Explicit, Dynamic Memory Management with Temporal and Spatial Guarantees

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Memory Management

- Allocation:
  - `malloc`

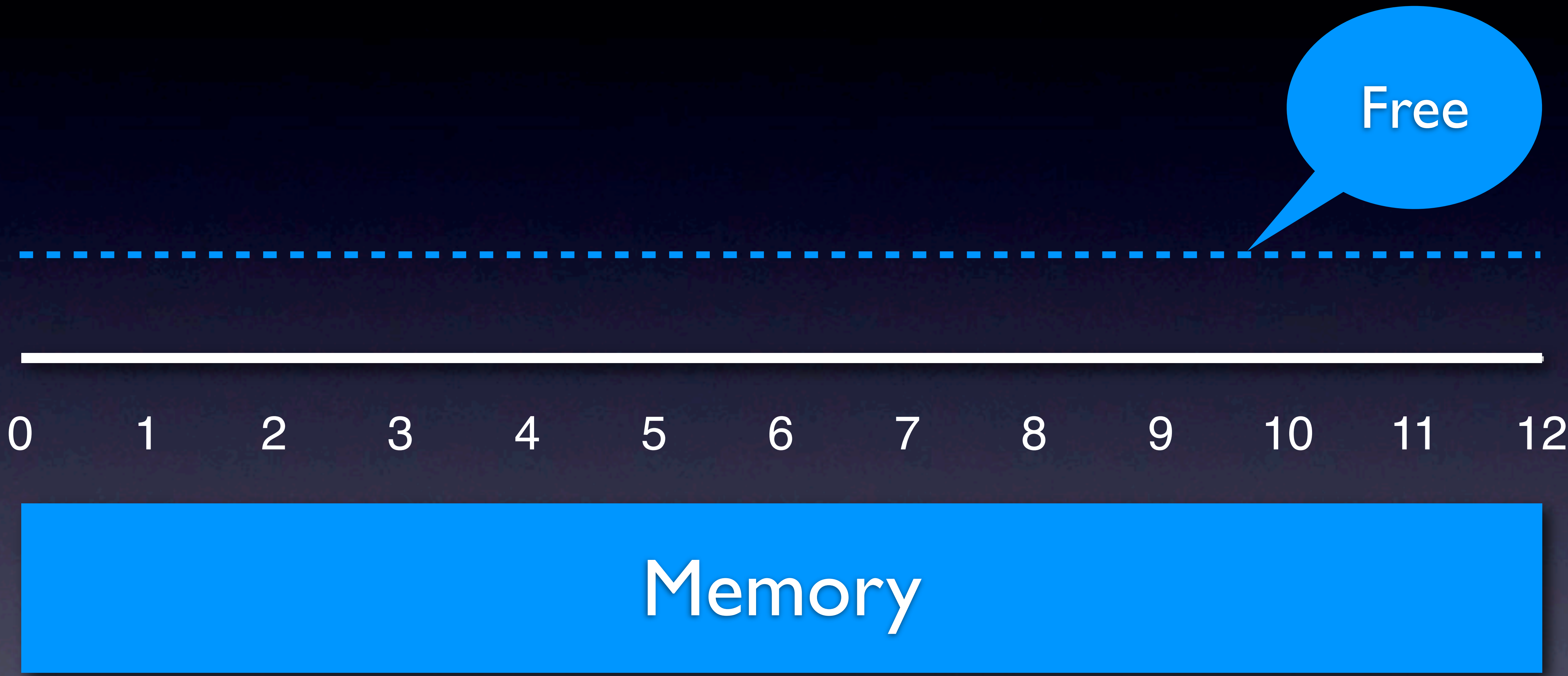
Memory Management

- Allocation:
  - `malloc`
- Deallocation:
  - `free`

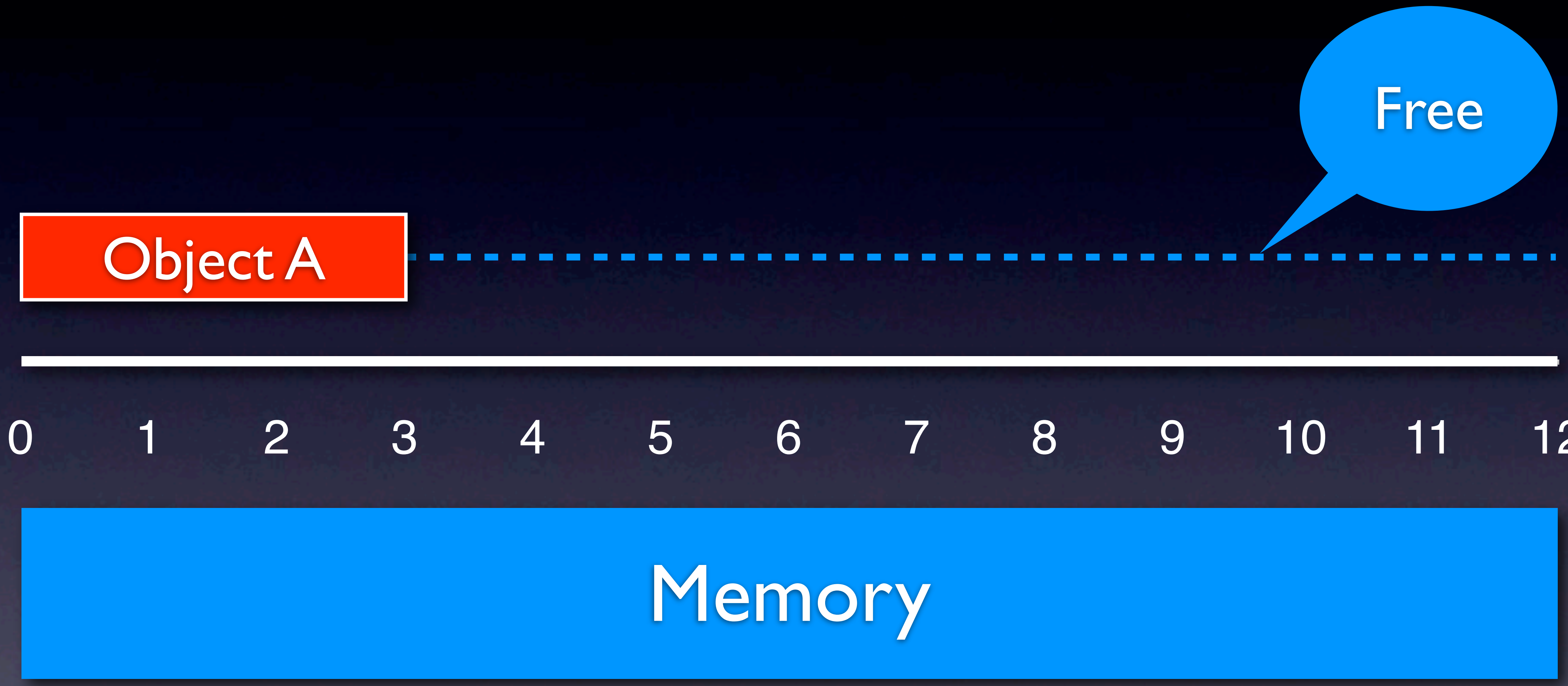
Memory Management

- Allocation:
  - `malloc`
- Deallocation:
  - `free`
- Access:
  - `read` and `write`

Allocation



Allocation



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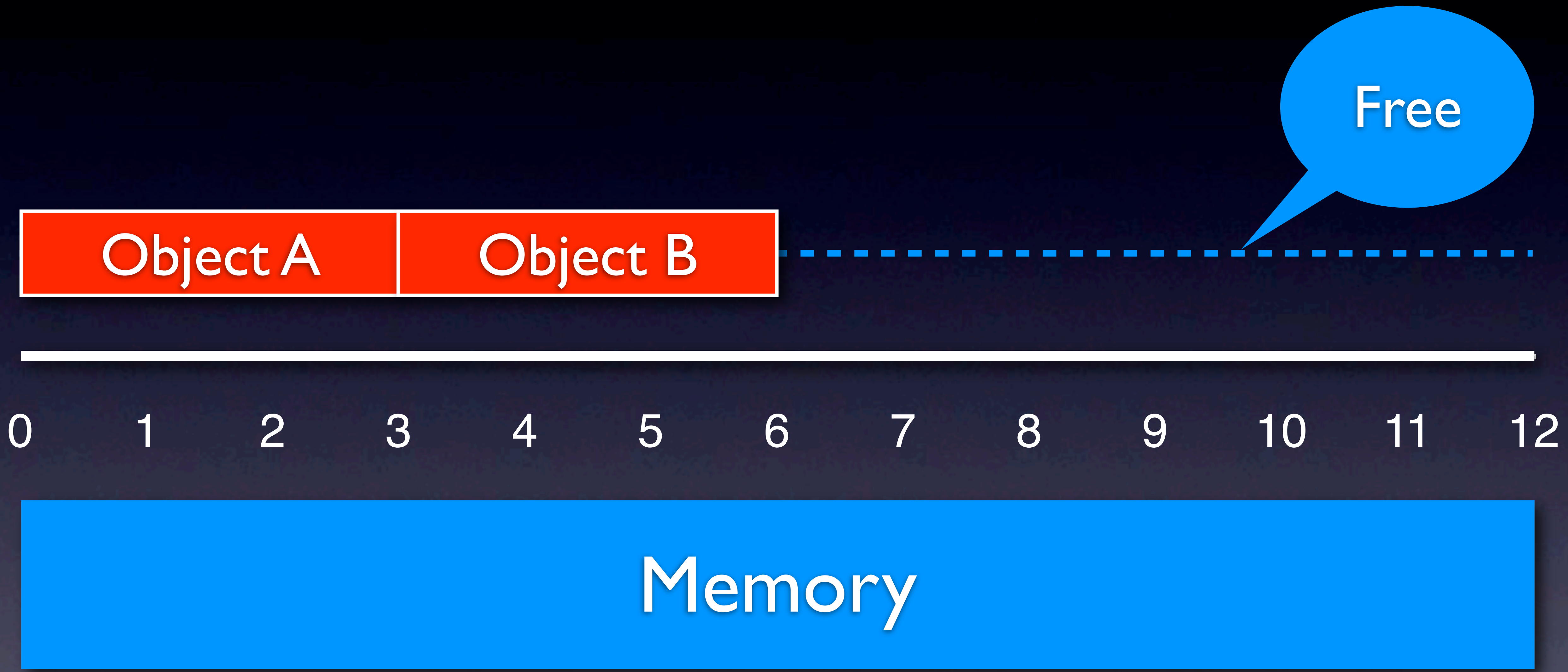
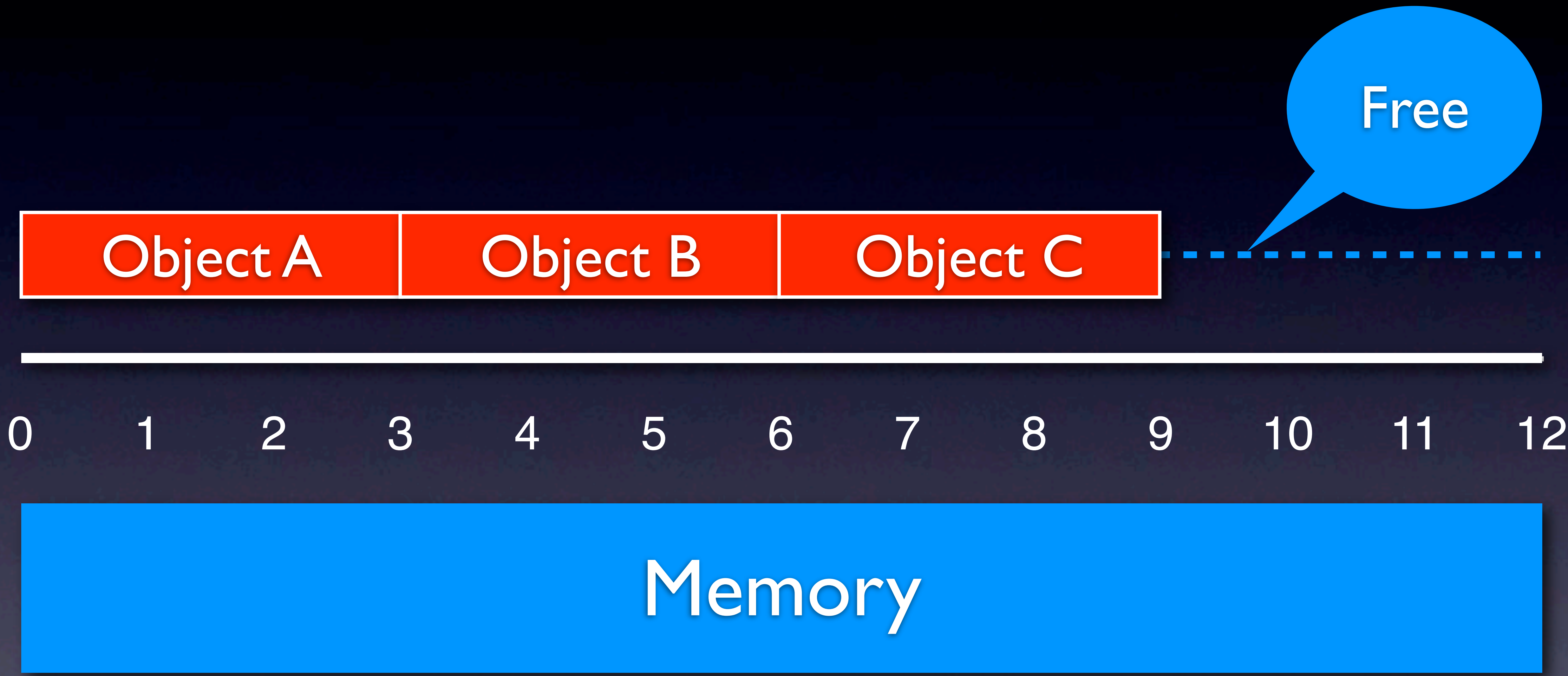


Allocation



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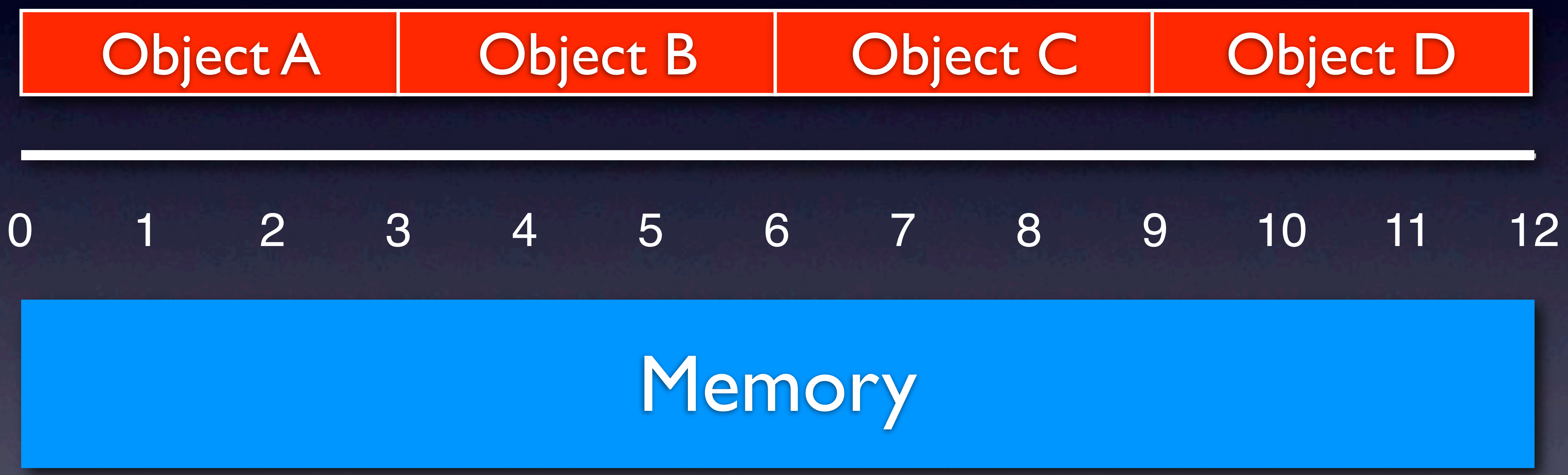
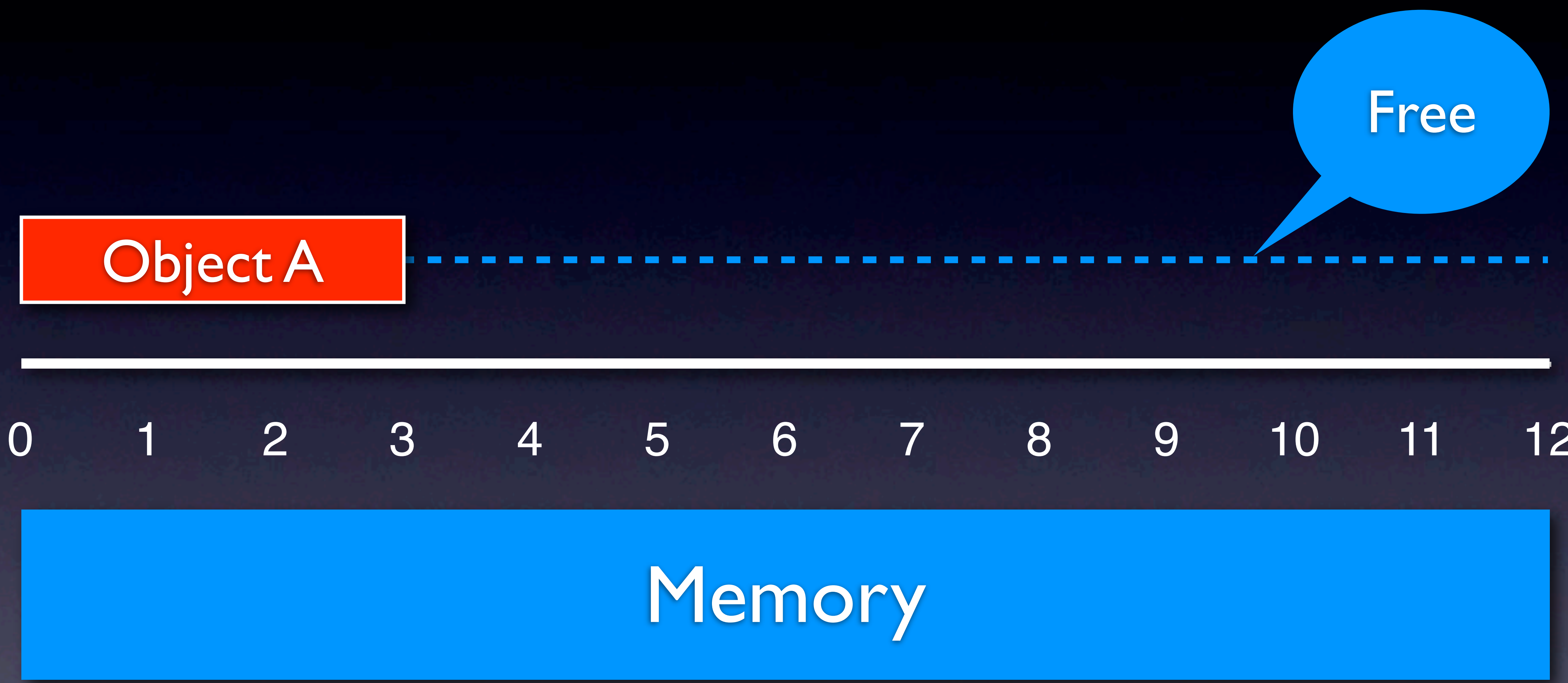

Object A
Object B
Object C
Object D
0 1 2 3 4 5 6 7 8 9 10 11 12
Memory

1. Assumption:

Objects may have different sizes

Allocation



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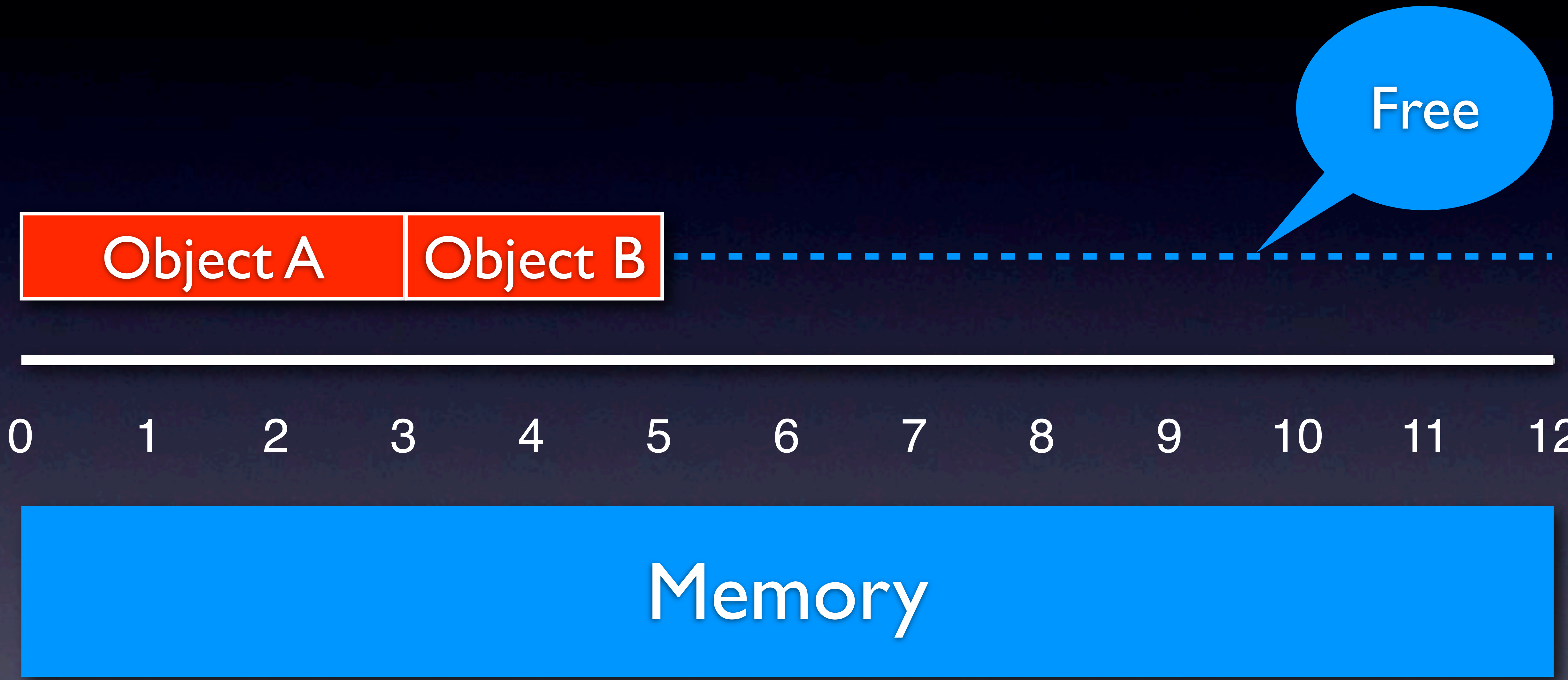

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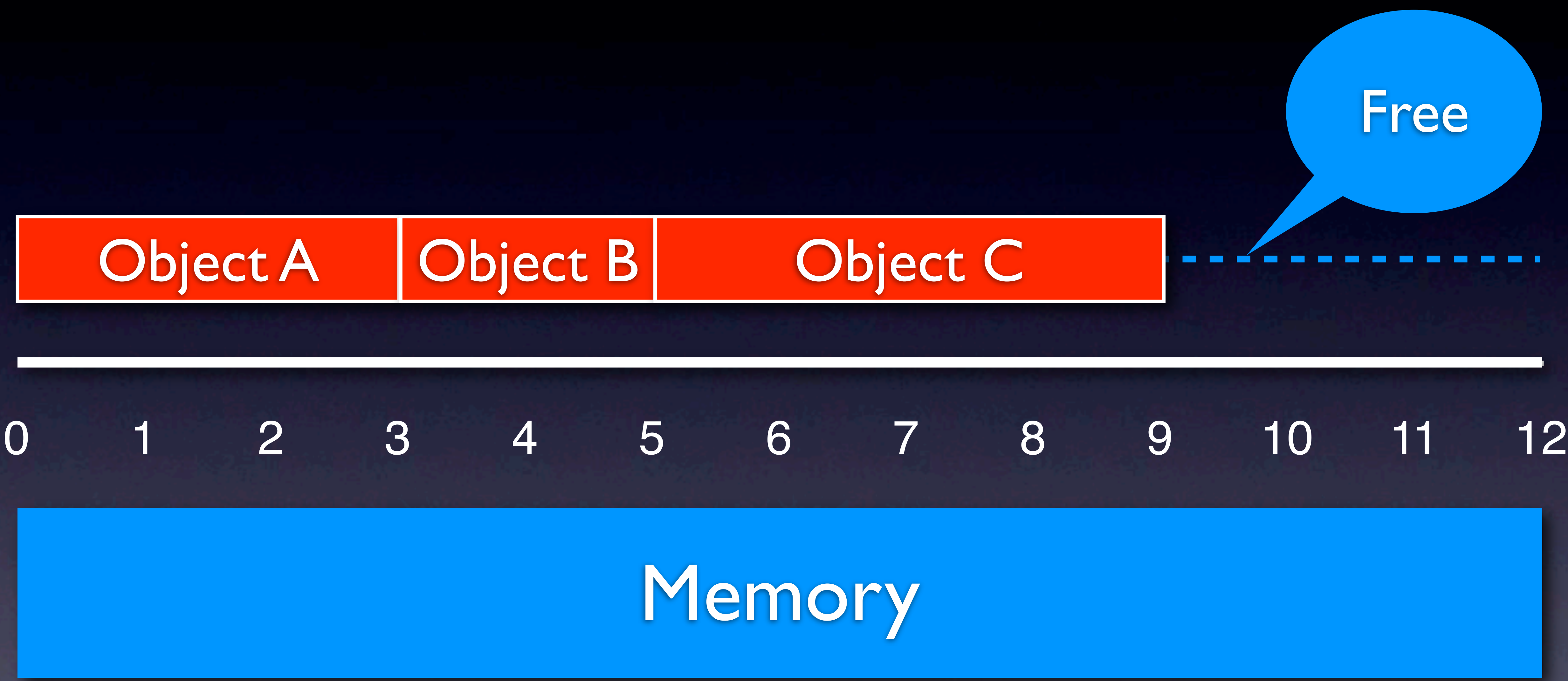

