utilizations
the utilization of the considered task
the number of available frequency levels
the distribution of the frequency levels in the interval [0, fmax]



Difference between the optimal and actual power consumption

Optimal is achieved with infinite frequency levels

The cost of power isolation depends on
the sum of task utilizations
the utilization of the considered task
the number of available frequency levels
the distribution of the frequency levels in the interval [0, fmax]

$$cE_i = (t_1 - t_0)c_1 f_{max}^{\omega} U_i (1 - U_i^{\omega - 1})$$

Cost of power isolation








Lower and upper bounds on the individual power consumption





Lower and upper bounds on the individual power consumption

Quality of power isolation



- Lower and upper bounds on the individual power consumption
- Quality of power isolation
- Cost of power isolation



- Lower and upper bounds on the individual power consumption
- Quality of power isolation
- Cost of power isolation

Full temporal [Craciunas 12], spatial [Craciunas08], and power isolation of tasks



- Lower and upper bounds on the individual power consumption
- Quality of power isolation
- Cost of power isolation

Full temporal [Craciunas 12], spatial [Craciunas08], and power isolation of tasks

The key insight is that there appears to be a fundamental trade-off between quality and cost of time, space, and power isolation



Thank you

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