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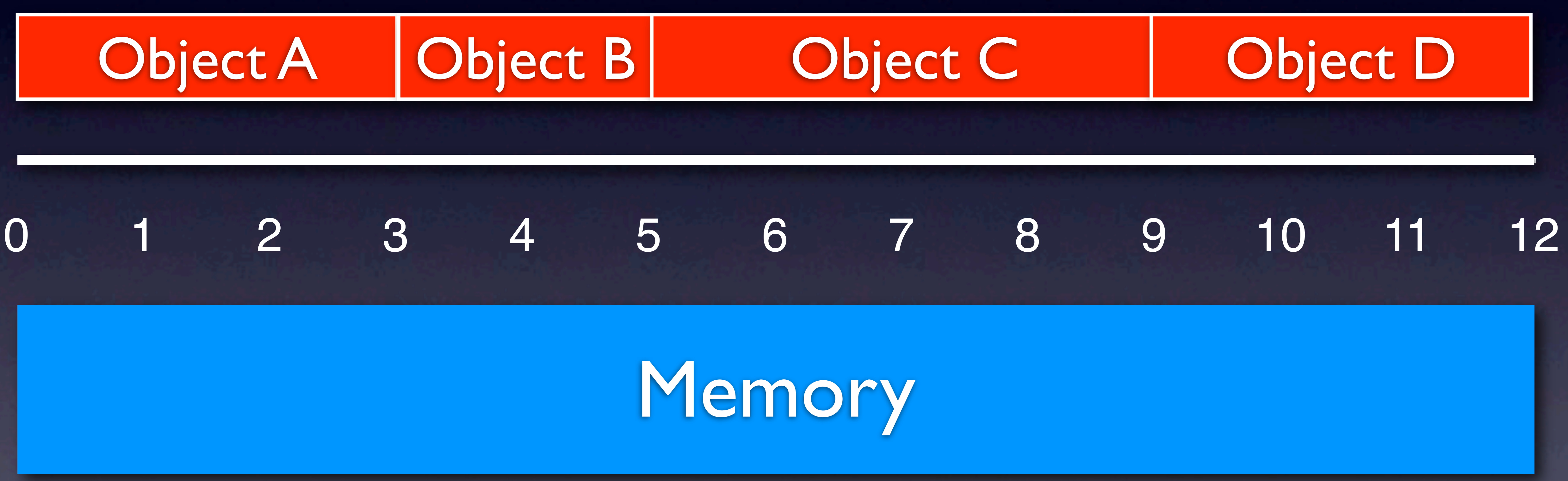
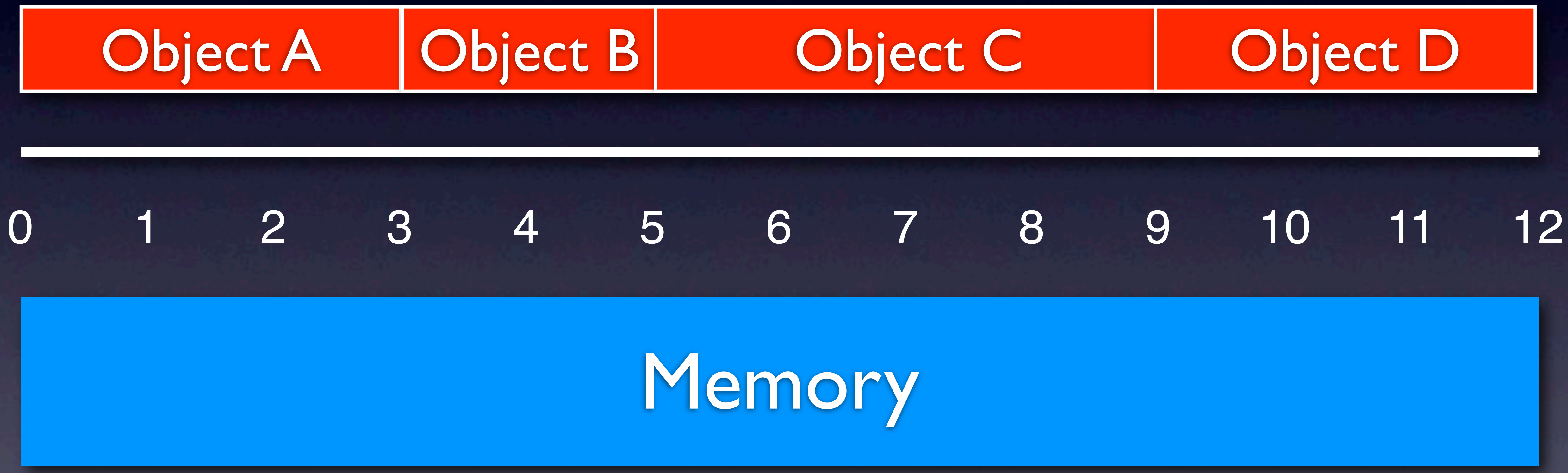
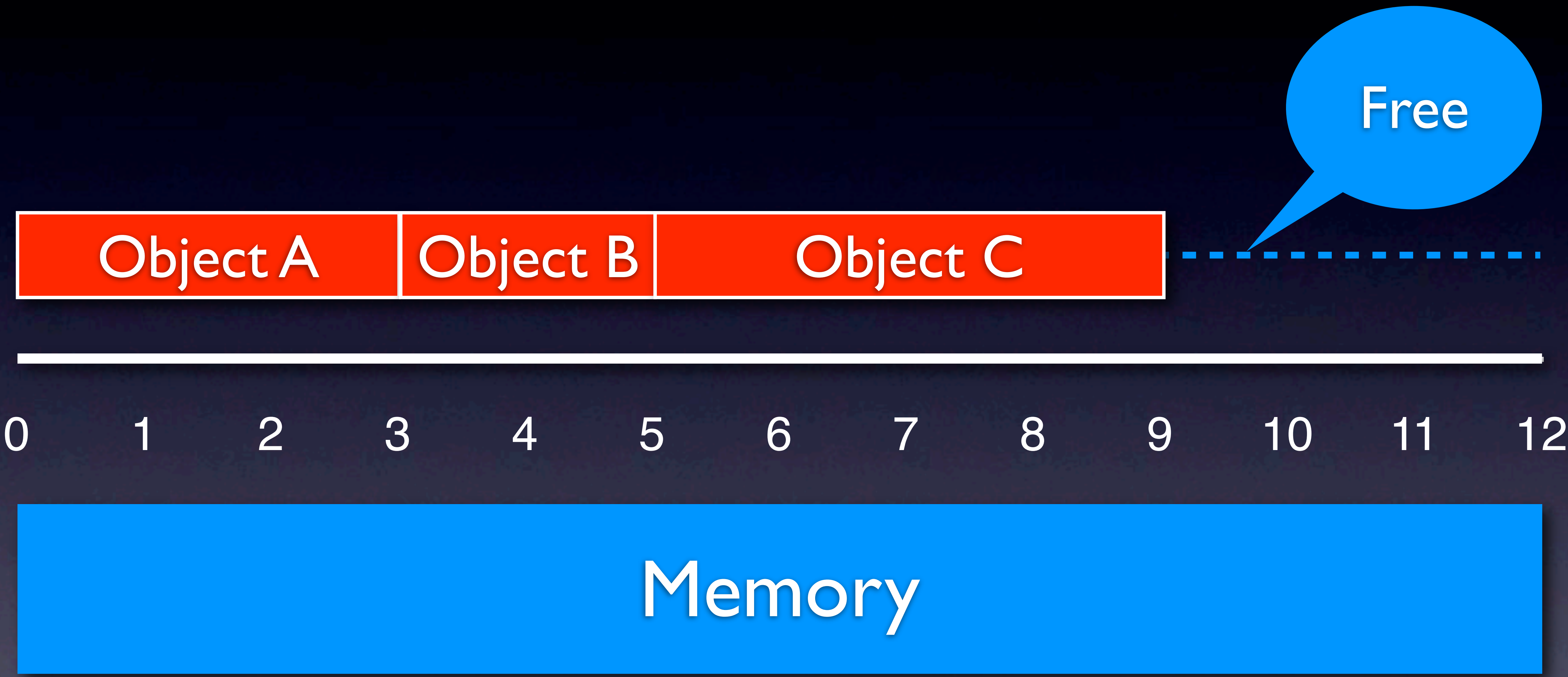

Object A
Object B
Object C
Object D
0 1 2 3 4 5 6 7 8 9 10 11 12
Memory

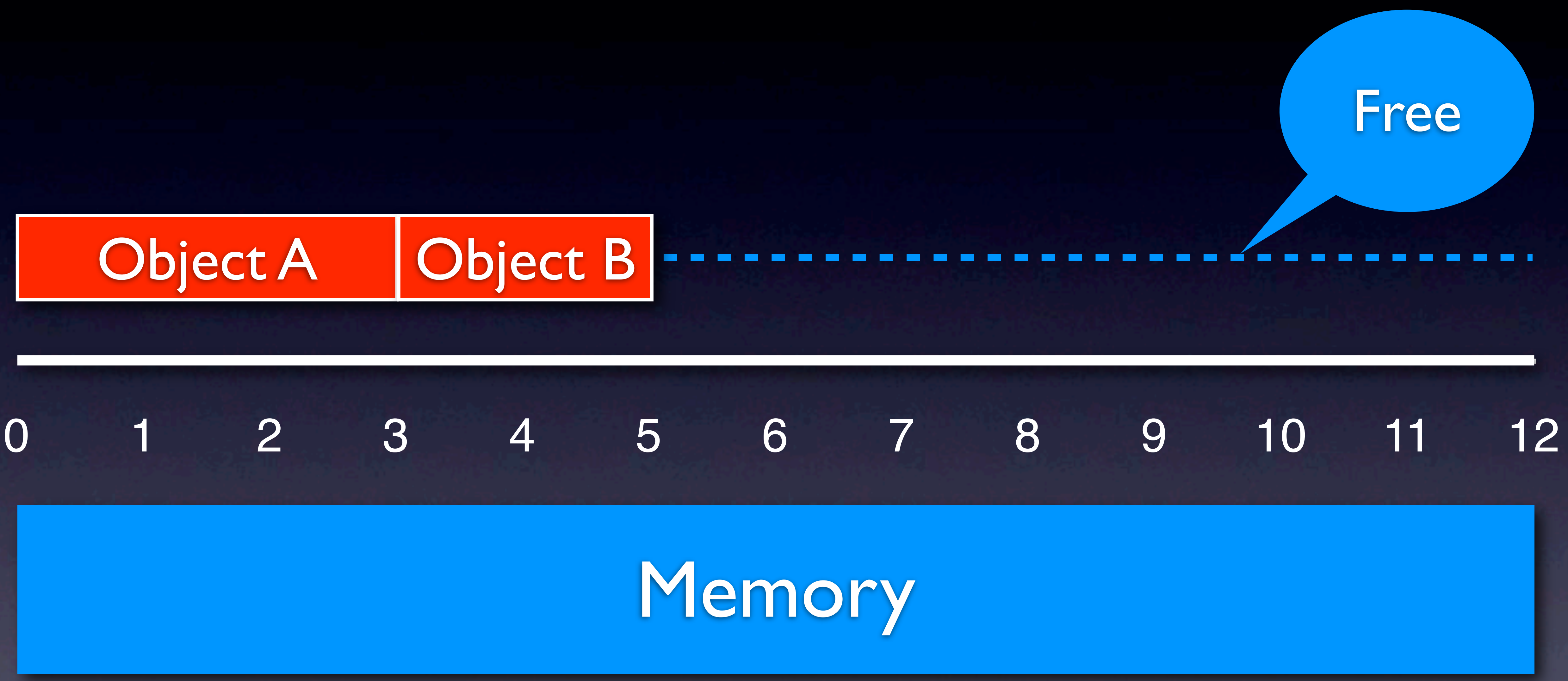
Deallocation



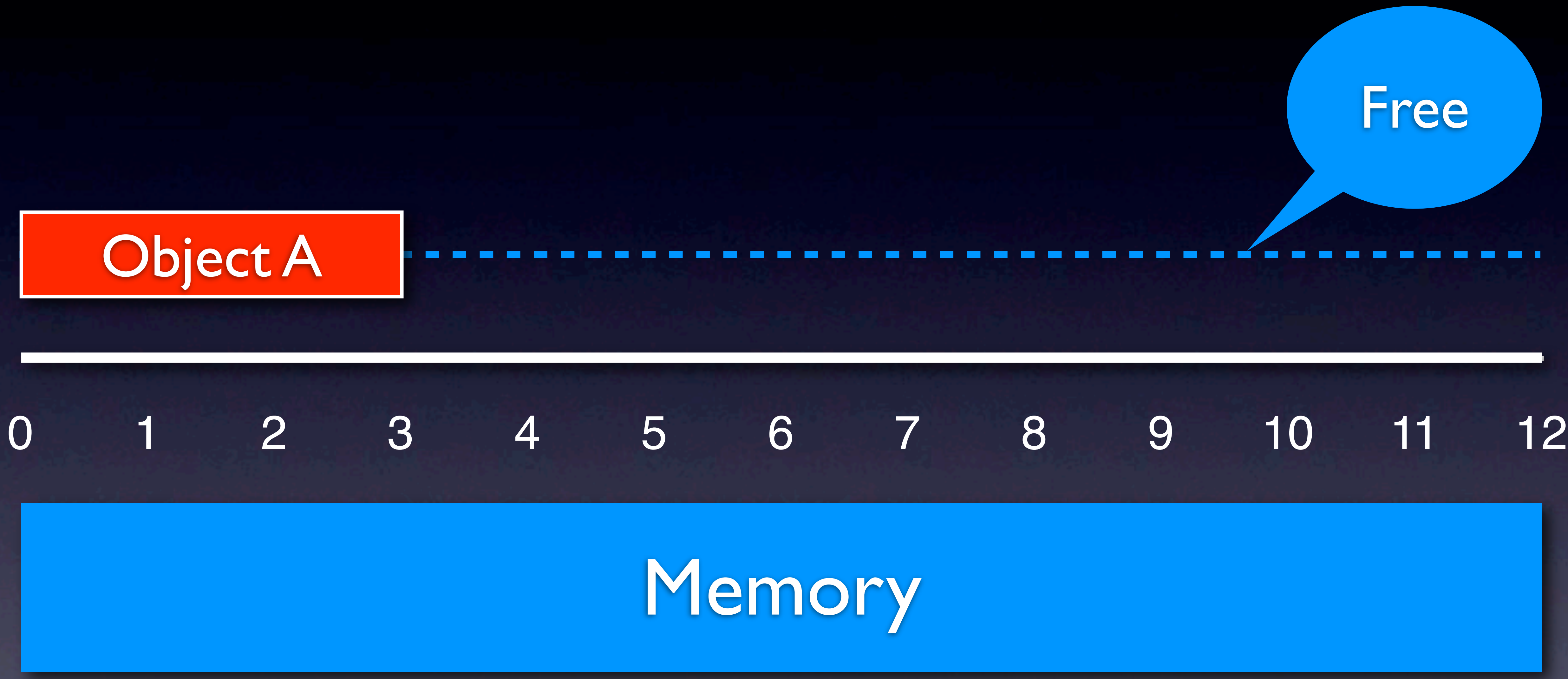
Deallocation



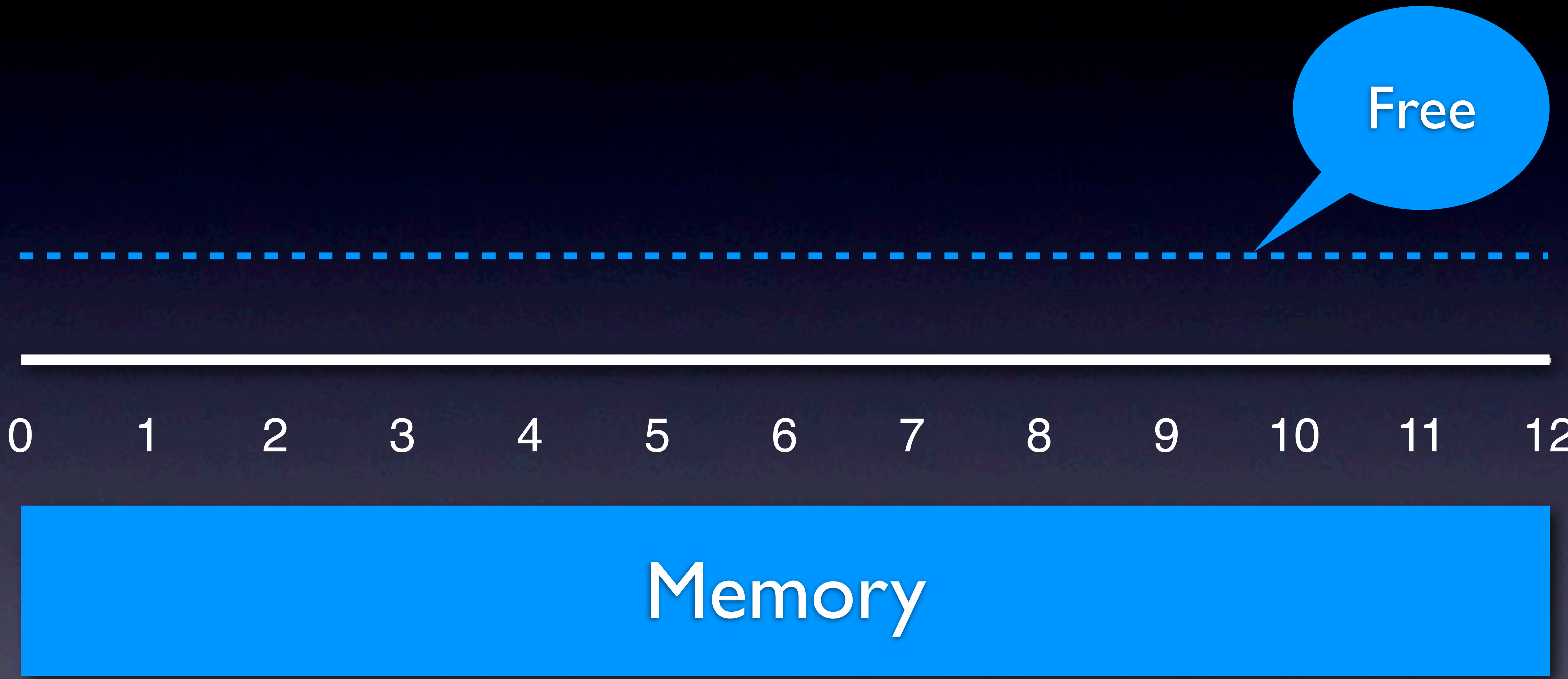
Deallocation



Deallocation



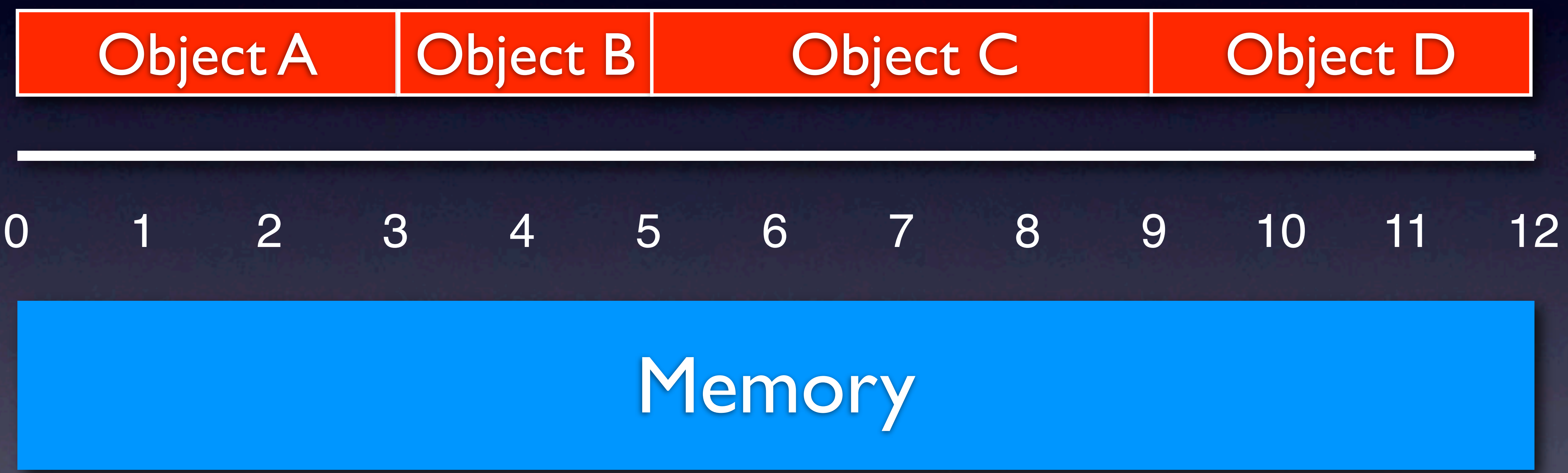
Deallocation



2. Assumption:

Objects may be allocated and deallocated in random order

Deallocation



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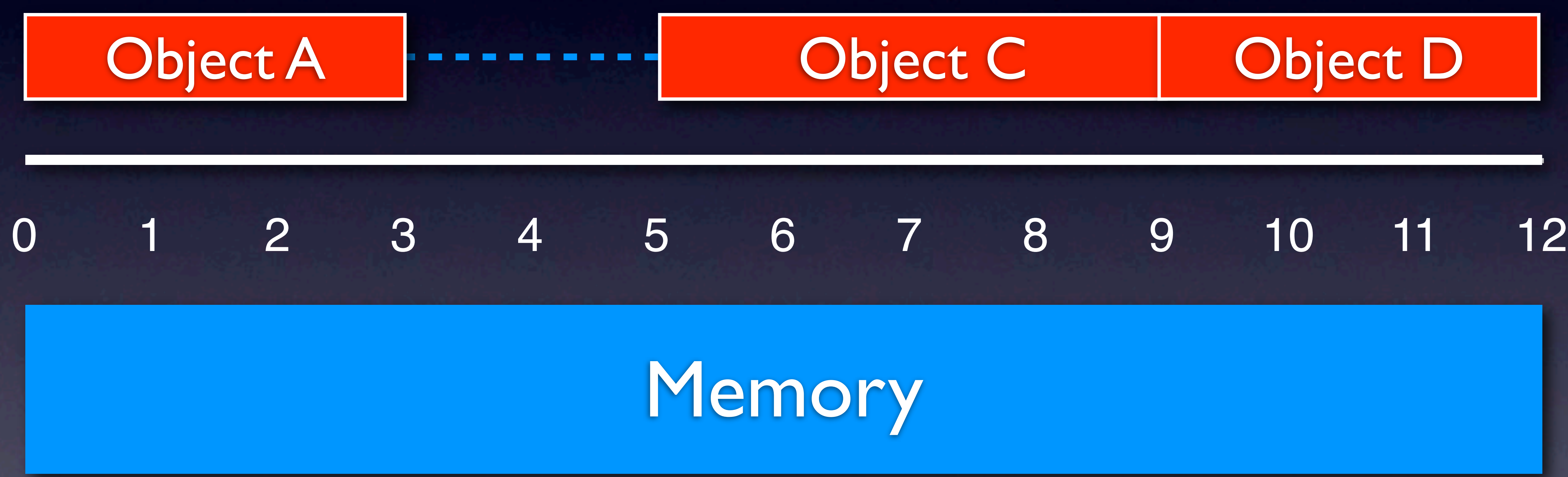
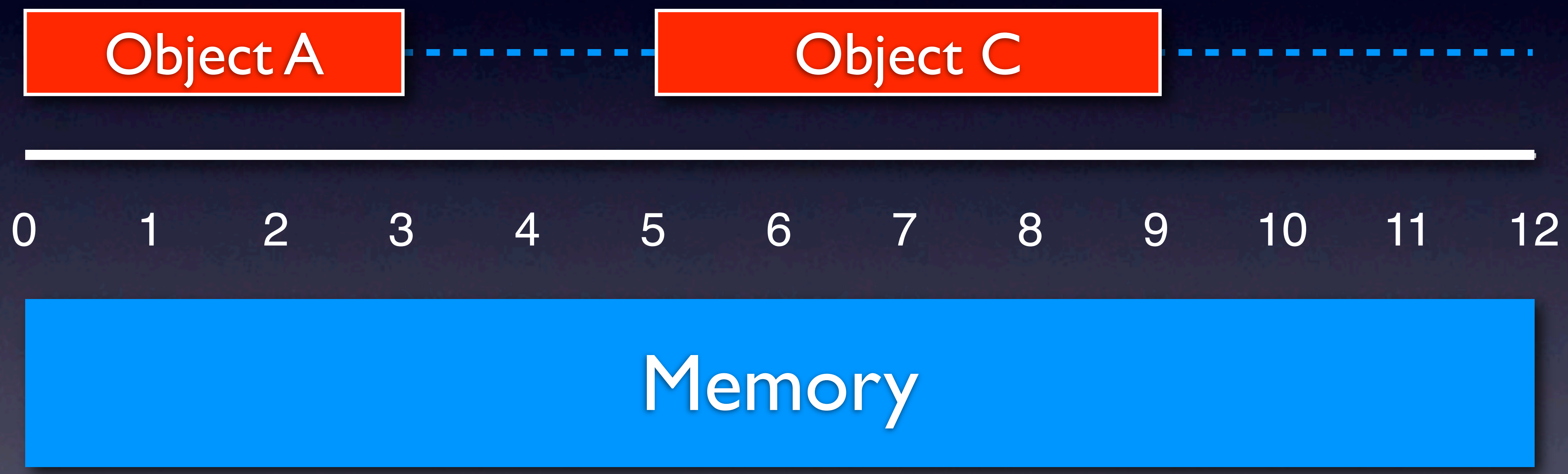
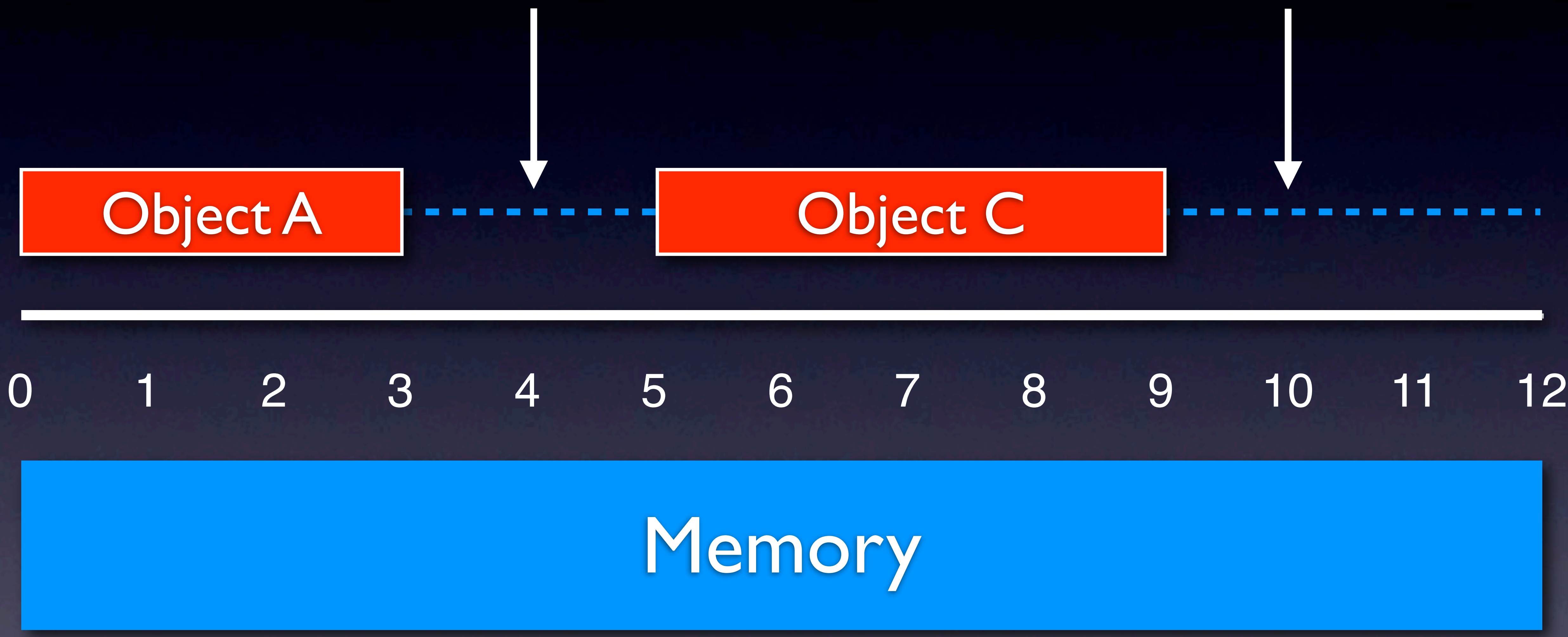


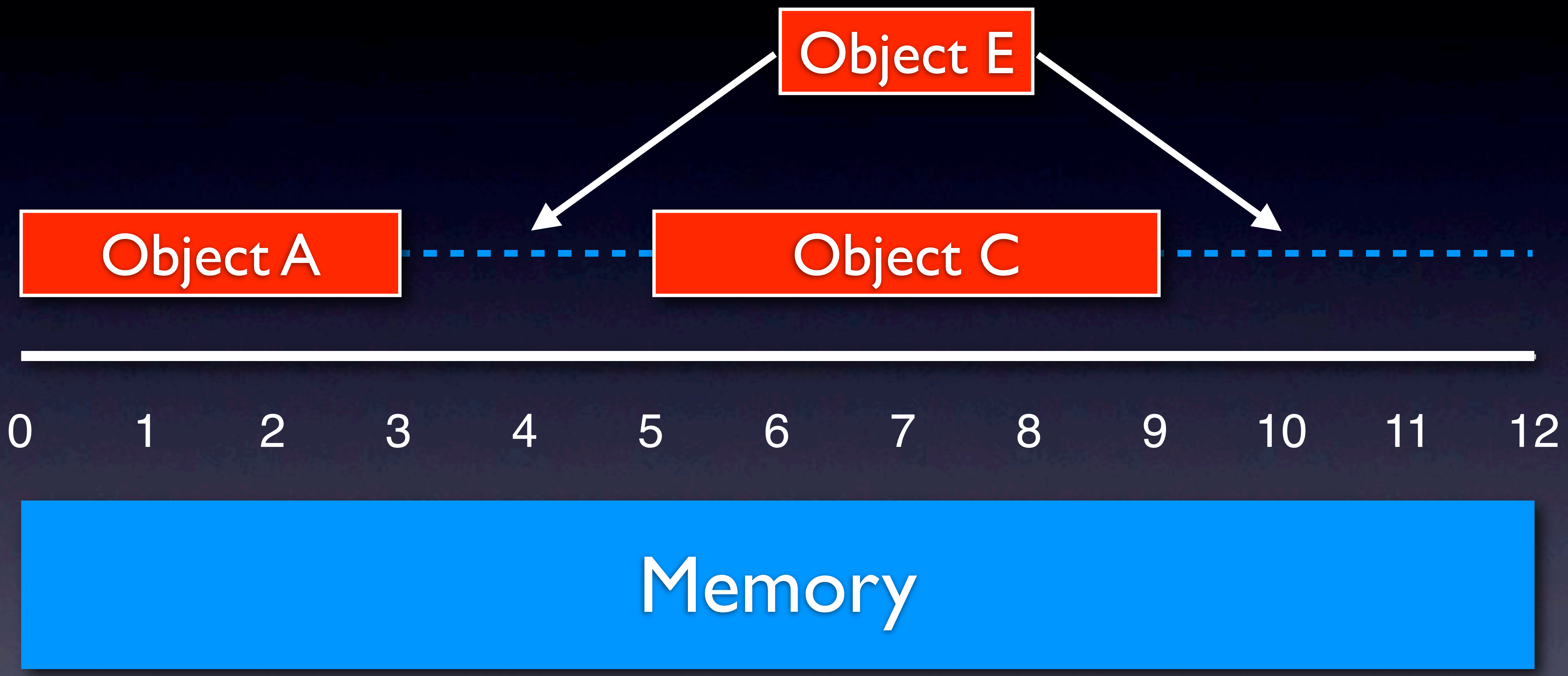
Deallocation



External Fragmentation



Allocation



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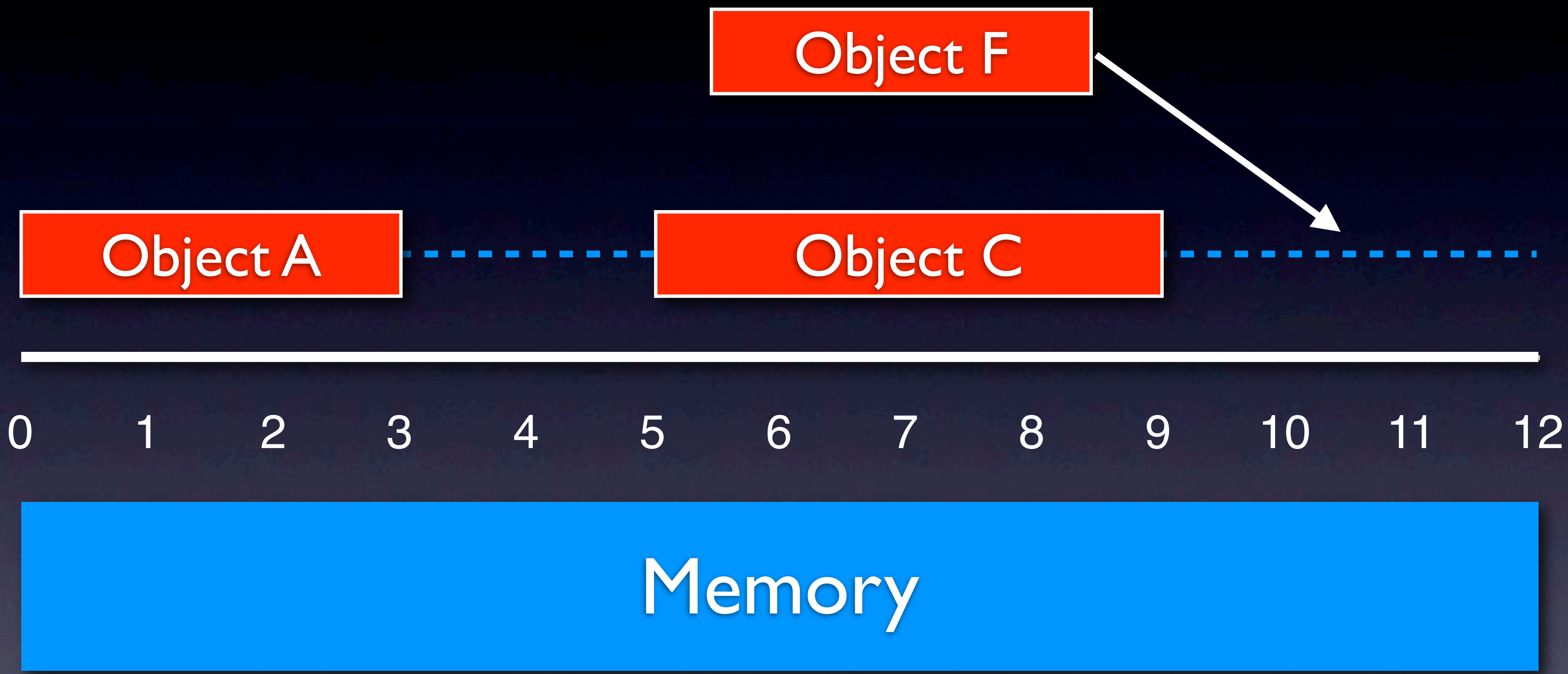

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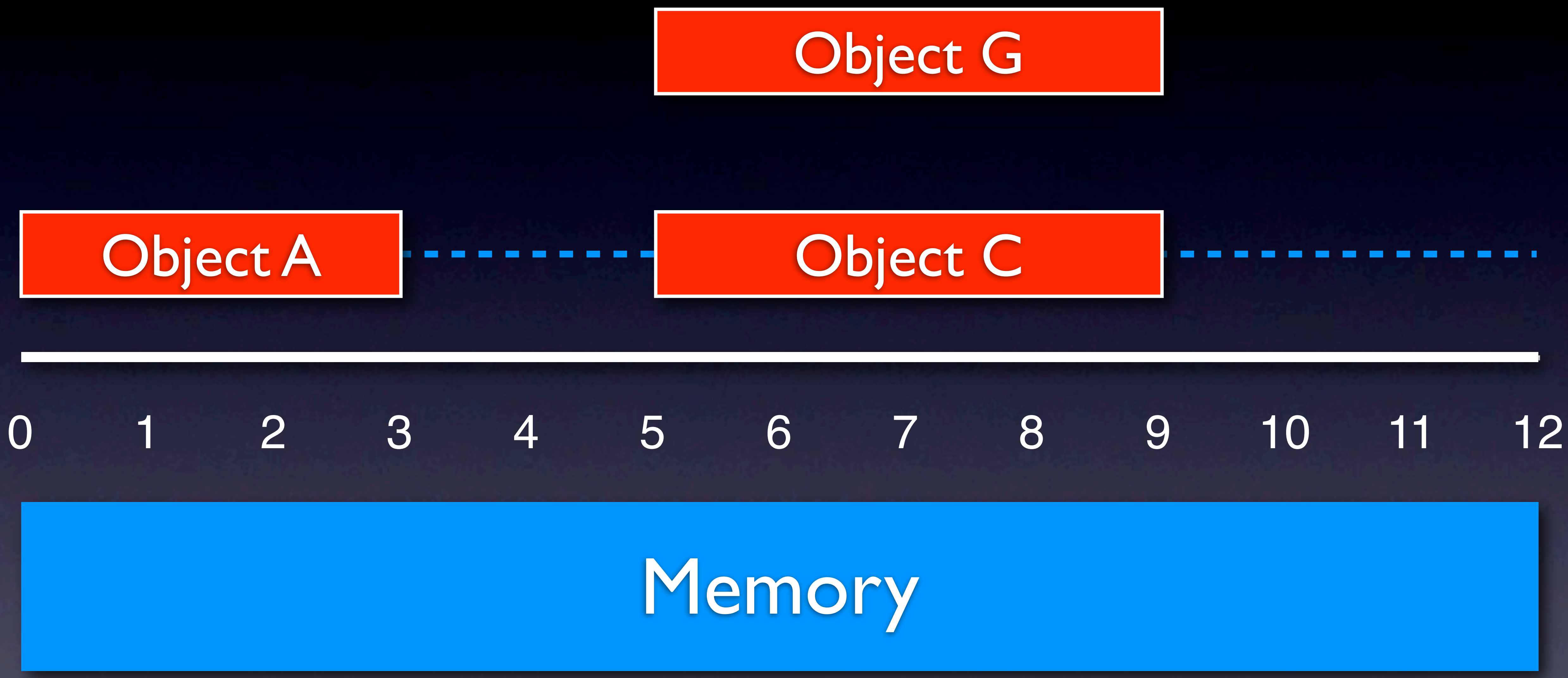

Memory is *fragmented* if the largest, contiguous piece of available space

is

smaller than

the total available space

Fragmentation

- Memory objects may have different sizes
- Memory objects may be allocated and deallocated in random order
  - creates the problem of memory fragmentation!

Explicit, Dynamic Memory Management with Temporal and Spatial Guarantees

Static versus Dynamic

- Static memory management:
  - Preallocate all memory at compile time

Static versus Dynamic

- Static memory management:
  - Preallocate all memory at compile time
- Dynamic memory management:
  - Allocate and deallocate memory at run time

Explicit, Dynamic Memory Management with Temporal and Spatial Guarantees

Implicit versus Explicit

- Implicit, dynamic memory management:
  - Garbage collector (GC) deallocates objects, not programmer (implicit `free` calls by GC)

Implicit versus Explicit

- Implicit, dynamic memory management:
  - Garbage collector (GC) deallocates objects, not programmer (implicit `free` calls by GC)
- Explicit, dynamic memory management:
  - Objects are deallocated by programmer (explicit `free` calls)

Programming Abstraction

Runtime Overhead

Implicit, Dynamic Memory Management

Explicit, Dynamic Memory Management

Static Memory Management

Programming Abstraction | Runtime Overhead

Implicit | Web, Safety

Explicit | Server, Performance

Static | Embedded, Real Time

Programming Abstraction

Runtime Overhead

Implicit	Web, Safety
Explicit	Server, Performance
Static	Embedded, Real Time

Programming Abstraction

Runtime Overhead

Implicit — Web, Safety

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Static — Embedded, Real Time

Temporal Performance

- Throughput:
  - 10MB/s allocation rate
  - 10MB/s deallocation rate

Temporal Performance

- Throughput:
  - 10MB/s allocation rate
  - 10MB/s deallocation rate
- Latency/Responsiveness:
  - 1ms execution time (malloc/free)
  - 0.1ms preemption time (malloc/free)

Spatial Performance

- Degree of fragmentation:
  - The number of contiguous pieces of memory of a given size that can still be allocated

Spatial Performance

- Degree of fragmentation:
  - The number of contiguous pieces of memory of a given size that can still be allocated
- Administrative space:
  - meta data structures (used, free lists)

There is a trade-off between temporal and spatial performance

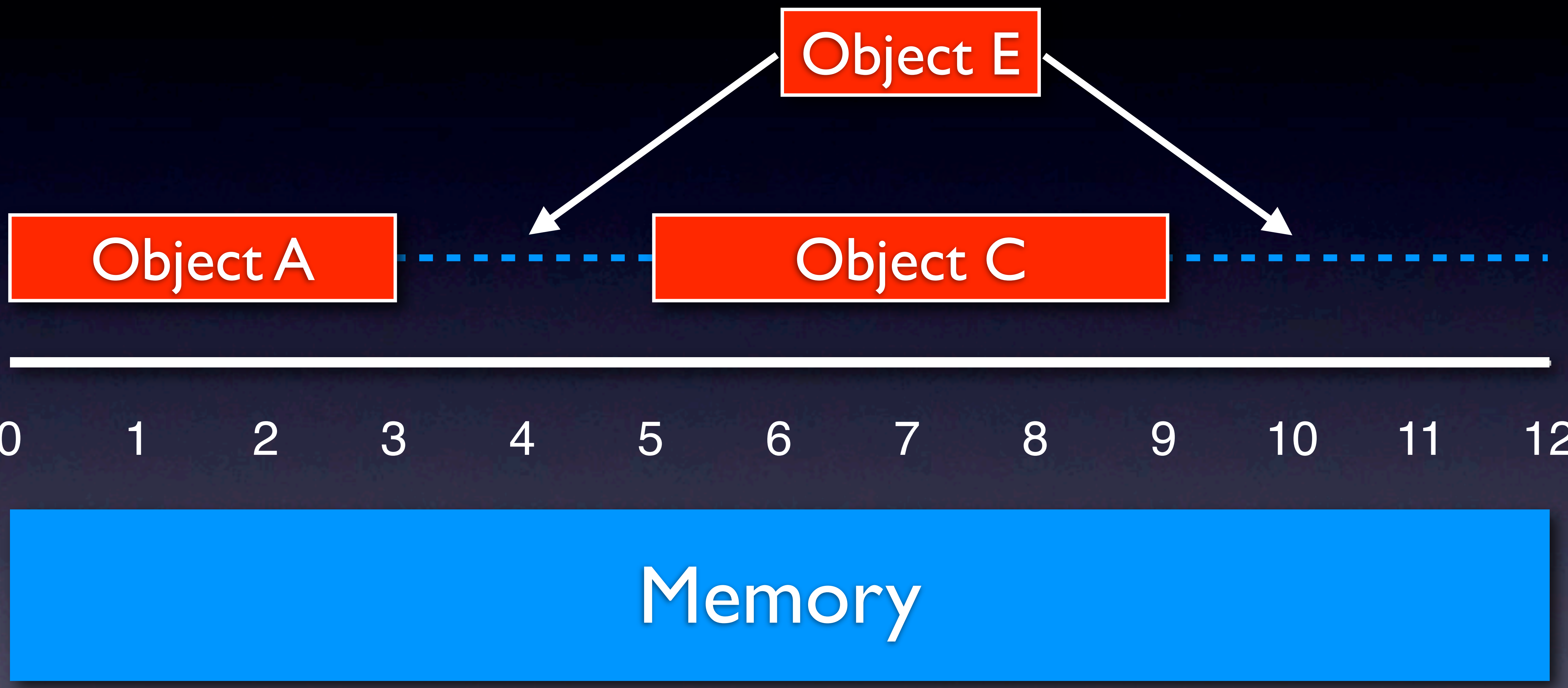
Temporal Predictability

- Unpredictable complexity (in terms of input):
  - allocation/deallocation may take time proportional to the total size of memory

Temporal Predictability

- Unpredictable complexity (in terms of input):
  - allocation/deallocation may take time proportional to the <u>total size of memory</u>
- Predictable complexity (in terms of input):
  - allocation/deallocation takes time at most proportional to the <u>size of involved object</u>
  - access takes time at most proportional to the <u>size of involved object</u>

Allocation Complexity



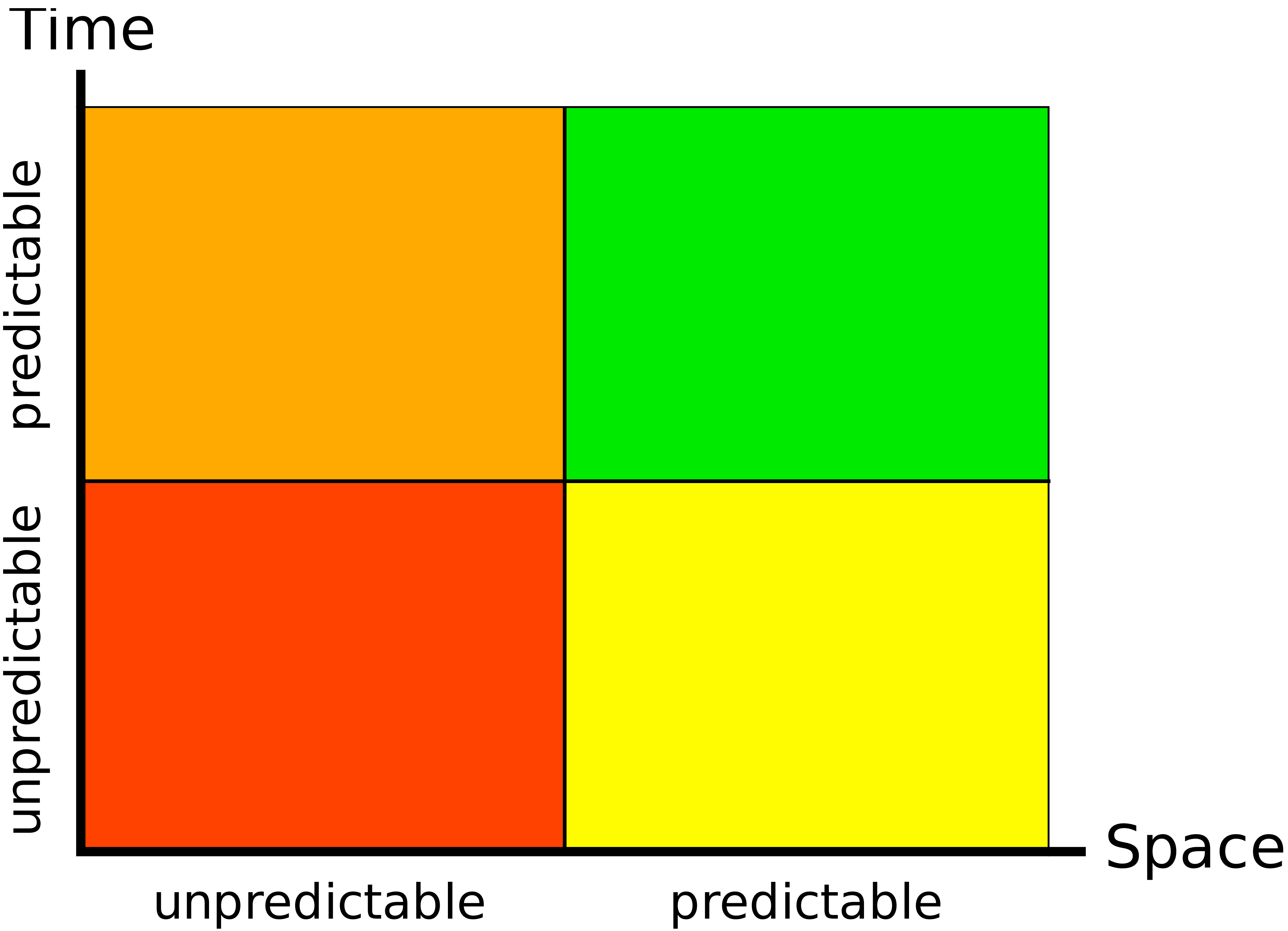
It may be <u>difficult</u> to improve
average performance
but it may still be <u>possible</u> to improve
predictability
without loosing too much
performance

Spatial Predictability

- Unpredictable fragmentation:
  - the degree of fragmentation may depend on the full allocation and deallocation history, i.e., the order of invocations

Spatial Predictability

- Unpredictable fragmentation:
  - the degree of fragmentation may depend on the full allocation and deallocation history, i.e., the order of invocations
- Predictable fragmentation:
  - the degree of fragmentation only depends on the number of allocations and deallocations, independently of the order of invocations

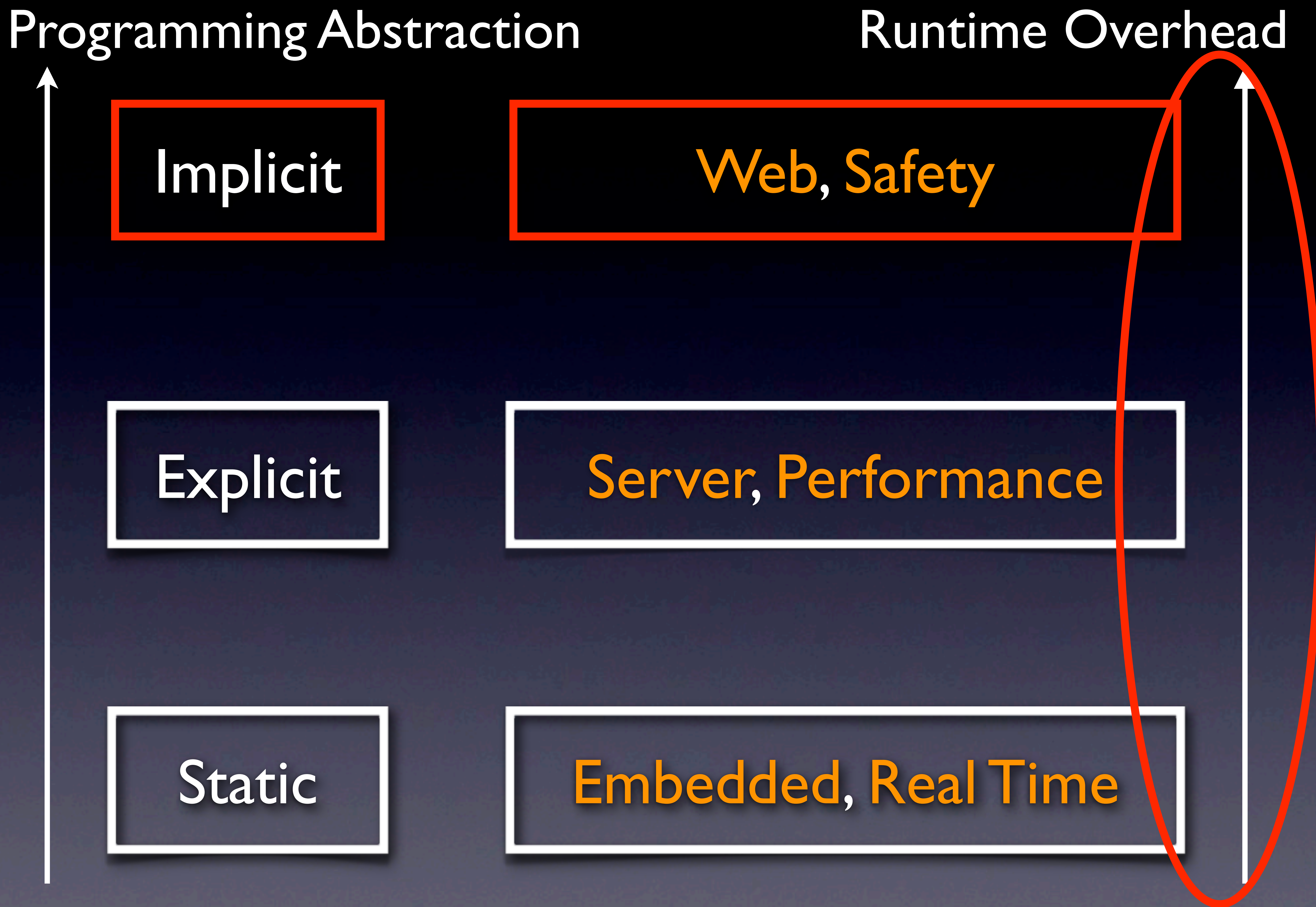

Time
predictable
unpredictable
Space
unpredictable
predictable

Explicit, Dynamic Memory Management with Temporal and Spatial Guarantees

Programming Abstraction

Runtime Overhead

Implicit	Web, Safety
Explicit	Server, Performance
Static	Embedded, Real Time


Programming Abstraction
Runtime Overhead
Implicit
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Programming Abstraction

Runtime Overhead

Programming Abstraction	Runtime Overhead
Implicit	Web, Safety
Explicit	Server, Performance
Static	Embedded, Real Time

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- Silviu Craciunas[#] (Programming Model)
- Andreas Haas (Memory Management)
- Hannes Payer[#] (Memory Management)
- Harald Röck (VM, Scheduling)
- Ana Sokolova* (Theoretical Foundation)

[#]Supported by a 2007 IBM Faculty Award, the EU ArtistDesign Network of Excellence on Embedded Systems Design, and Austrian Science Fund Project P18913-N15.
*Supported by Austrian Science Fund Project V00125.

Tiptoe

- Tiptoe is a microkernel-based virtual machine and process monitor for embedded systems

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- Tiptoe virtualizes the host platform (system VM) and provides infrastructure to run process VMs and processes in real time

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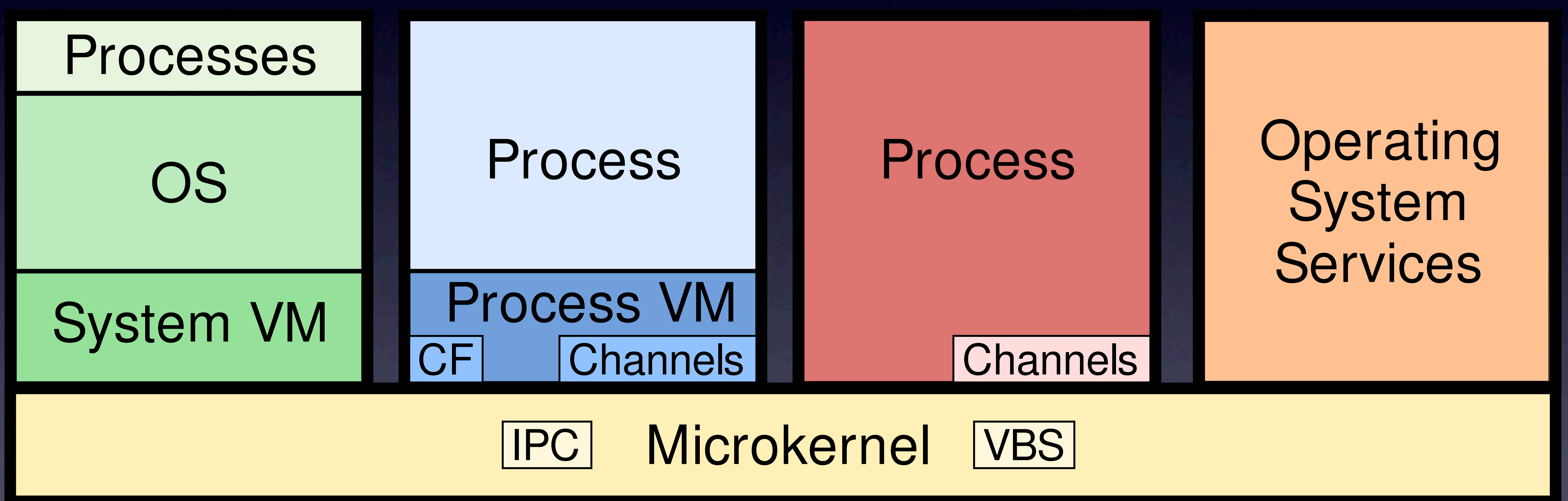
- Tiptoe virtualizes the host platform (system VM) and provides infrastructure to run process VMs and processes in real time
- Tiptoe controls throughput and latency of CPU, memory, and I/O
- I/O is multiplexed through IPC to a system VM running Linux

Tiptoe



Processes
OS
System VM
Process
Process VM
CF
Channels
Process
Channels
Operating System Services
IPC
Microkernel
VBS

The JAviator

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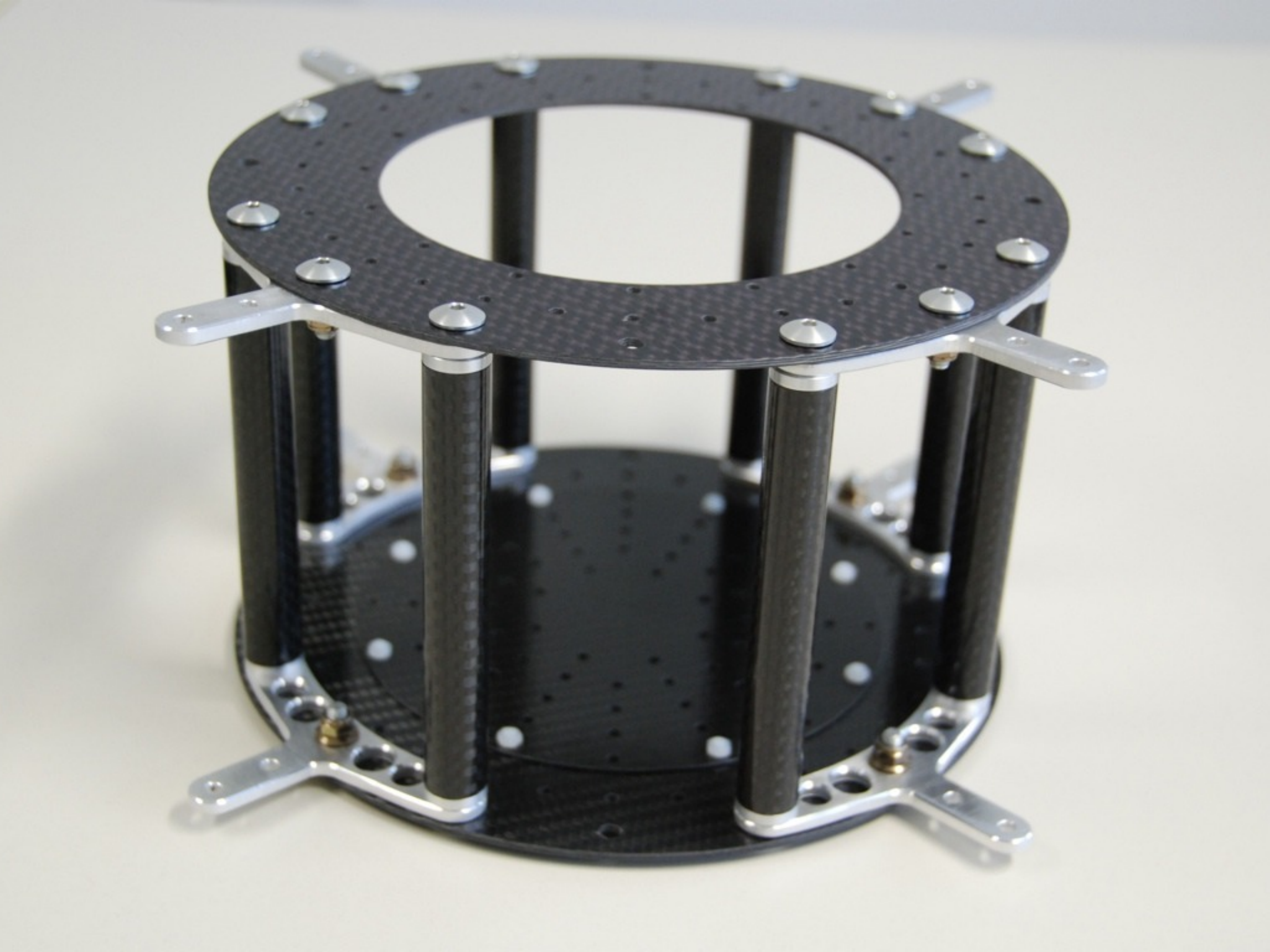
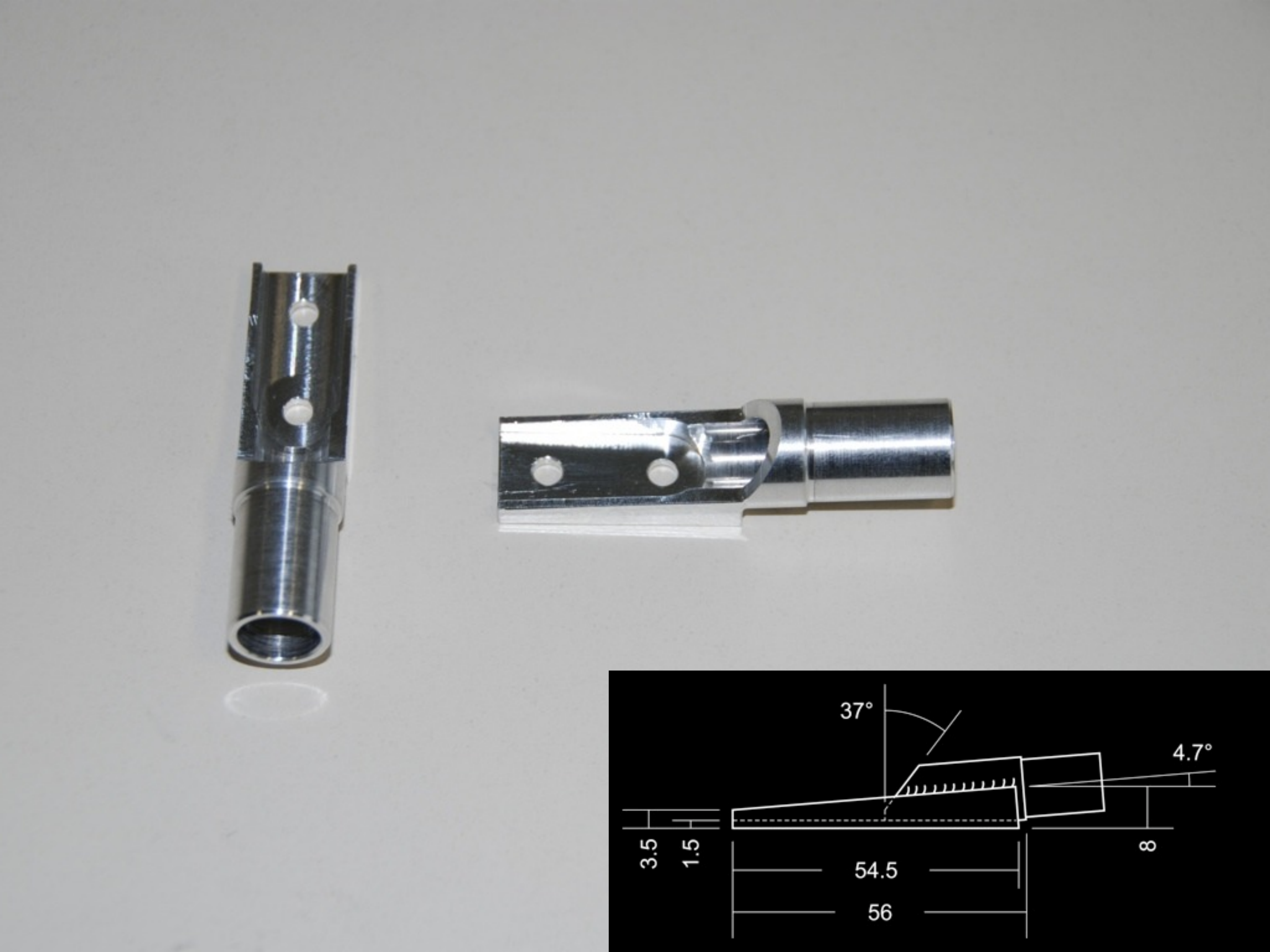
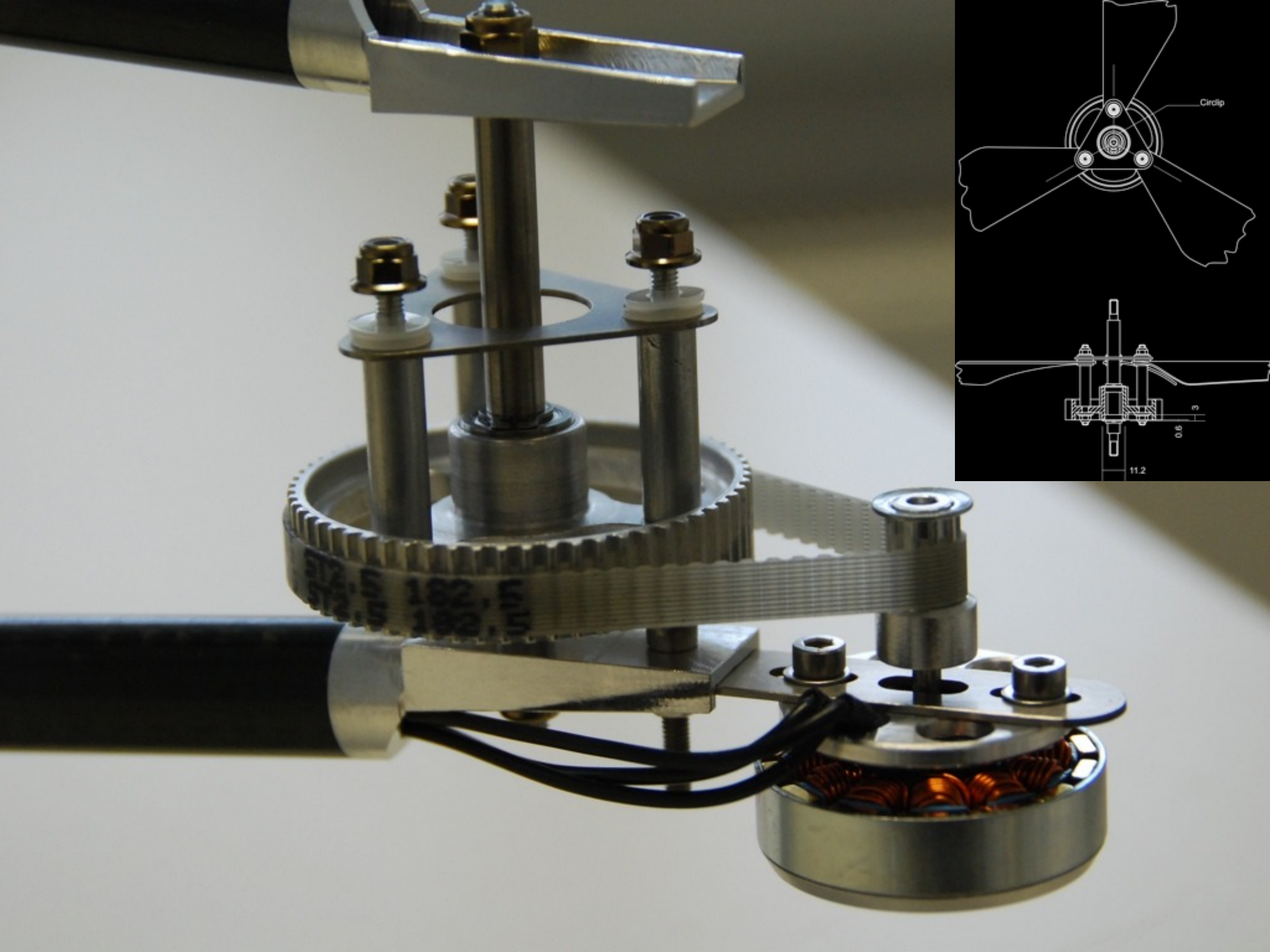
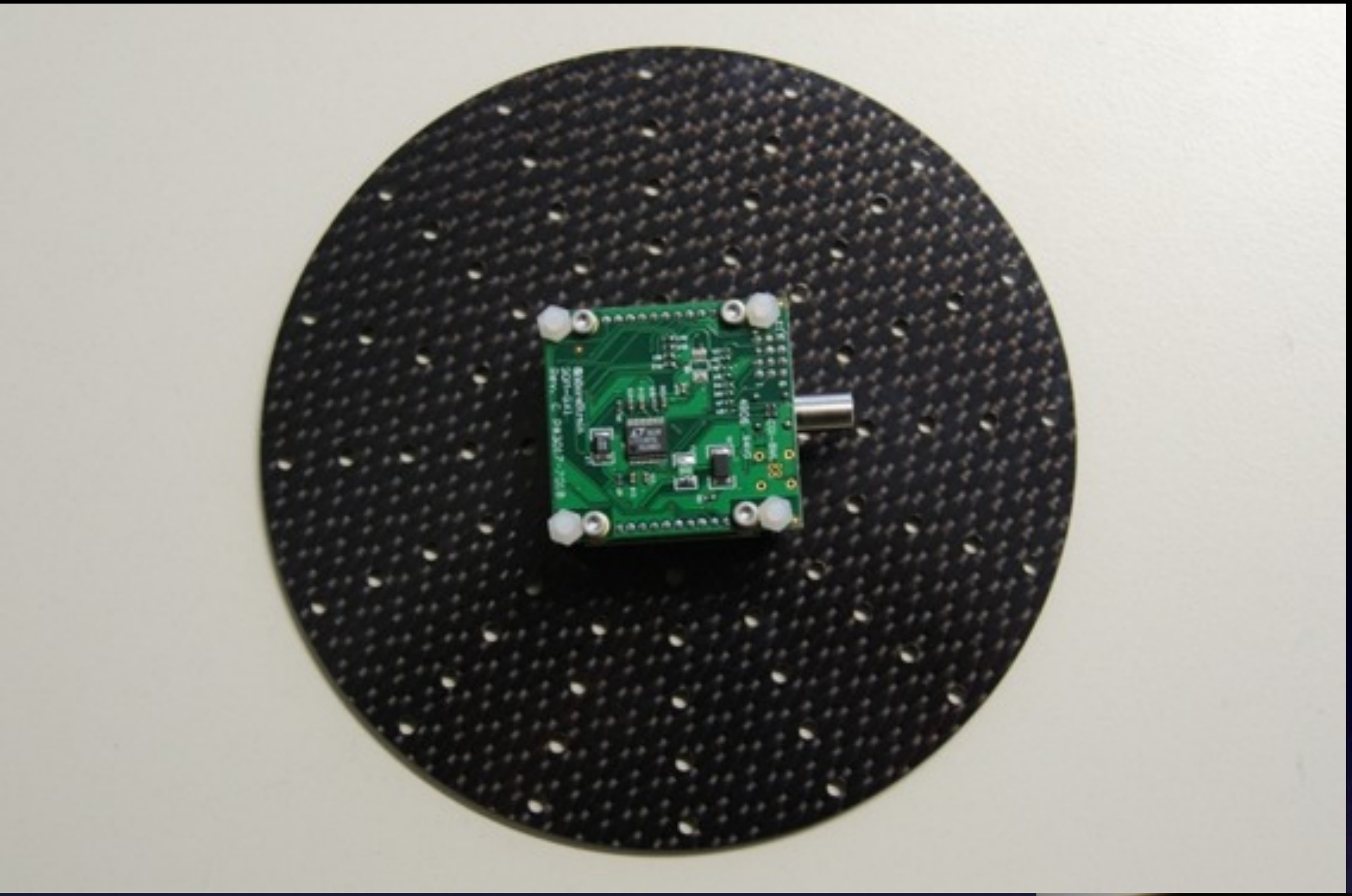
Quad-Rotor Helicopter




37°
4.7°
3.5
1.5
54.5
56
8

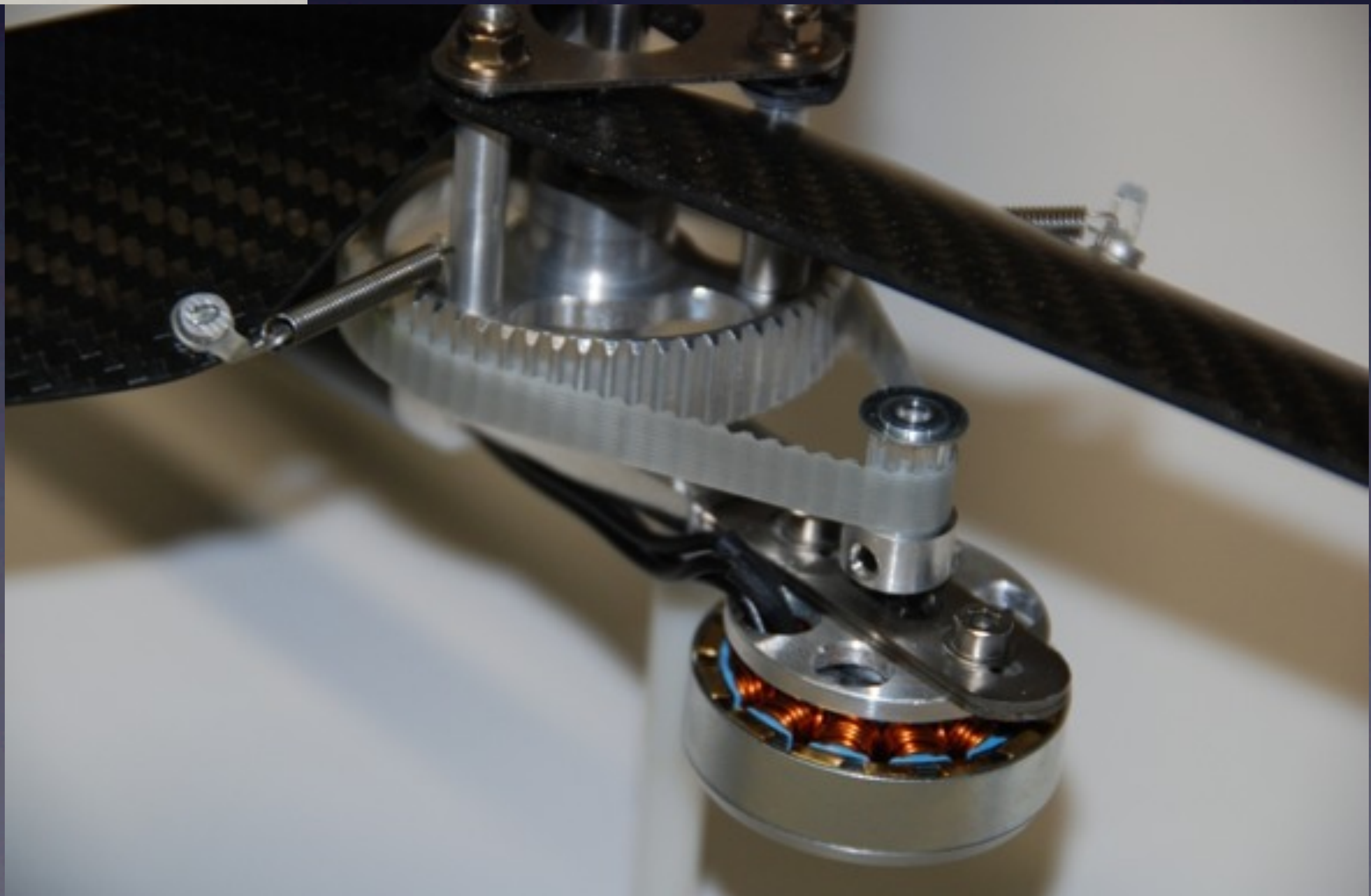

Circlip
3
0.6
11.2
ST2.5 182.5


25/12/2005

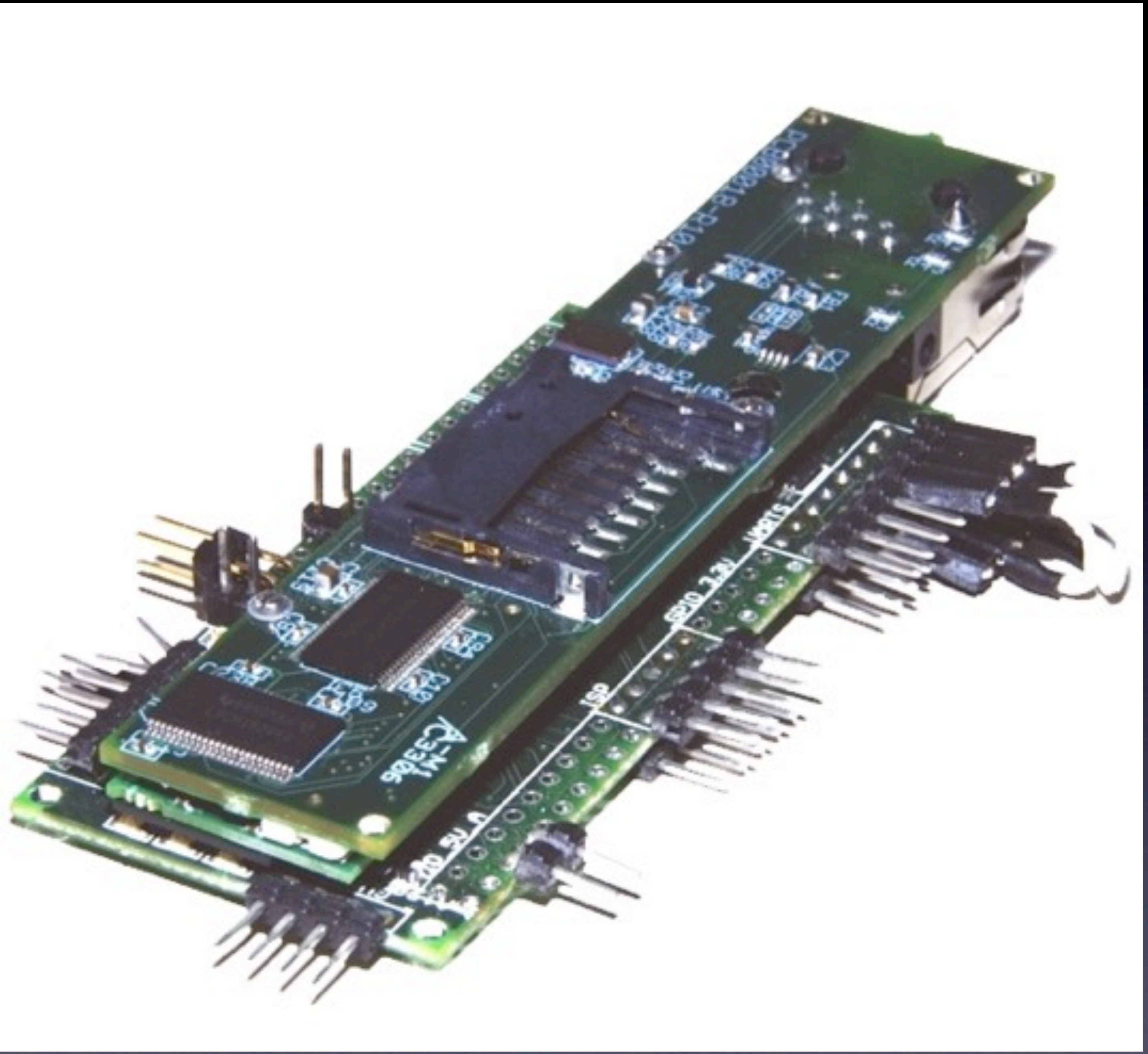


Gyro

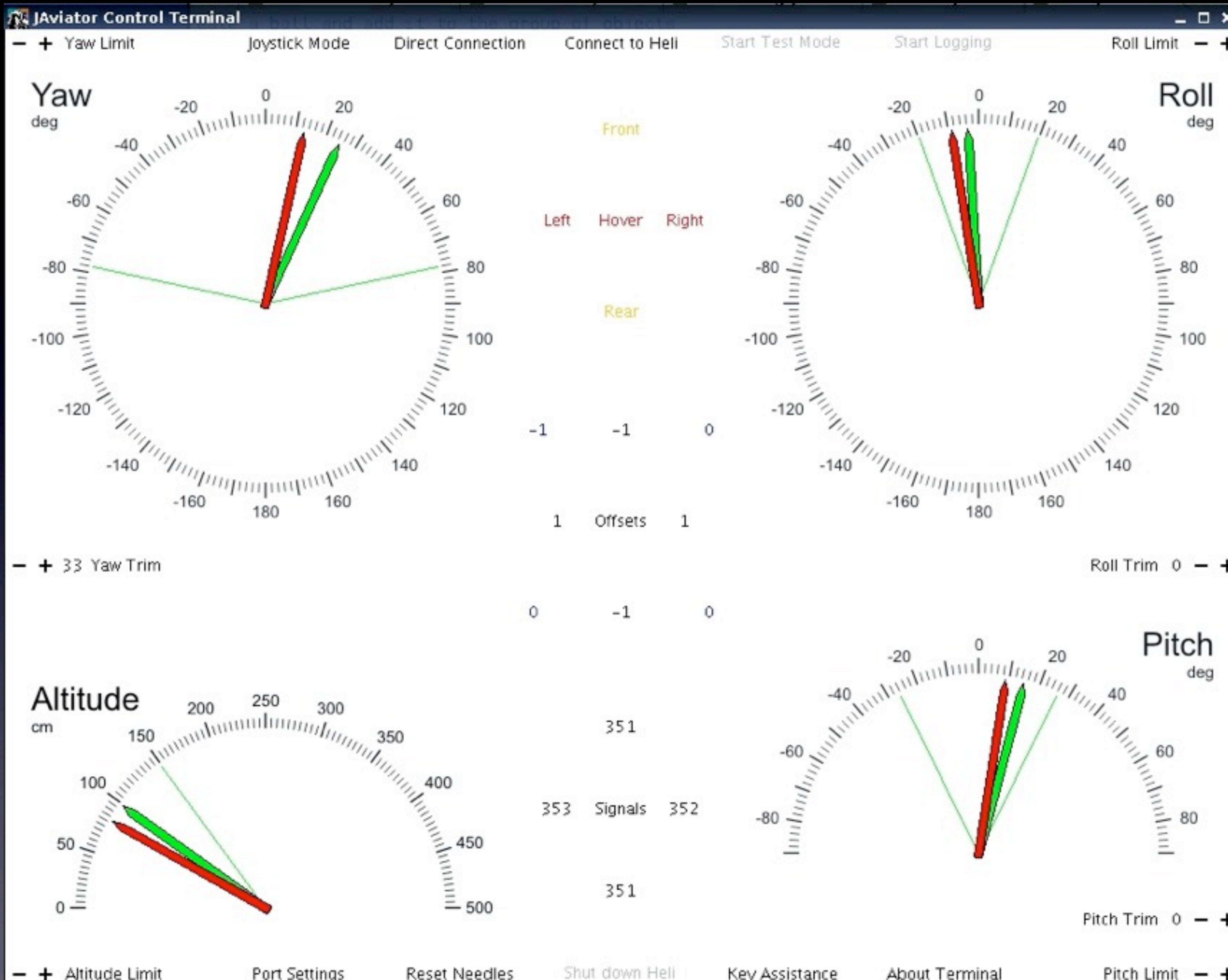
Propulsion

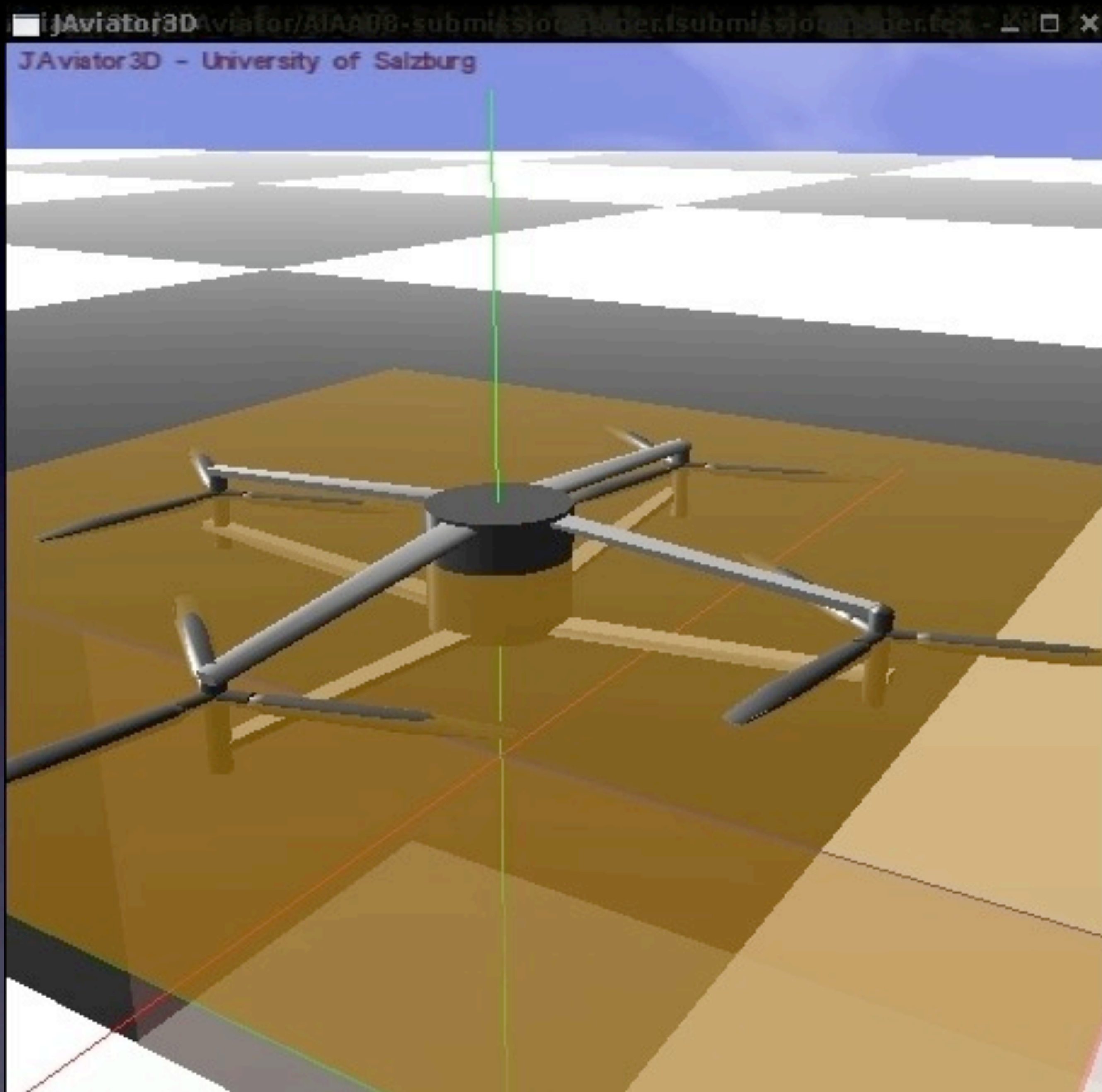


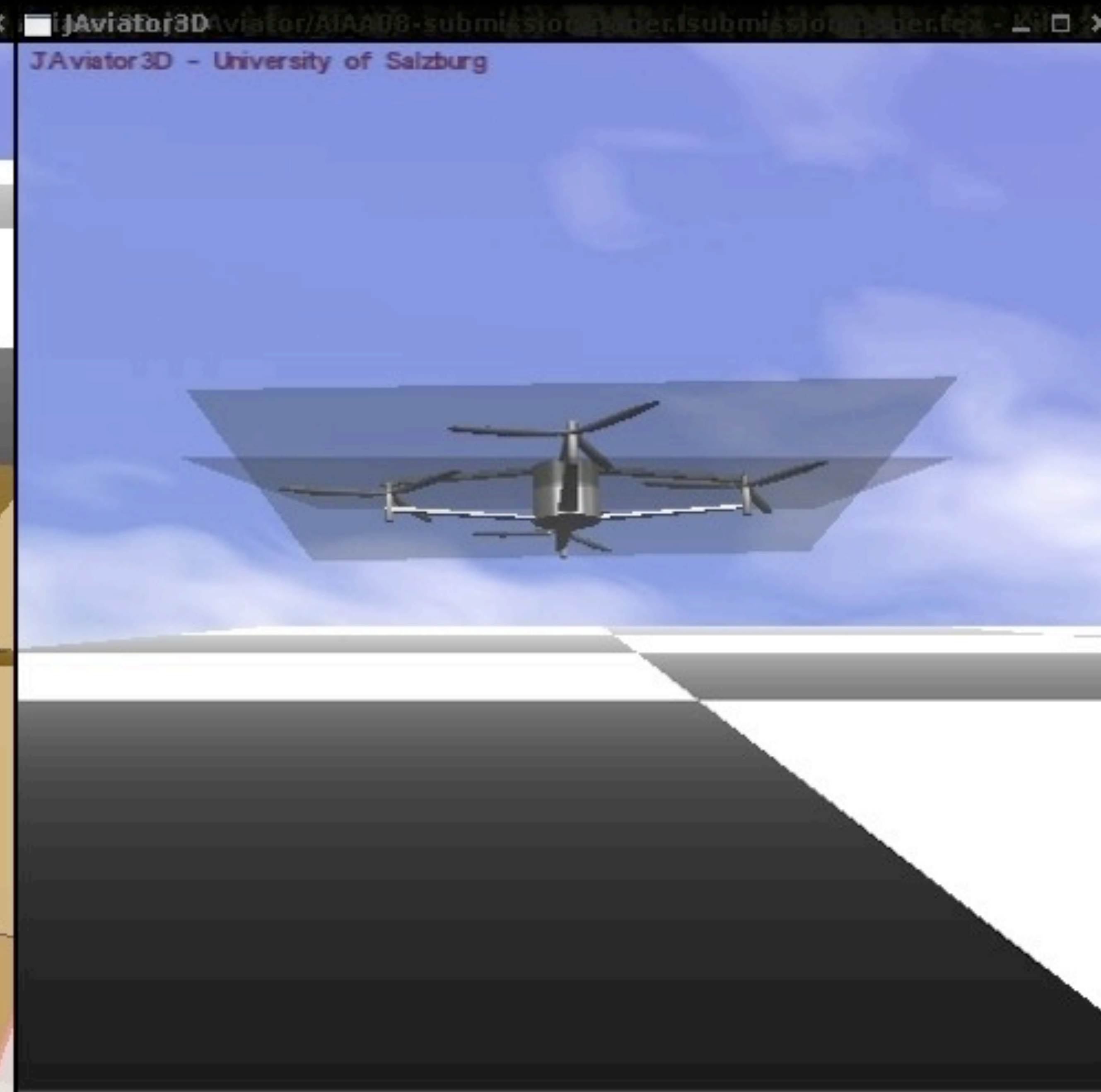
Gumstix



600MHz XScale, 128MB RAM, WLAN, Atmega uController

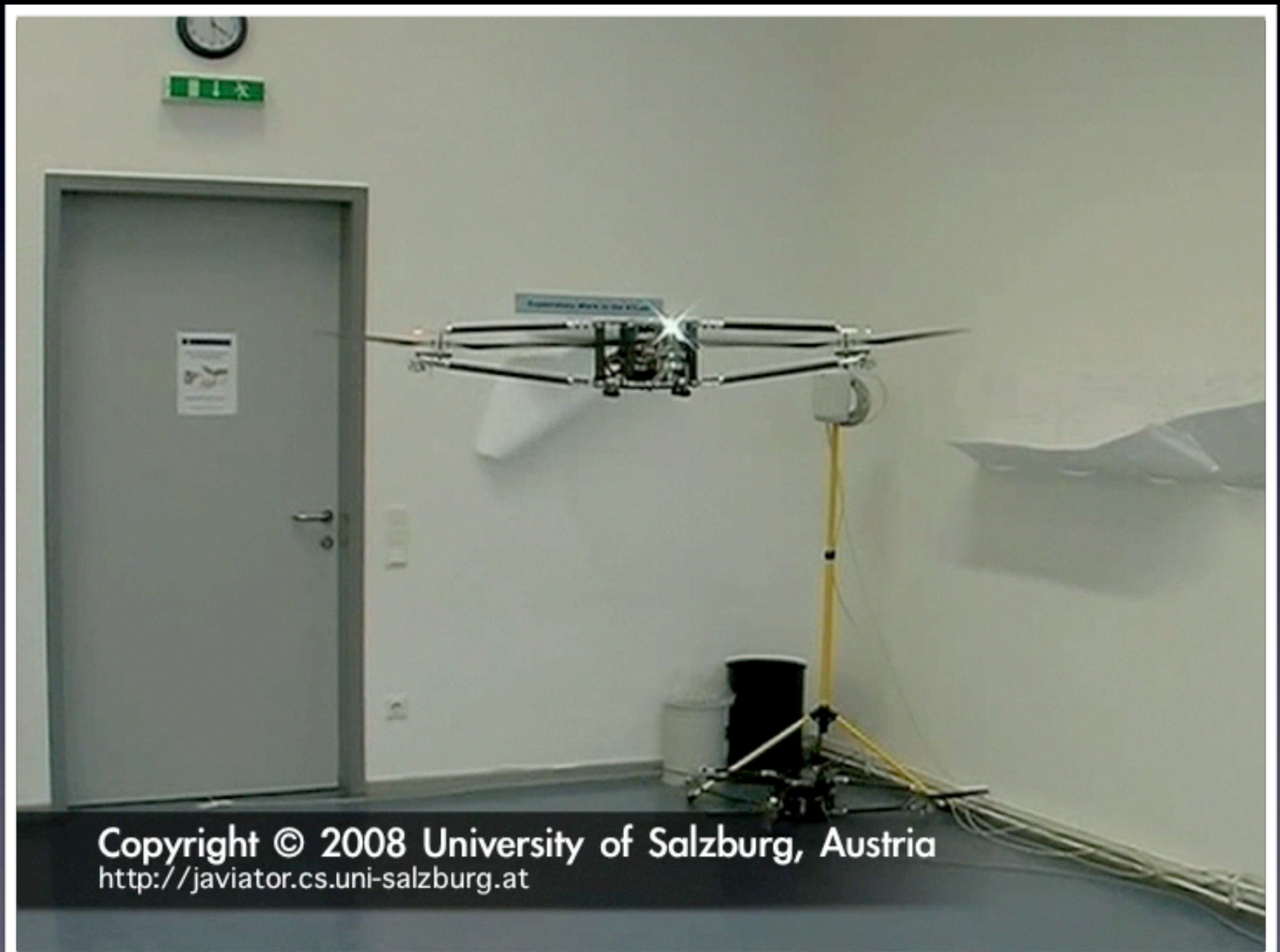

JAviator Control Terminal
Yaw Limit
Joystick Mode
Direct Connection
Connect to Heli
Start Test Mode
Start Logging
Roll Limit
Yaw
deg
Roll
deg
Front
Left
Hover
Right
Rear
-1
-1
0
1
Offsets
1
33 Yaw Trim
Roll Trim 0
0
-1
0
Altitude
cm
Pitch
deg
351
353
Signals
352
351
Pitch Trim 0
Altitude Limit
Port Settings
Reset Needles
Shut down Heli
Key Assistance
About Terminal
Pitch Limit


JAviator3D
JAviator3D - University of Salzburg


JAviator3D
JAviator3D - University of Salzburg

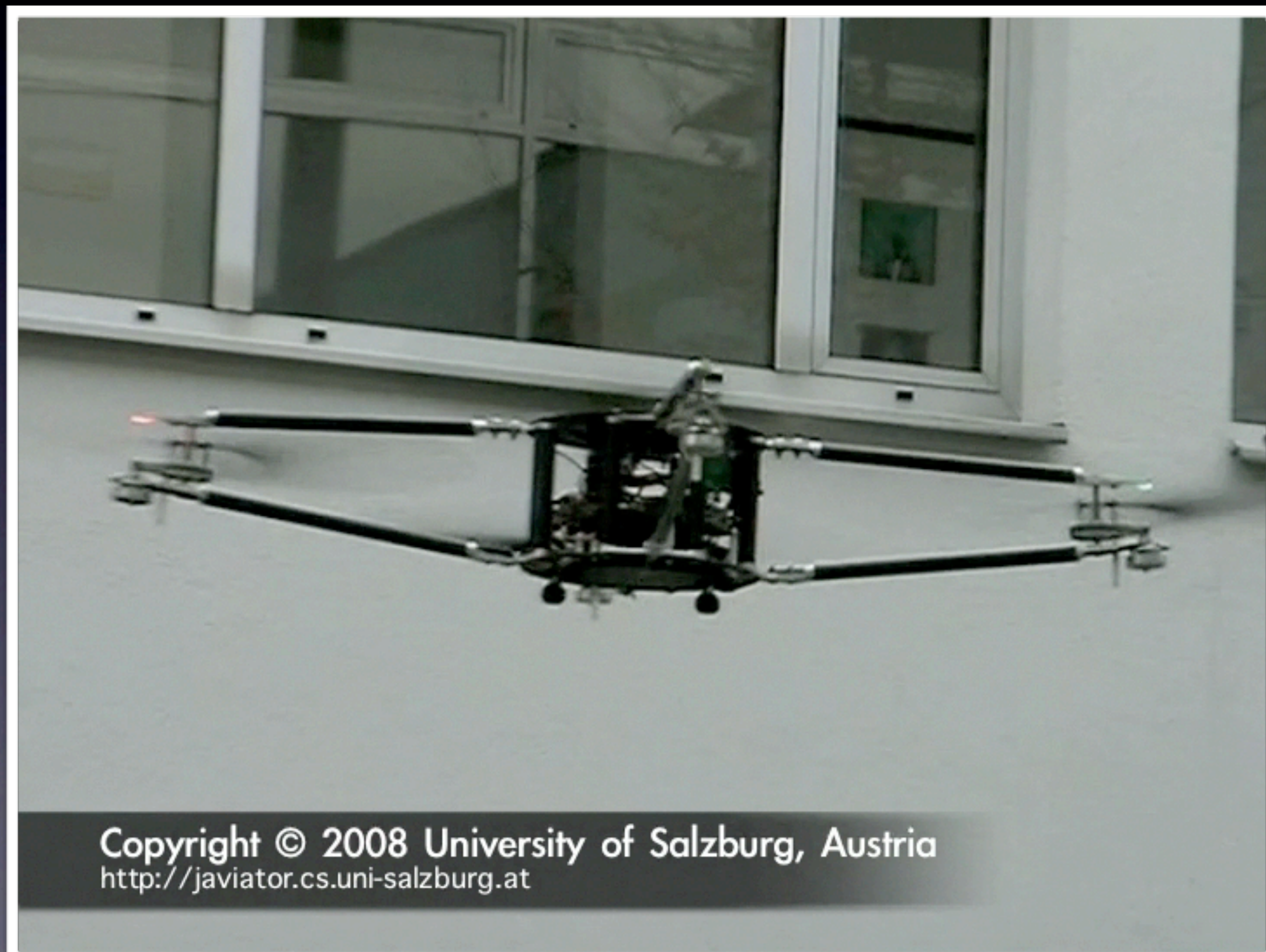
Indoor Flight
STARMAC Controller

Indoor Flight STARMAC Controller



Outdoor Flight STARMAC Controller

Outdoor Flight STARMAC Controller



Outdoor Flight
Salzburg Controller

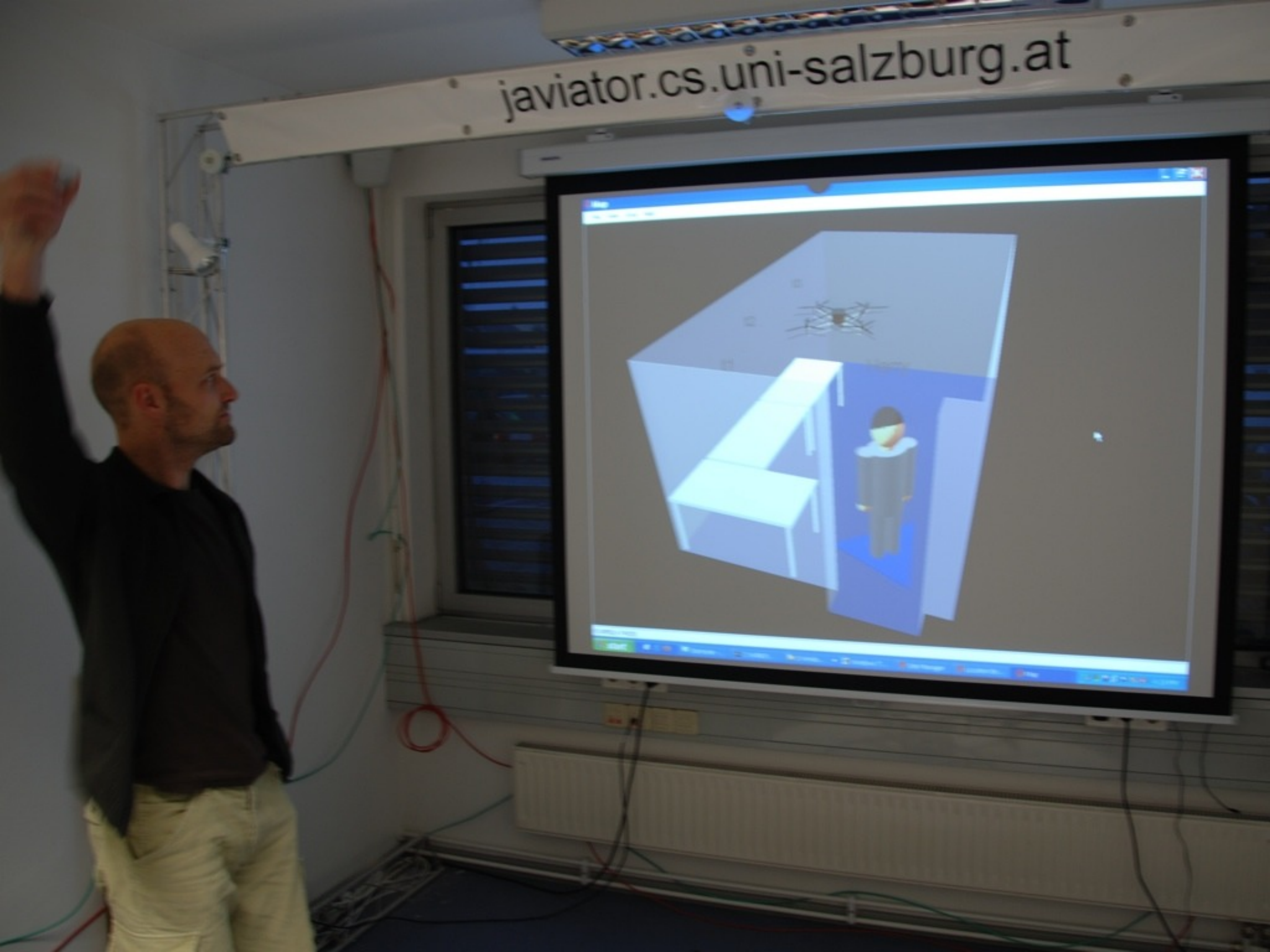
Outdoor Flight
Salzburg Controller

What's next?

- Autonomous single-vehicle flights
  - position controller
  - waypoint controller

What's next?

- Autonomous single-vehicle flights
  - position controller
  - waypoint controller
- Autonomous multi-vehicle flights
  - mission controller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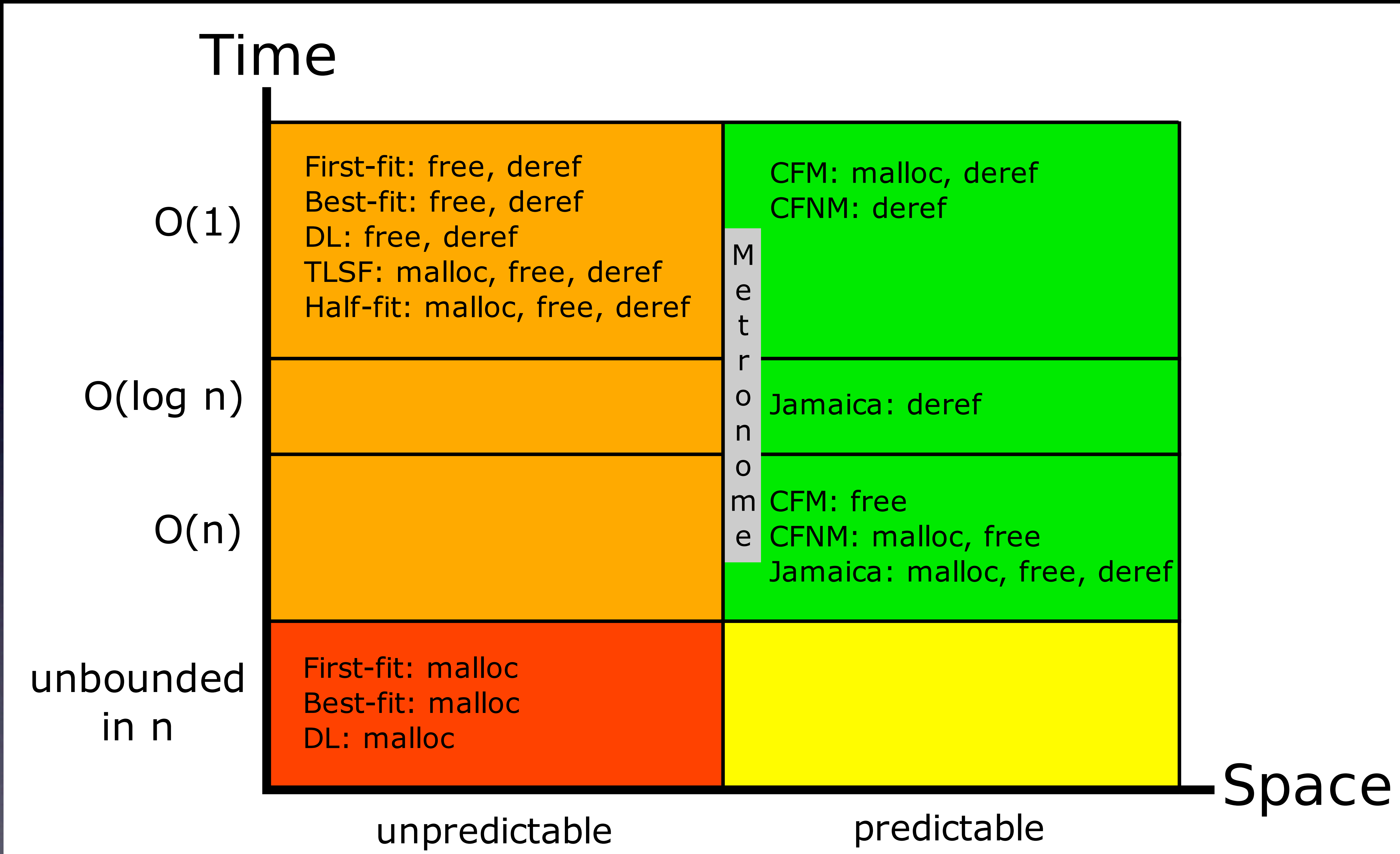
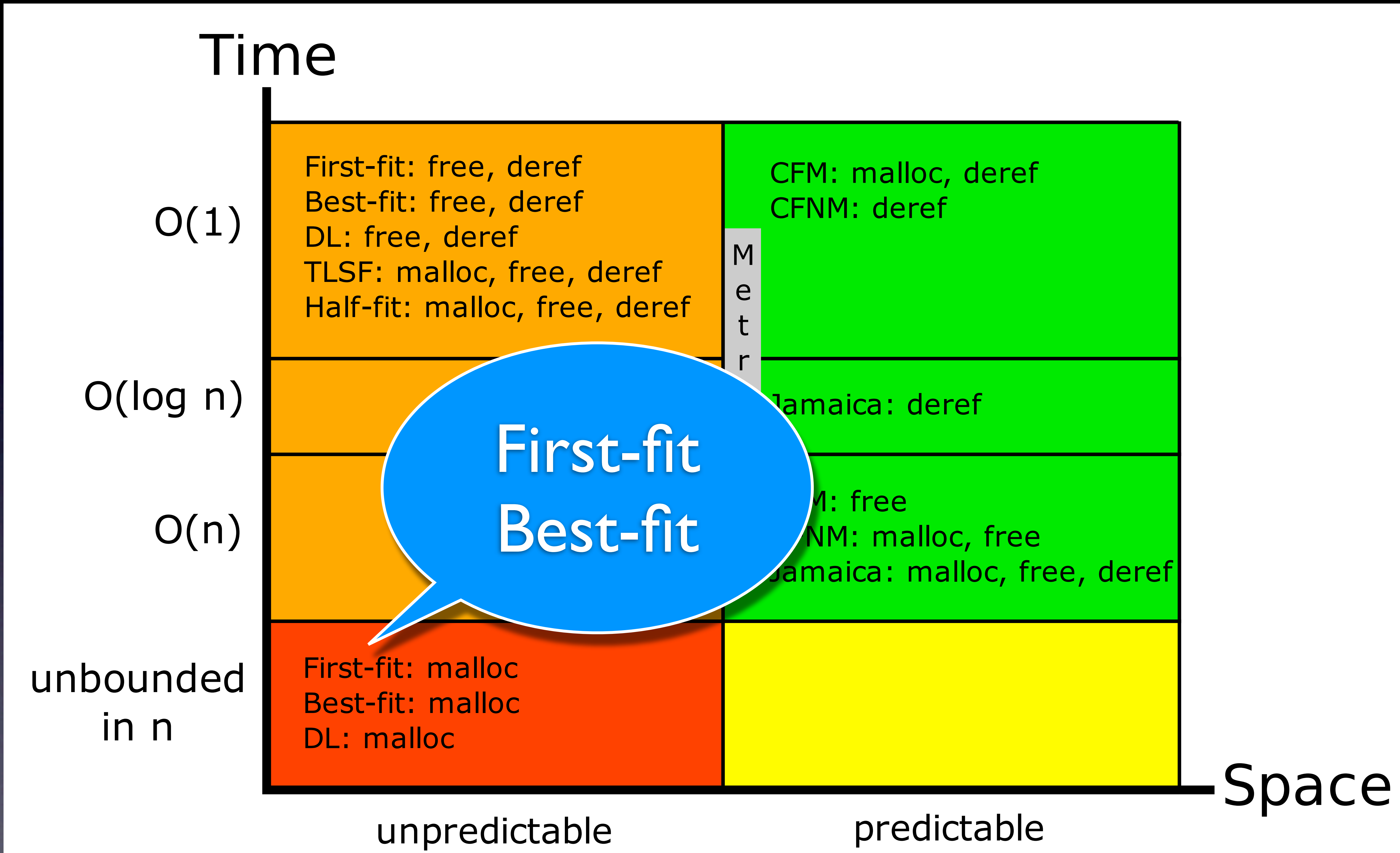
Salzburg Soft Walls Controller on JJ

Salzburg Soft Walls Controller on JJ

Memory Management Systems Overview


Time
O(1)
O(log n)
O(n)
unbounded in n
First-fit: free, deref
Best-fit: free, deref
DL: free, deref
TLSF: malloc, free, deref
Half-fit: malloc, free, deref
CFM: malloc, deref
CFNM: deref
Metronome
Jamaica: deref
CFM: free
CFNM: malloc, free
Jamaica: malloc, free, deref
First-fit: malloc
Best-fit: malloc
DL: malloc
Space
unpredictable
predictable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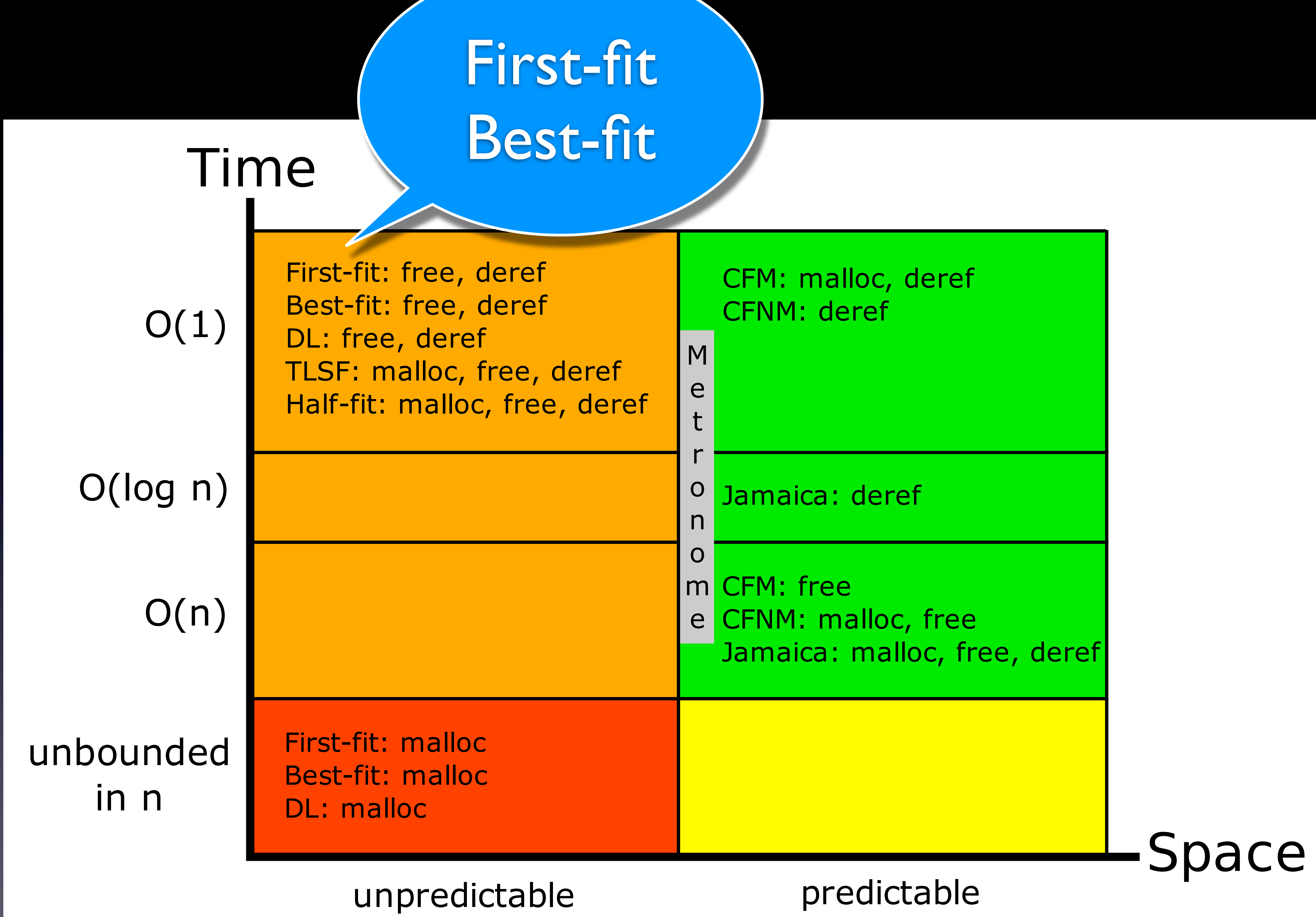
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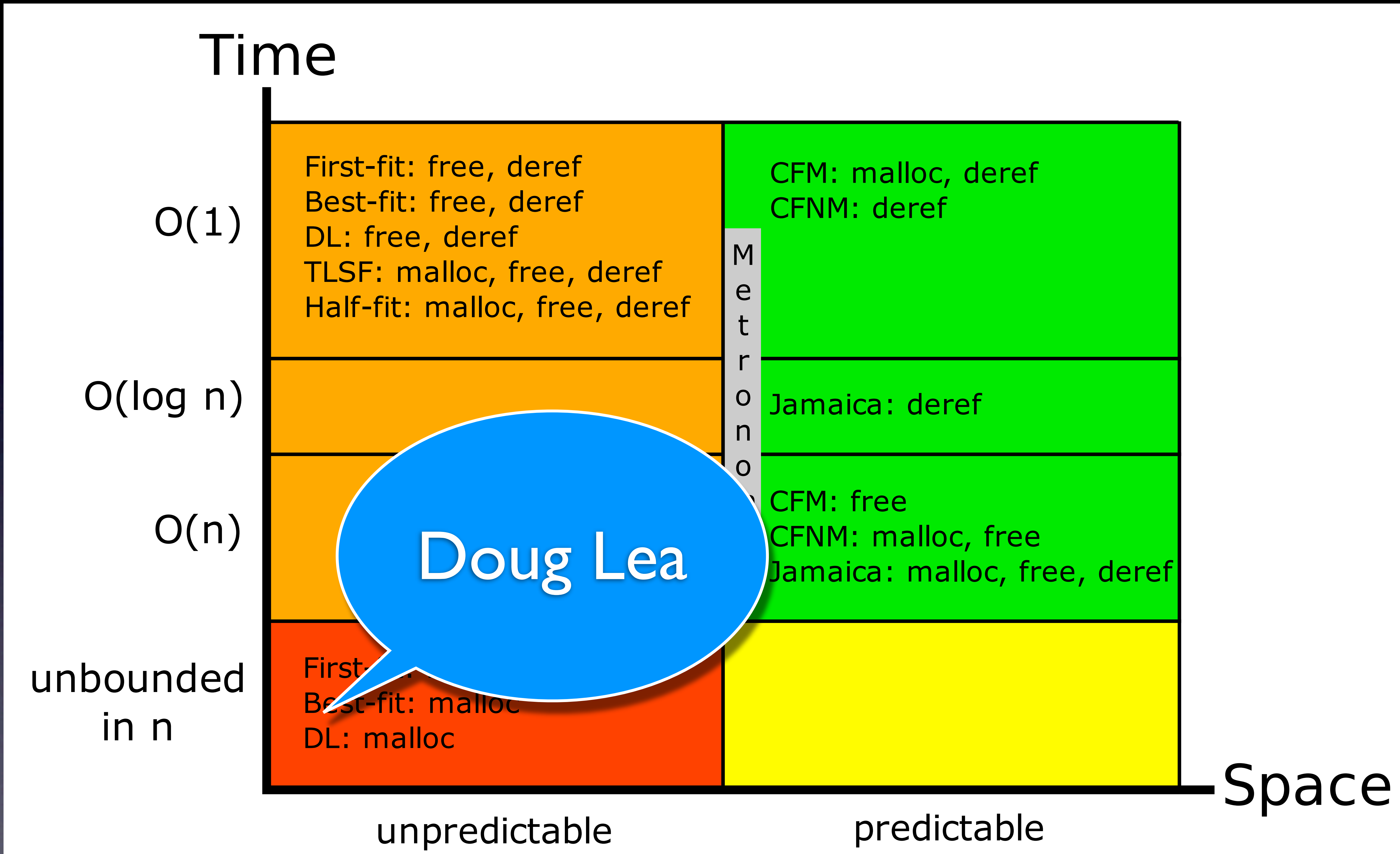
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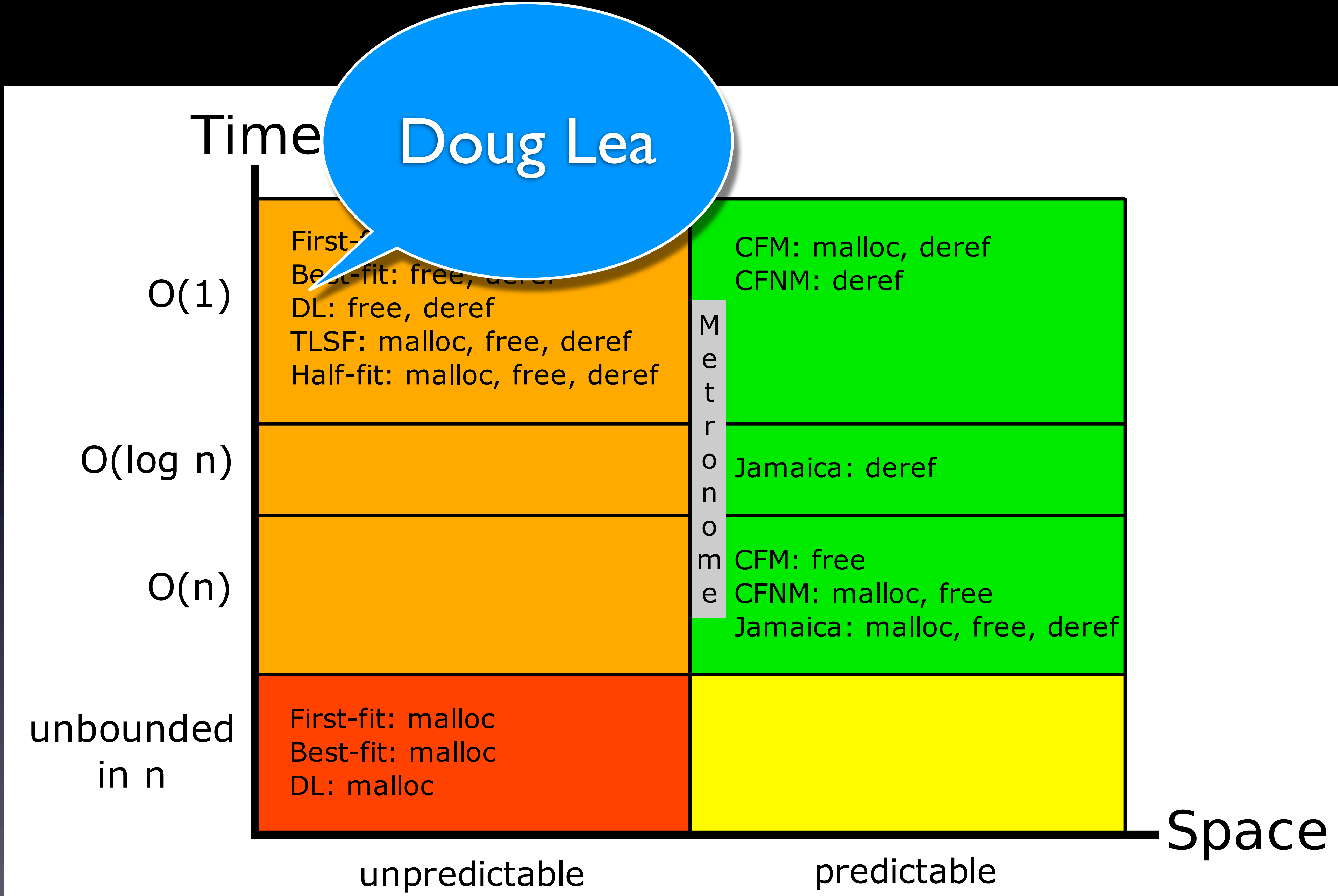
Jamaica: malloc, free, deref
DL: malloc
Doug Lea
unpredictable
predictable
Space


Doug Lea
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in n
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DL: free, deref
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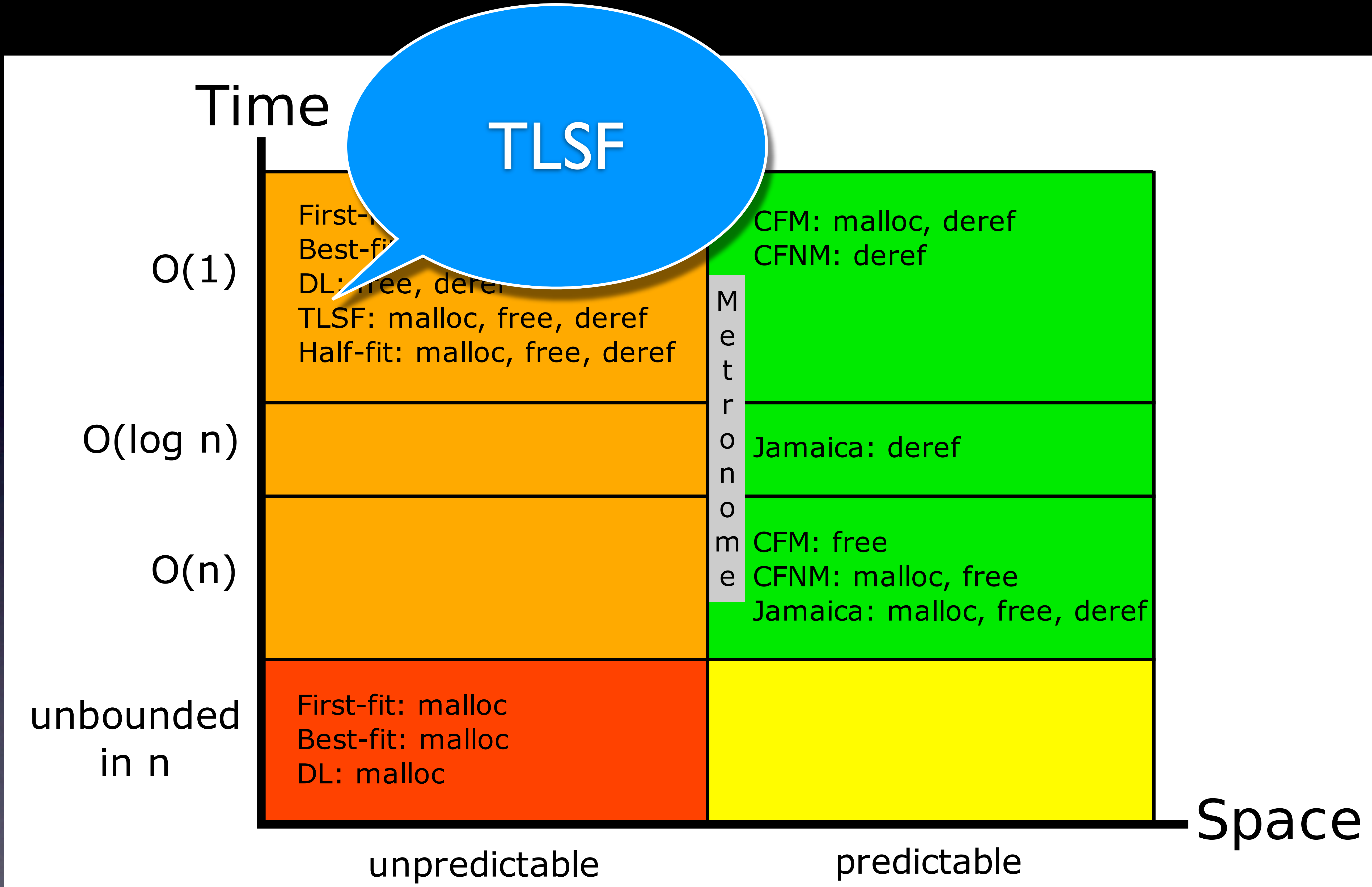

TLSF
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DL:
TLSF: malloc, free, deref
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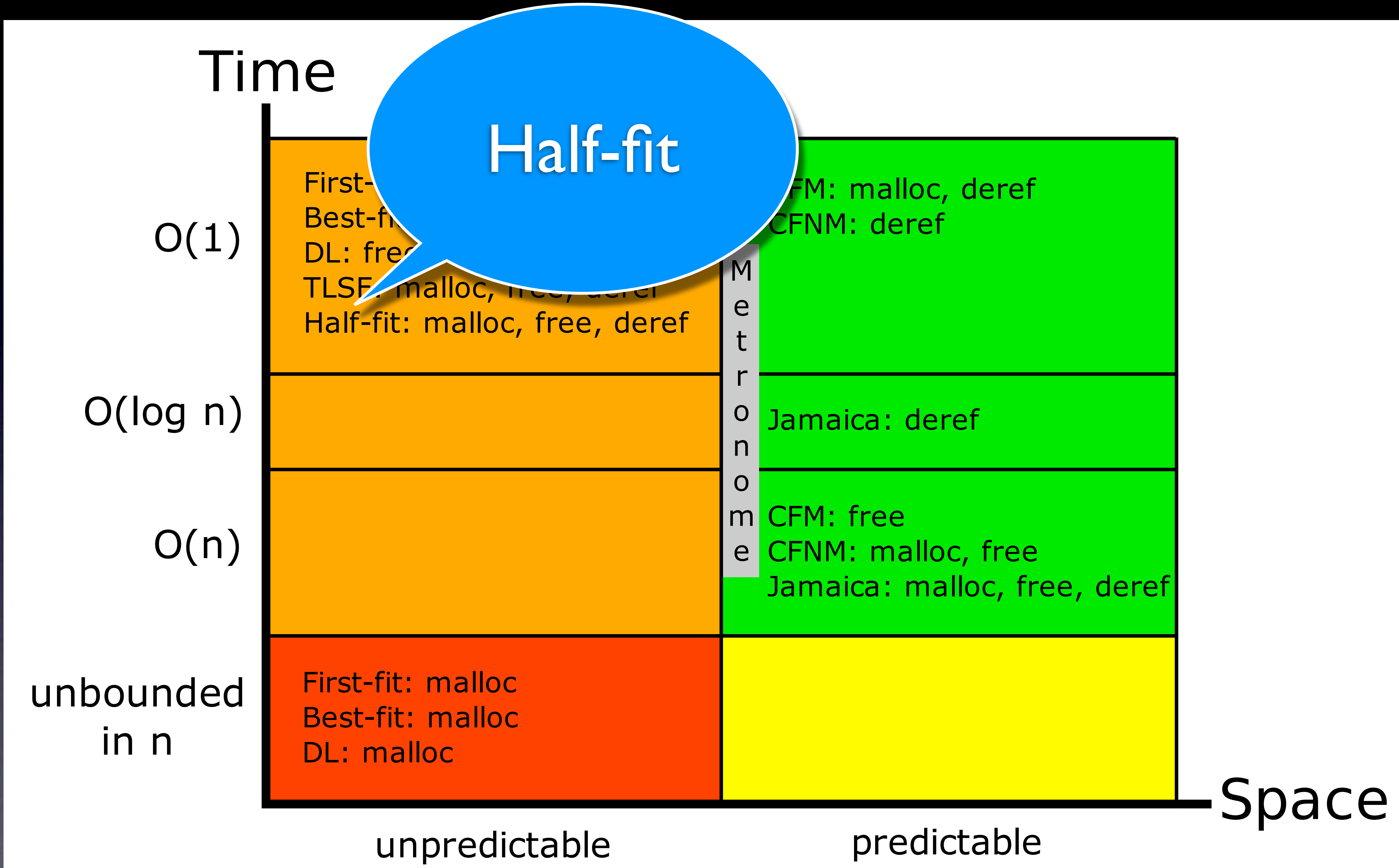

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DL: fre
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First-fit: malloc
Best-fit: malloc
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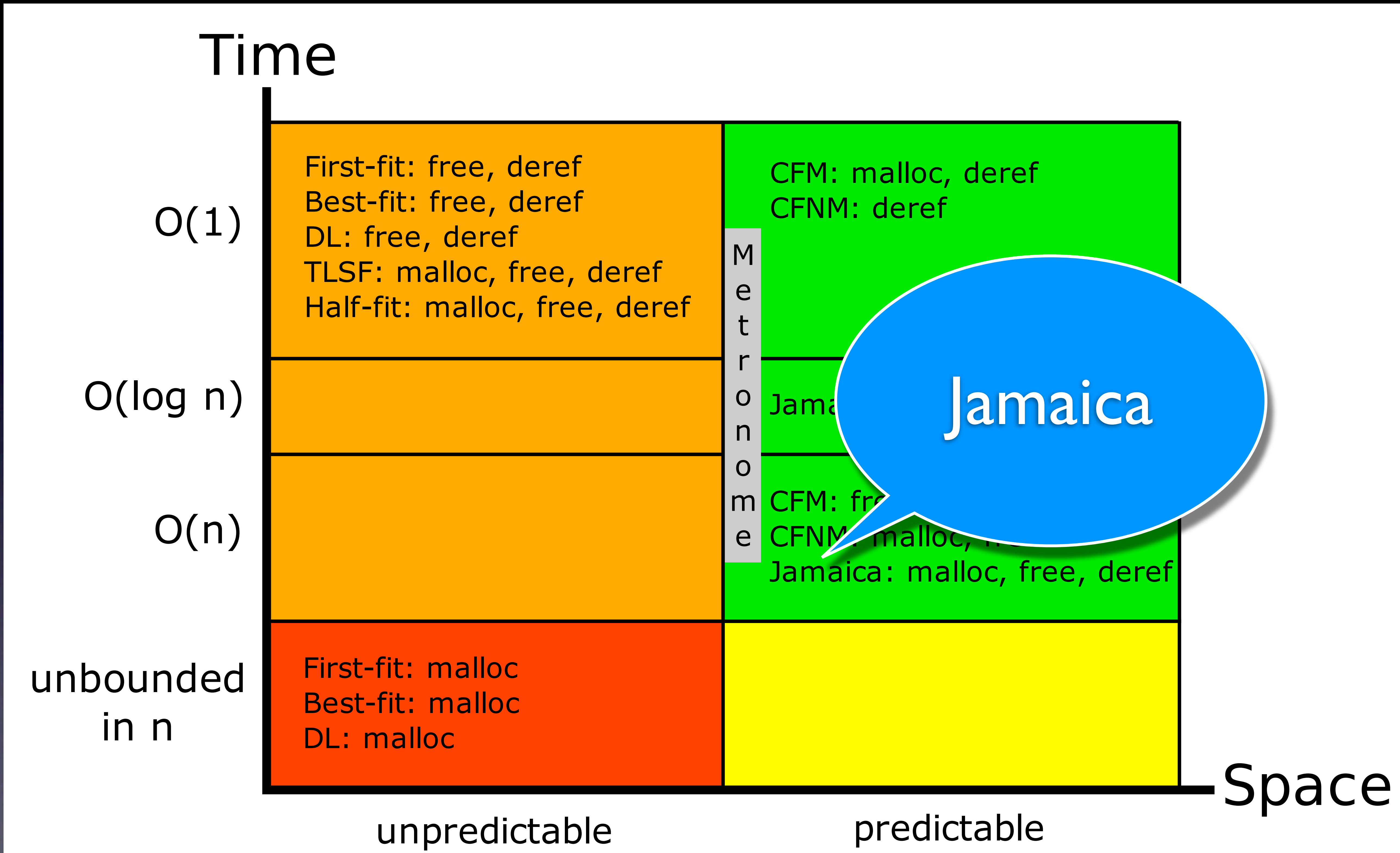
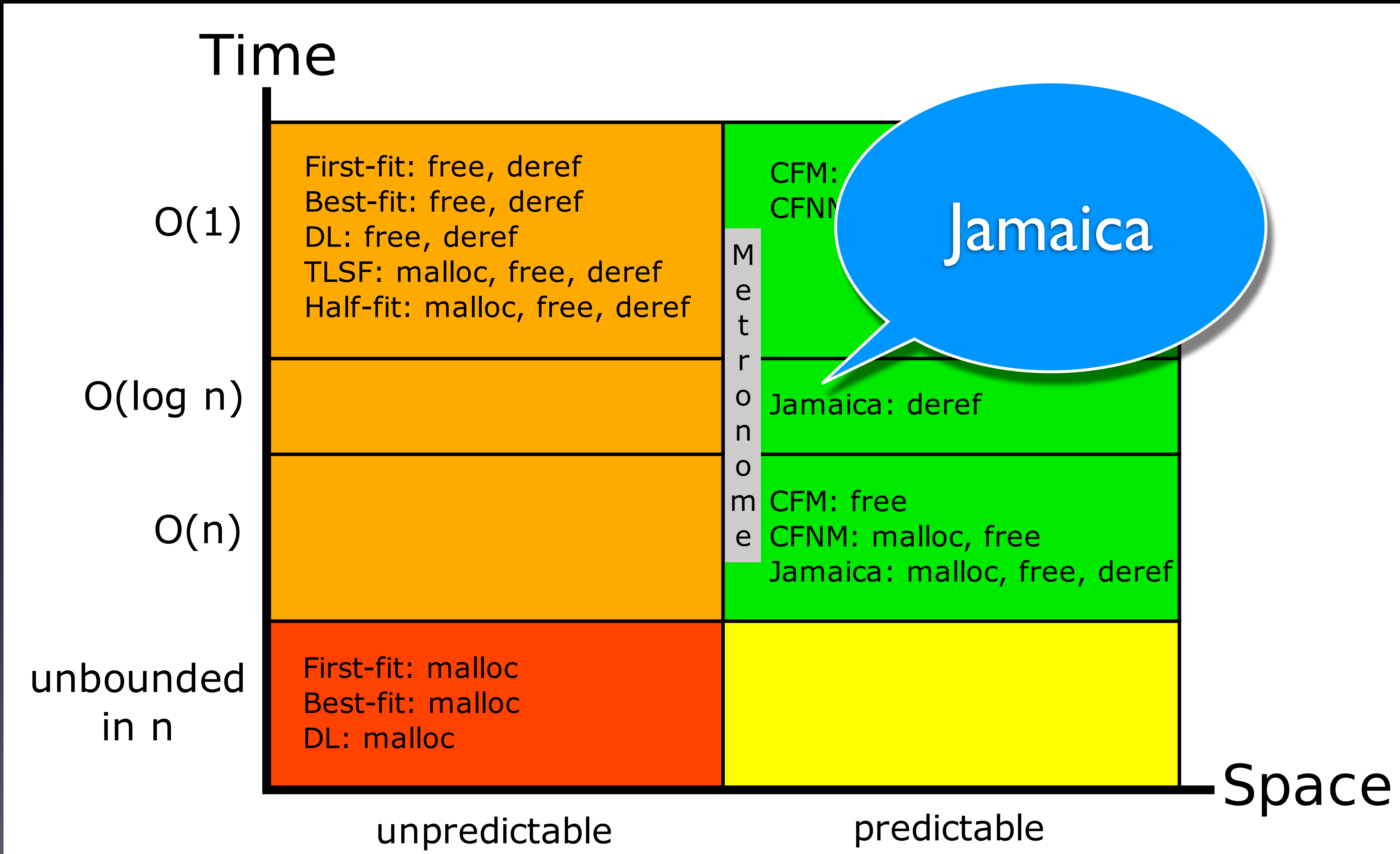
unpredictable
predictable


Time
O(1)
O(log n)
O(n)
unbounded in n
First-fit: free, deref
Best-fit: free, deref
DL: free, deref
TLSF: malloc, free, deref
Half-fit: malloc, free, deref
CFM: malloc, deref
CFNM: deref
Metronome
Jamaica
Jamaica: malloc, free, deref
First-fit: malloc
Best-fit: malloc
DL: malloc
unpredictable
predictable
Space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CFN
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CFM: free
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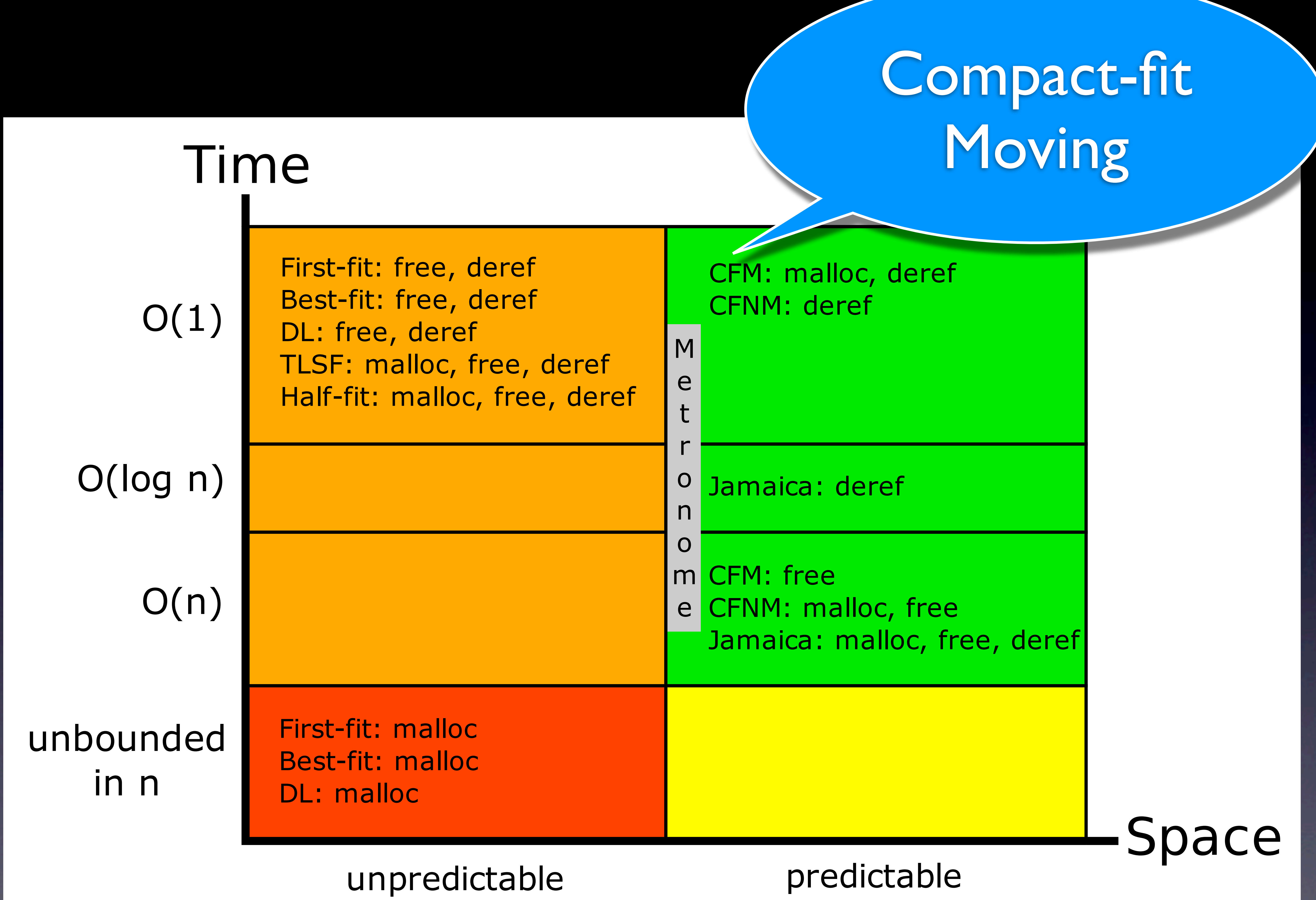
Space
unpredictable
predictable


Compact-fit
Moving
Time
O(1)
O(log n)
O(n)
unbounded in n
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DL: free, deref
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First-fit: malloc
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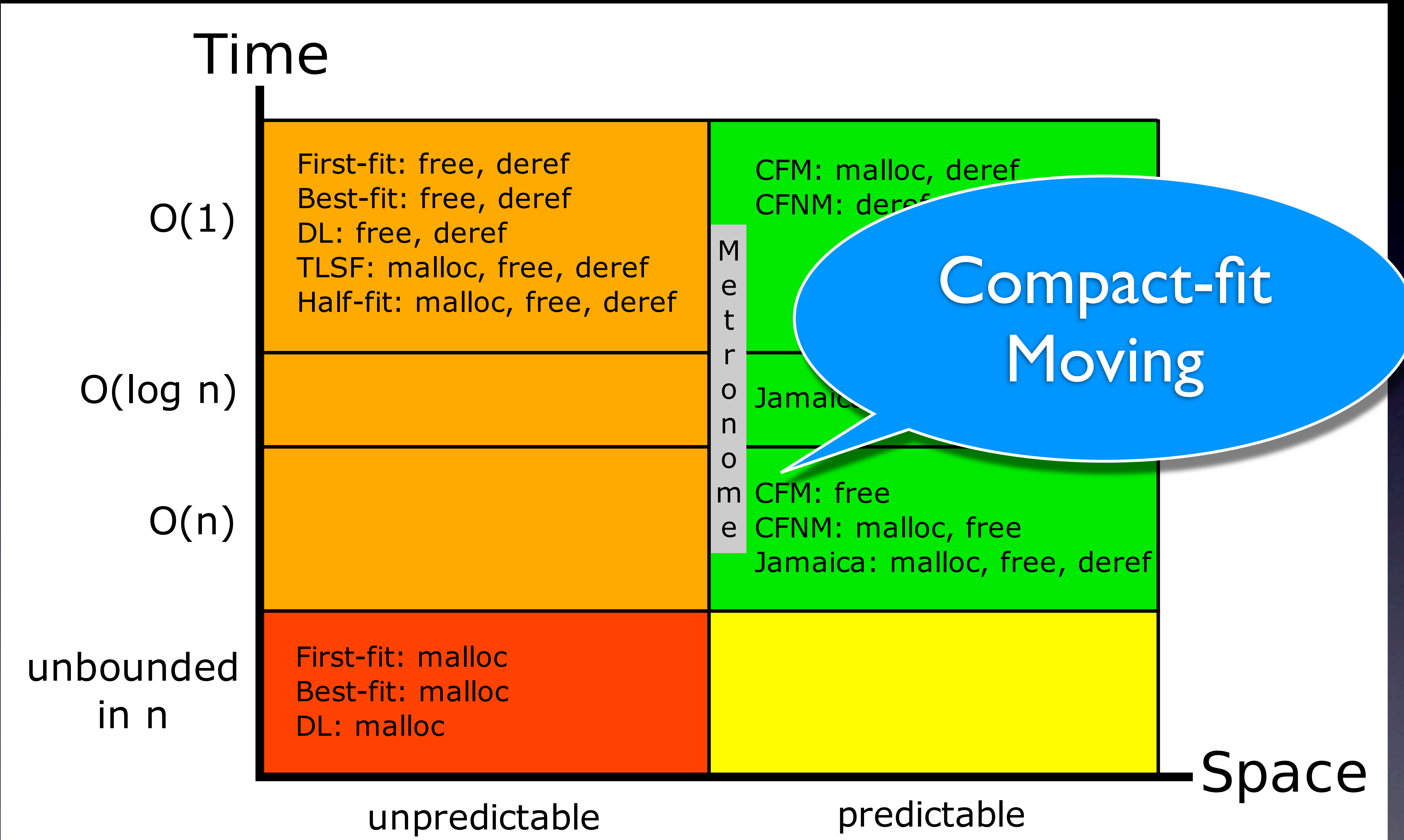
Space
unpredictable
predictable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First-fit: malloc
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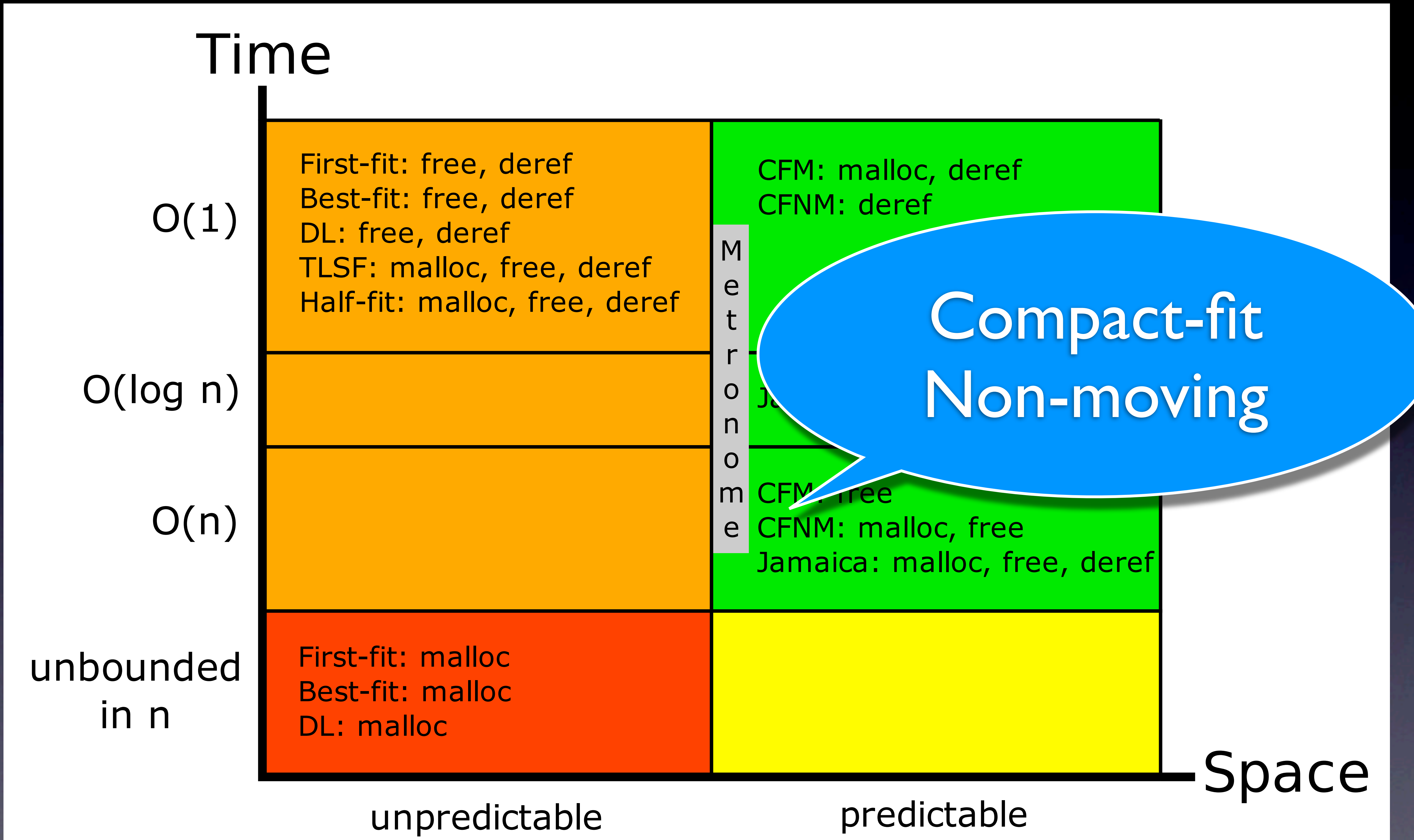
Compact-fit
Moving
Space
unpredictable
predictable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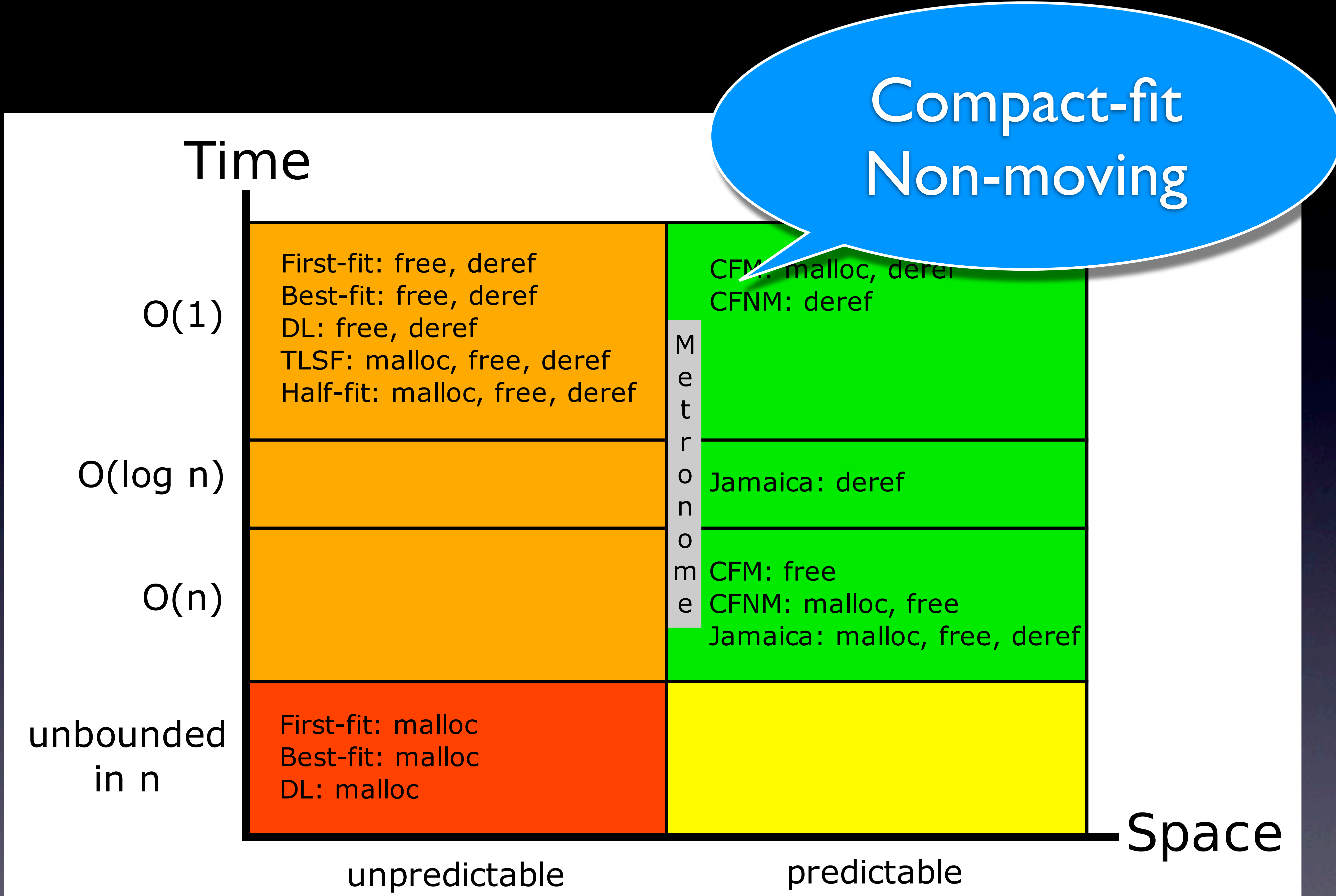
Jamaica: malloc, free, deref
First-fit: malloc
Best-fit: malloc
DL: malloc
Compact-fit
Non-moving
Space
unpredictable
predictable


Compact-fit
Non-moving
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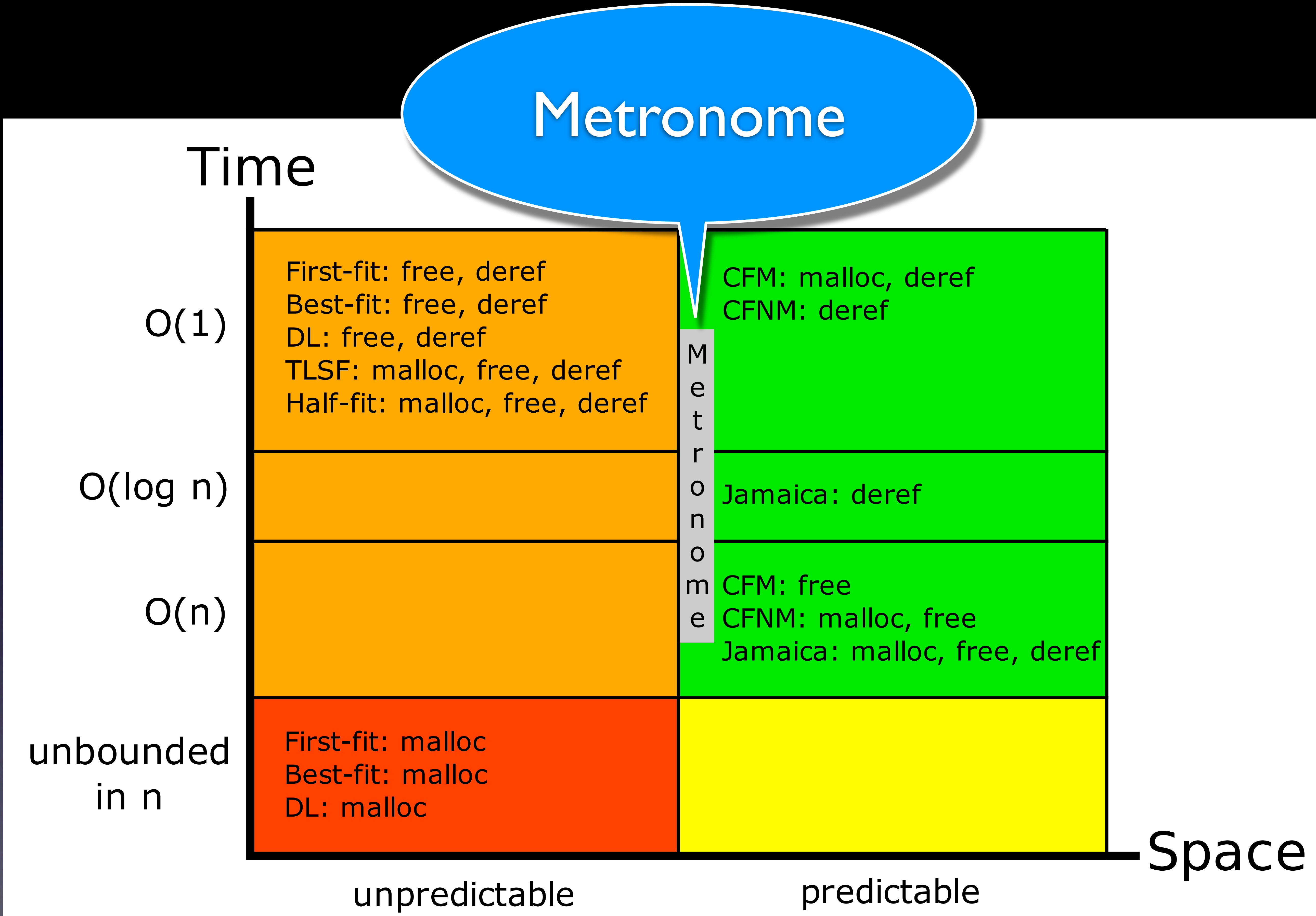
Space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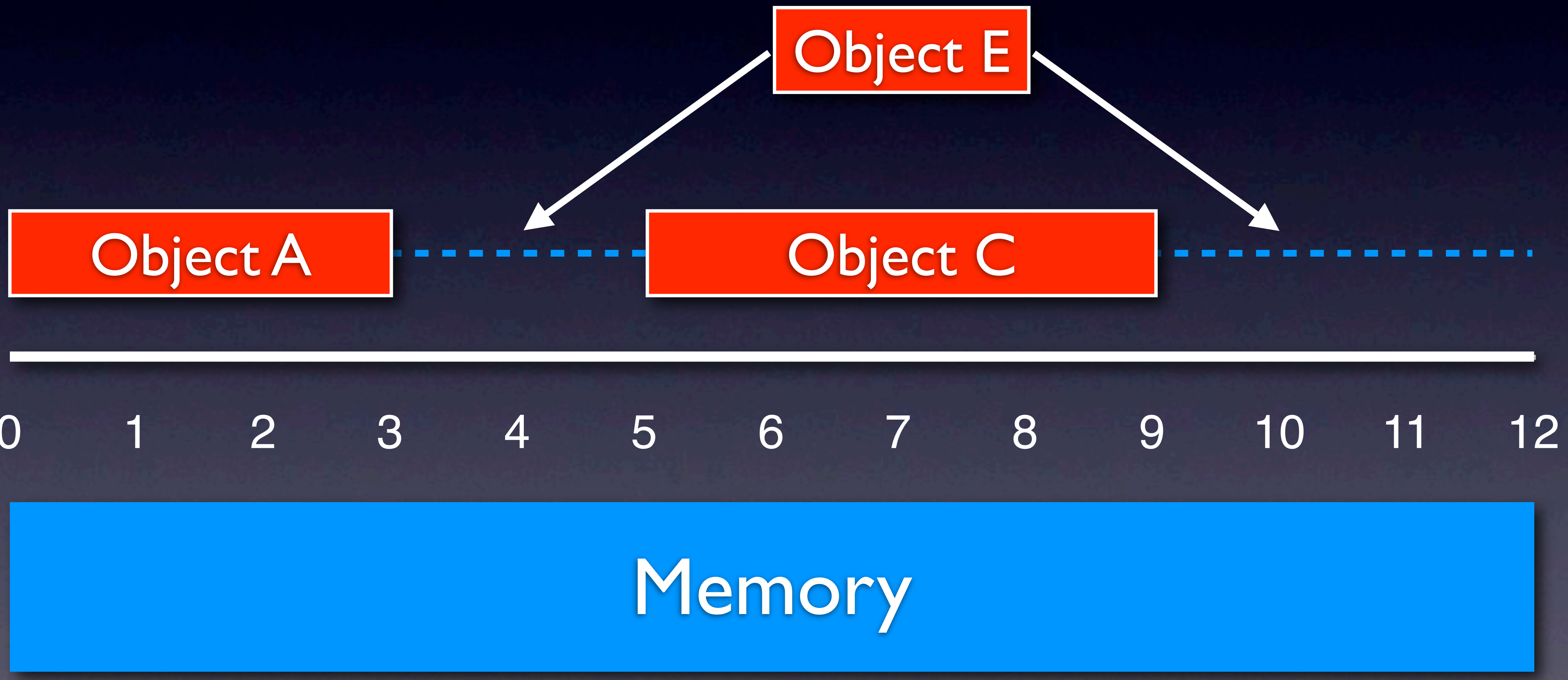
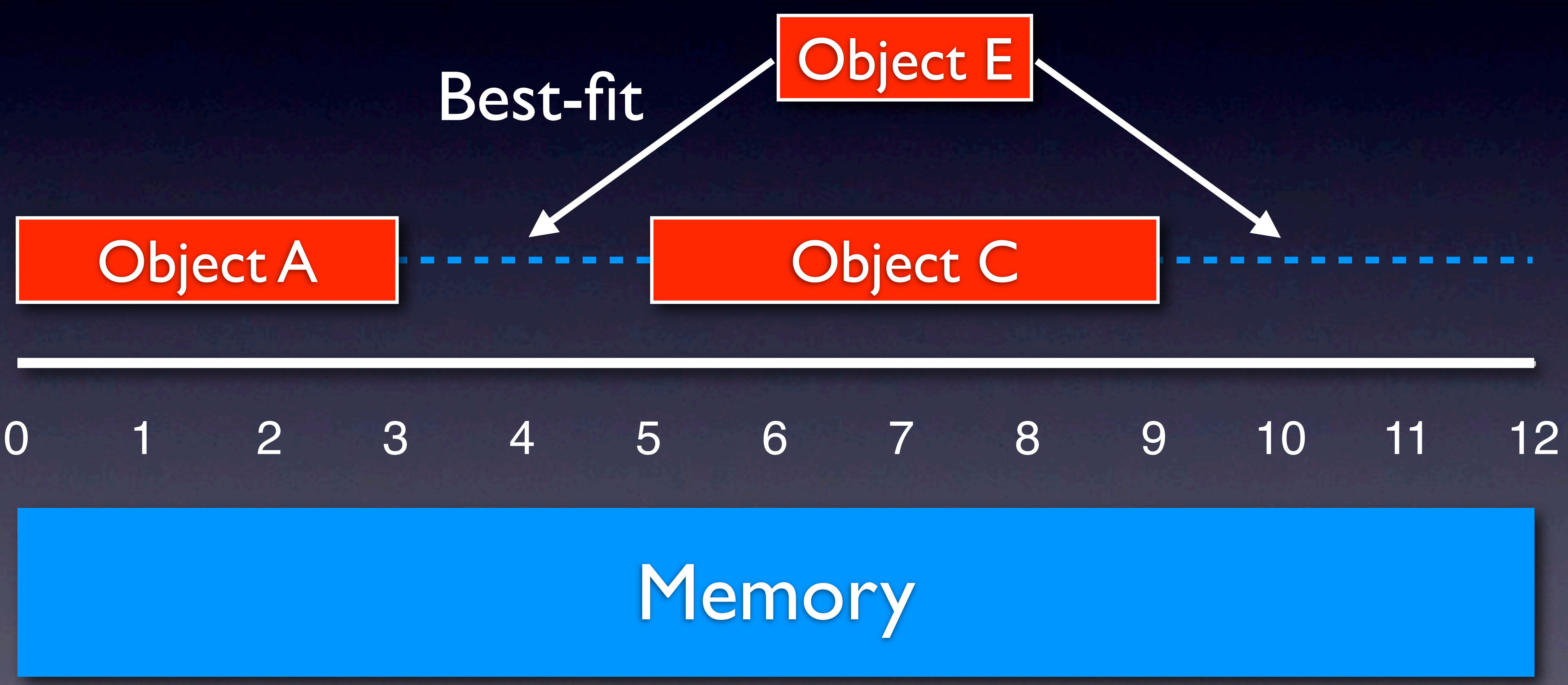
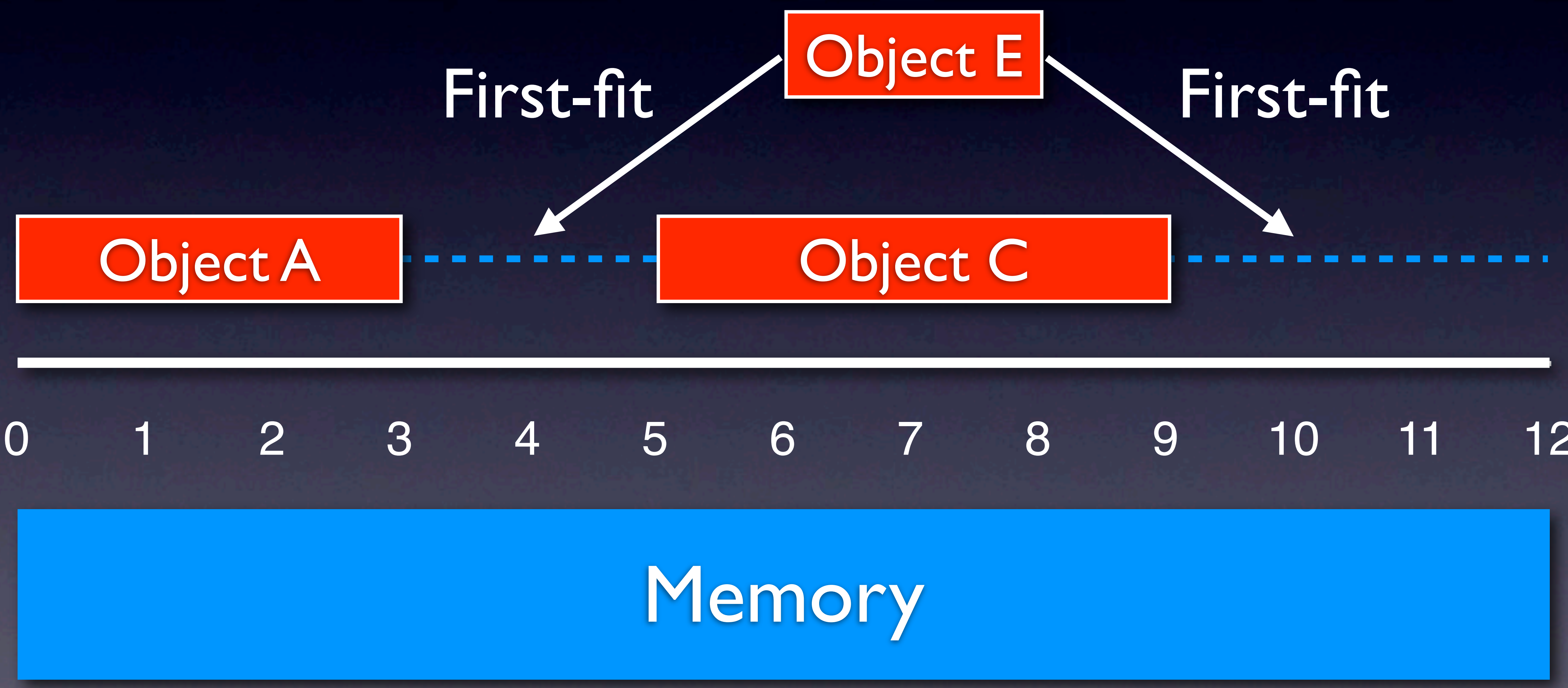
Best-fit versus First-fit



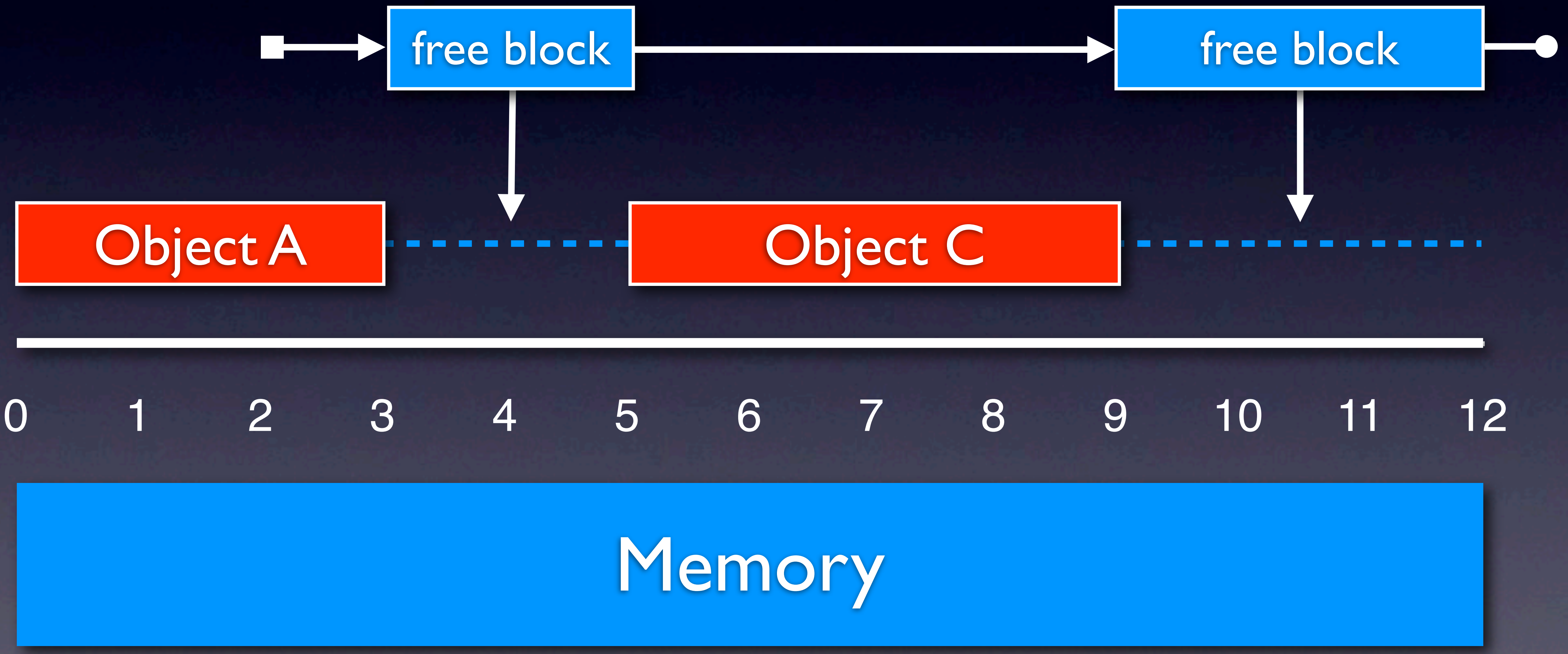
Best-fit versus First-fit



Best-fit versus First-fit



Free List


free block
free block
Object A
Object C
0
1
2
3
4
5
6
7
8
9
10
11
12
Memory

Best-fit, First-fit Complexity

- Allocation:
  - `malloc` may take time proportional to heap size

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- Allocation:
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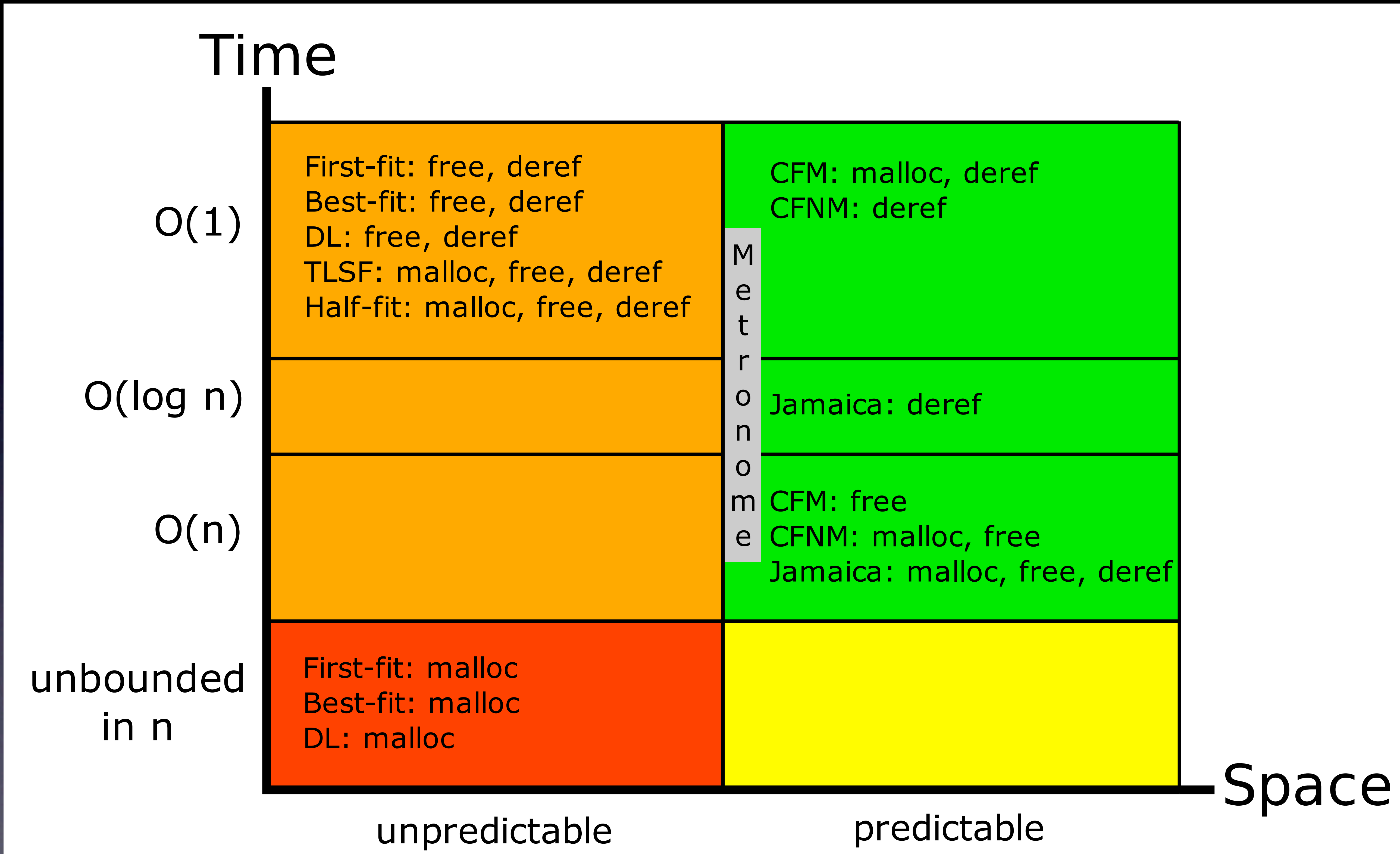
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First-fit: malloc
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Space
unpredictable
predictable

Free List Operations

- Select:
  - `malloc`

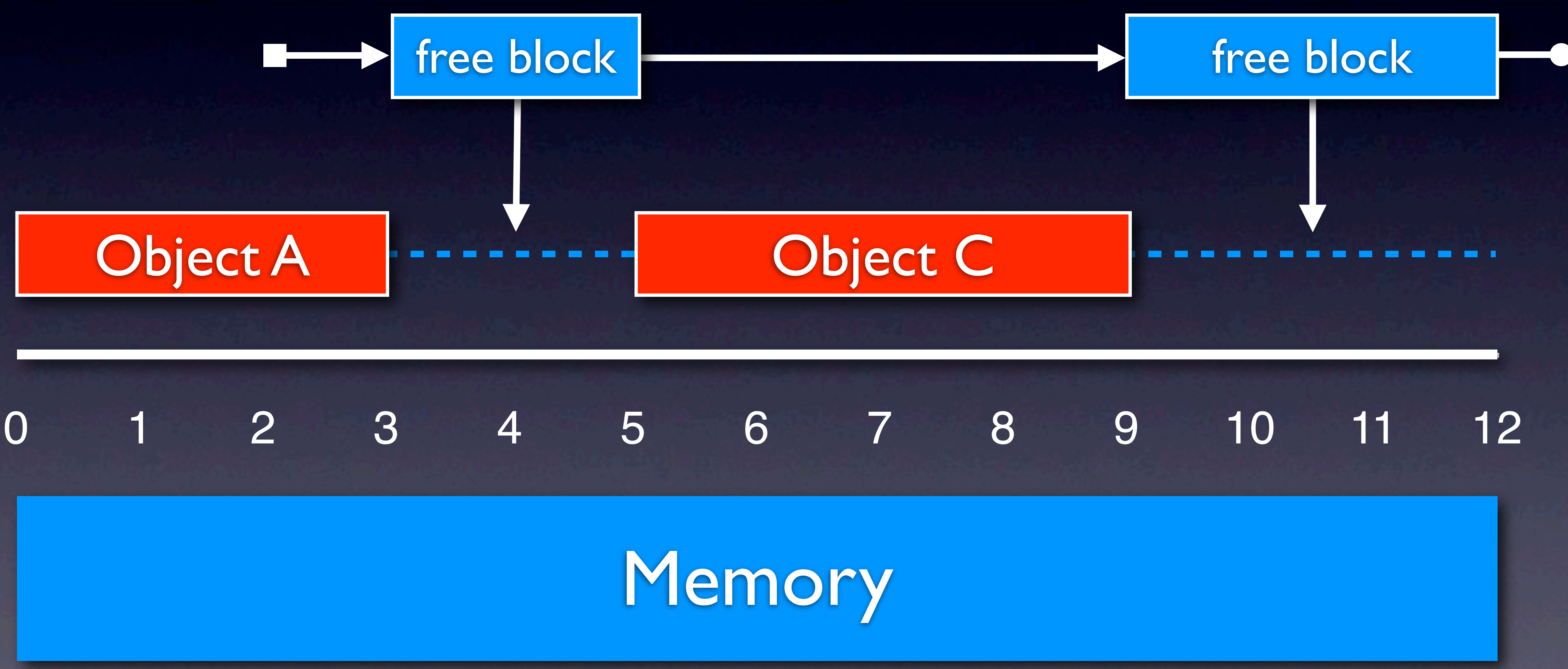
Free List Operations

- Select:
  - `malloc`
- Insert:
  - `free`

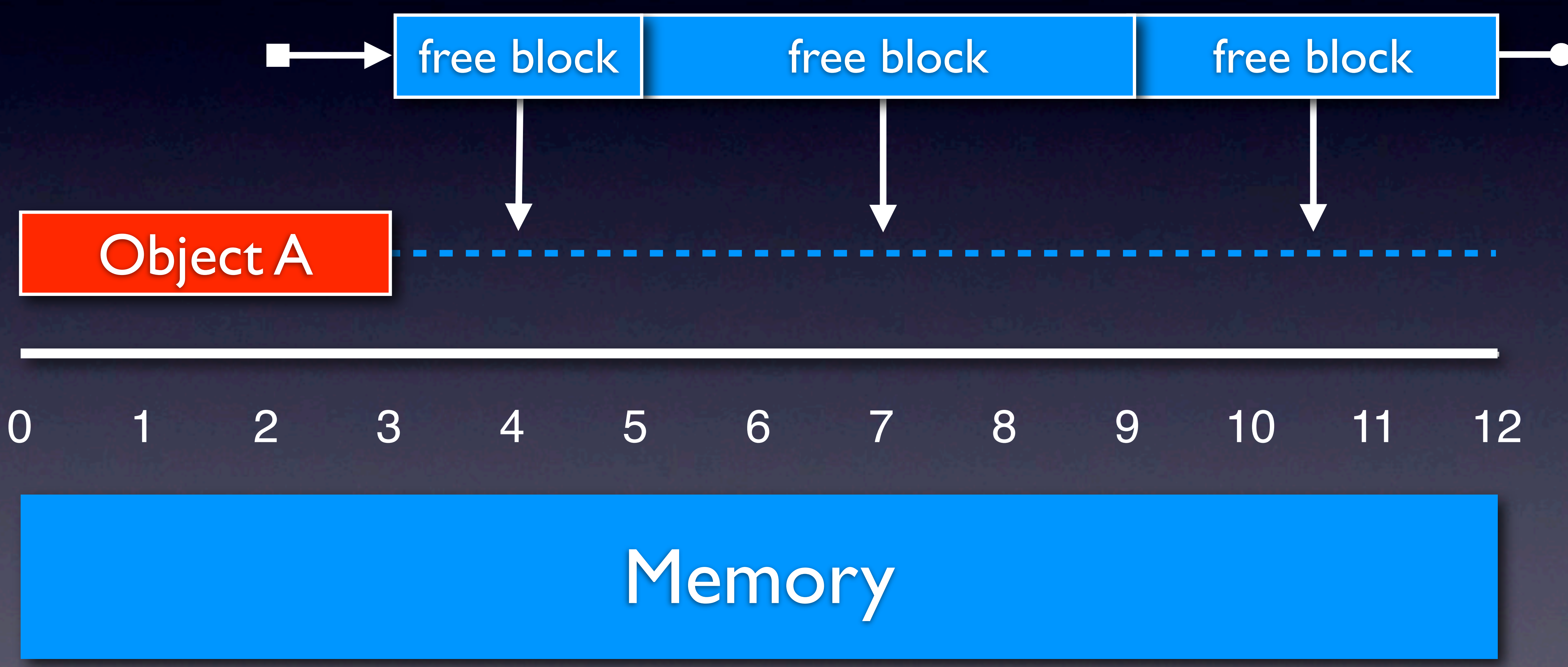
Free List Operations

- Select:
  - `malloc`
- Insert:
  - `free`
- Delete:
  - coalescing

Coalescing

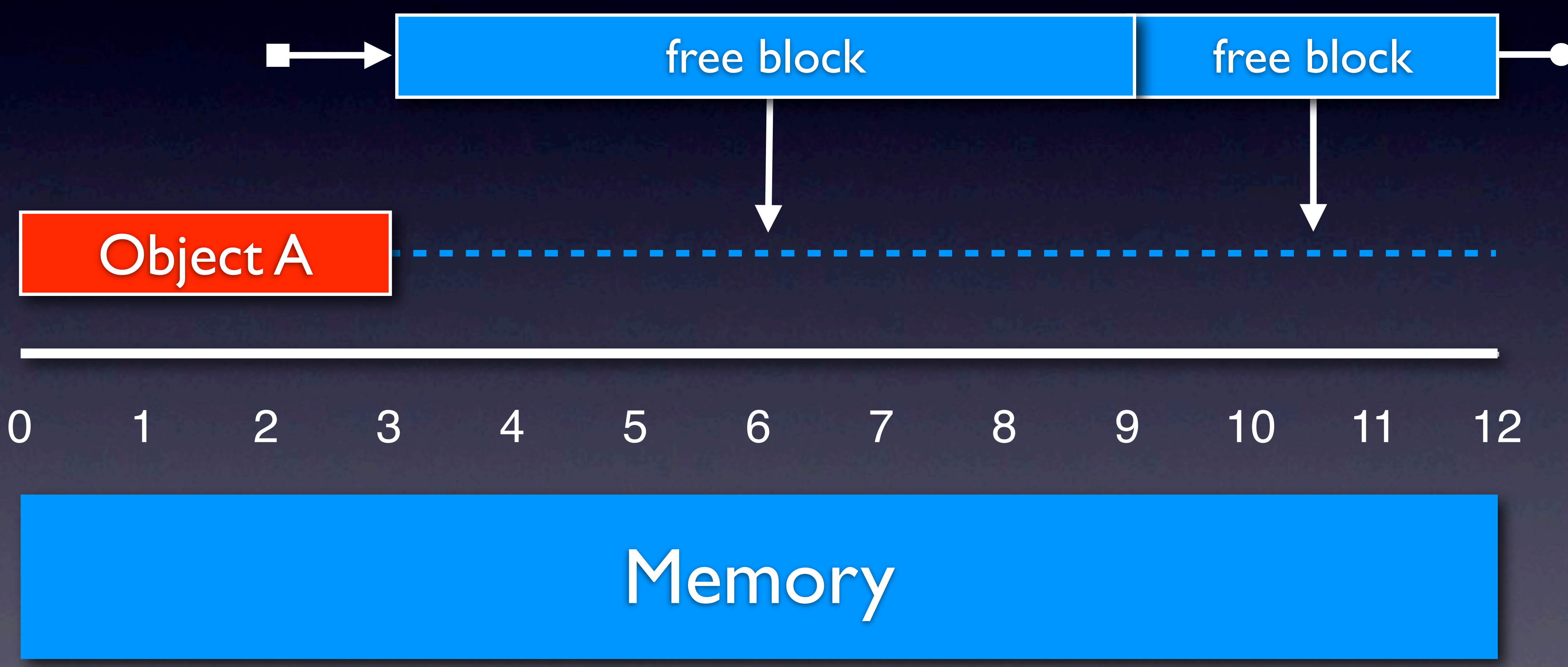

free block
free block
Object A
Object C
0
1
2
3
4
5
6
7
8
9
10
11
12
Memory

Coalescing

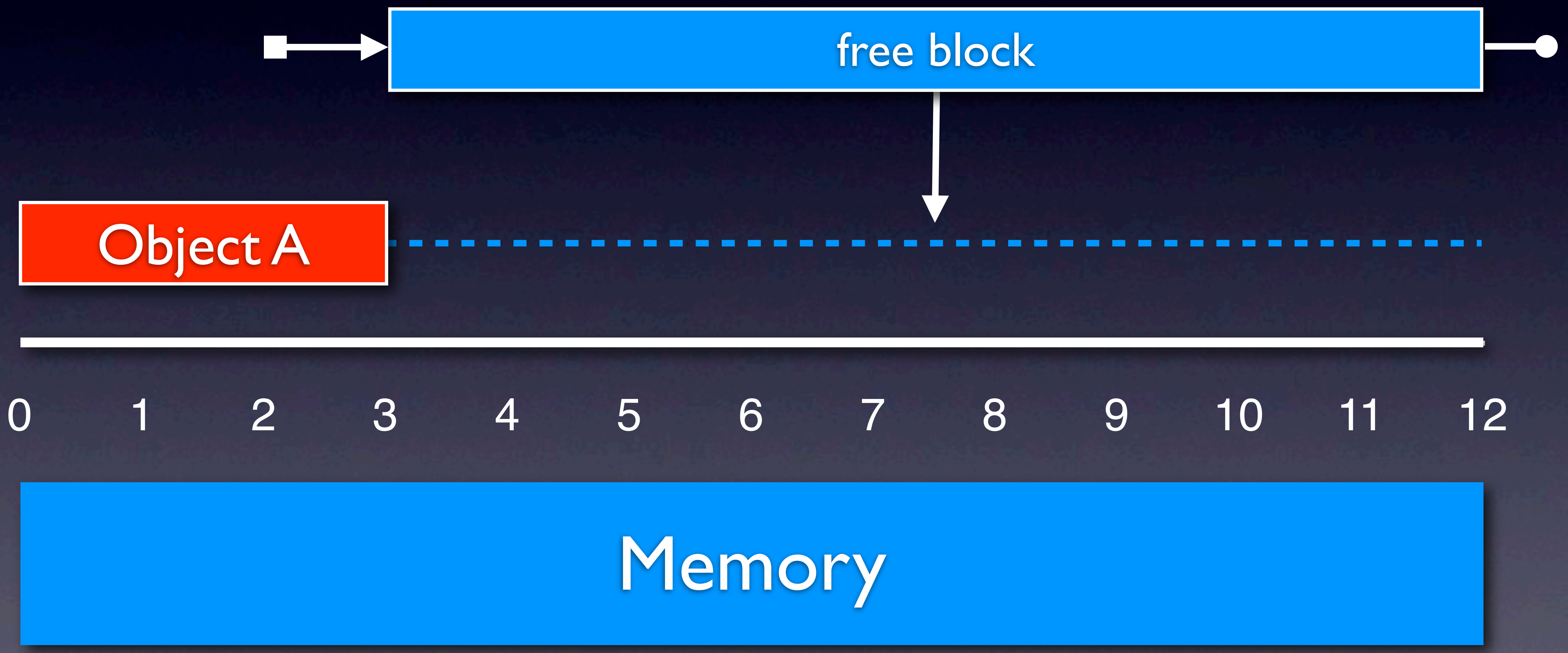


free block
free block
free block
Object A
0 1 2 3 4 5 6 7 8 9 10 11 12
Memory

Coalescing



Coalescing



List Representations

- List: singly-linked or doubly-linked (using boundary tags)

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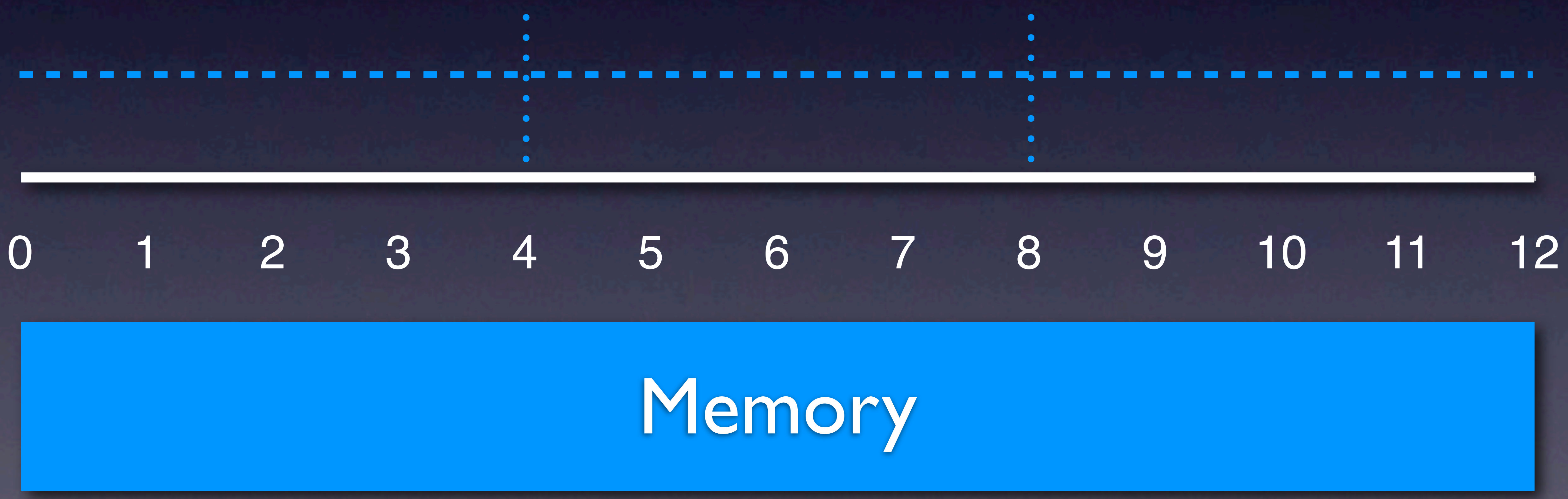
- Segregated lists: array of lists for different sizes
- Buddy systems: split blocks in powers of two (called buddies if same size)
- Indexed lists: trees, bitmaps
- Hybrid: Doug Lea's allocator

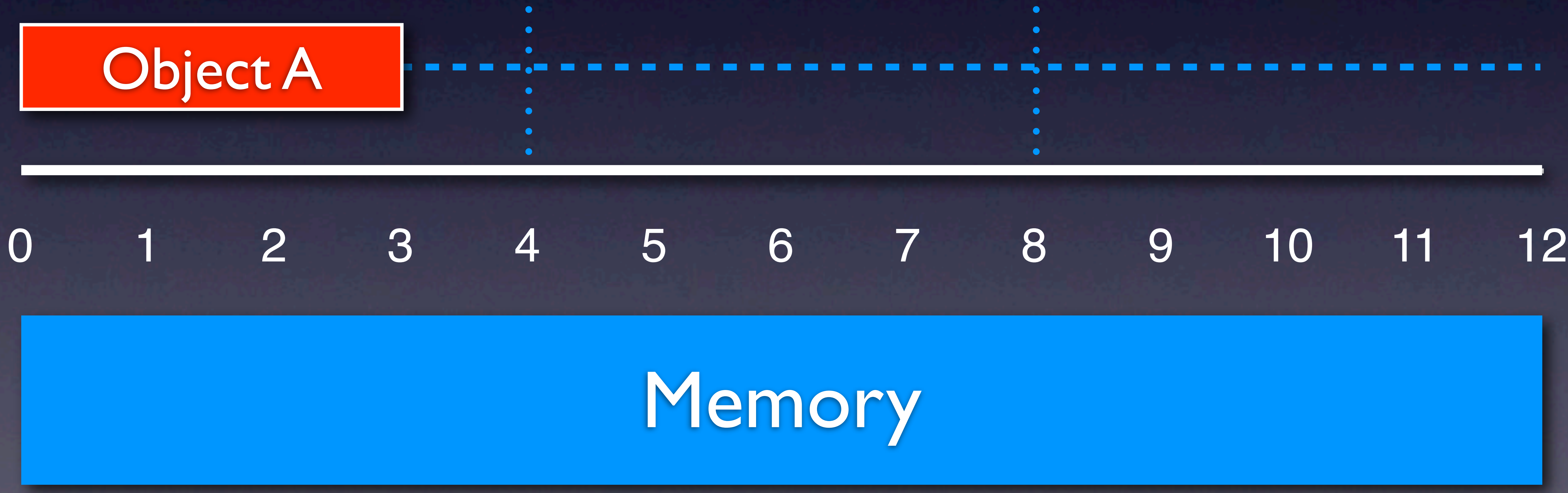
DL Complexity

- Allocation:
  - `malloc` may take time proportional to heap size
- Deallocation:
  - `free` takes constant time
- Access:
  - `read` and `write` take constant time
- <u>Unpredictable</u> fragmentation

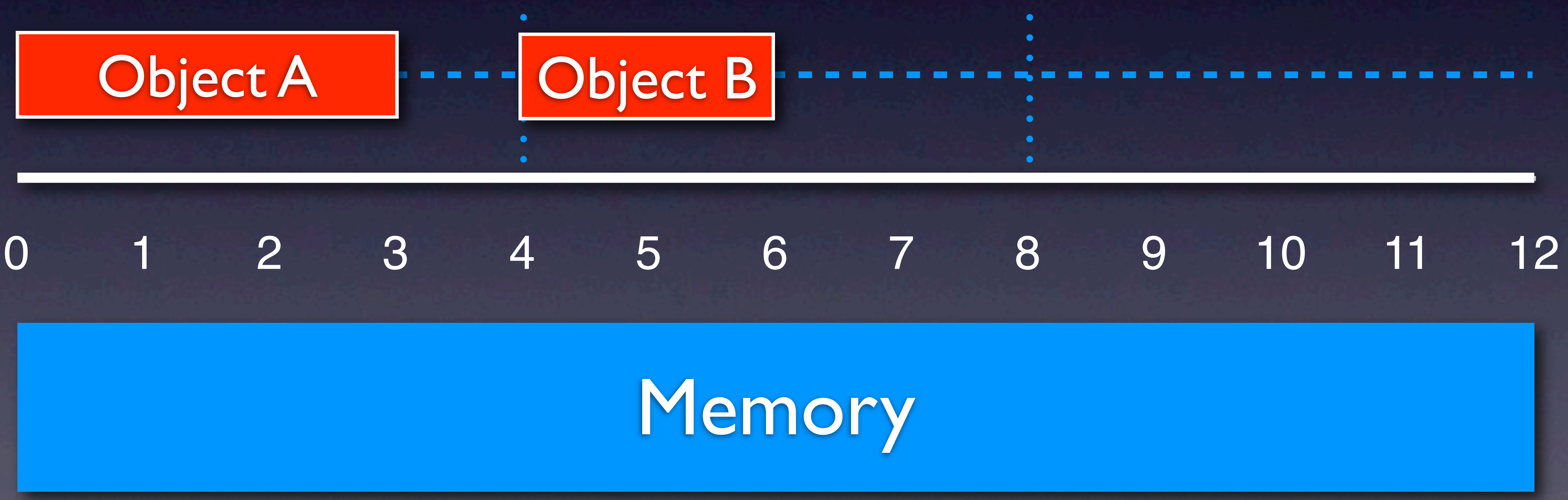
Partitioning



Partitioning



Partitioning



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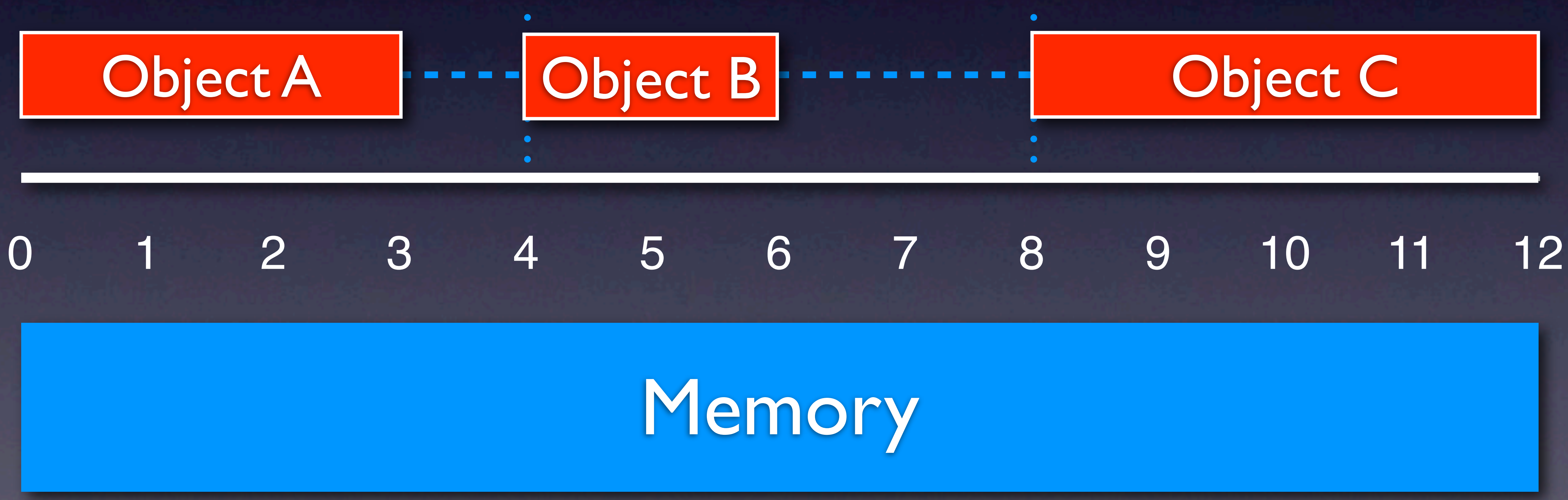
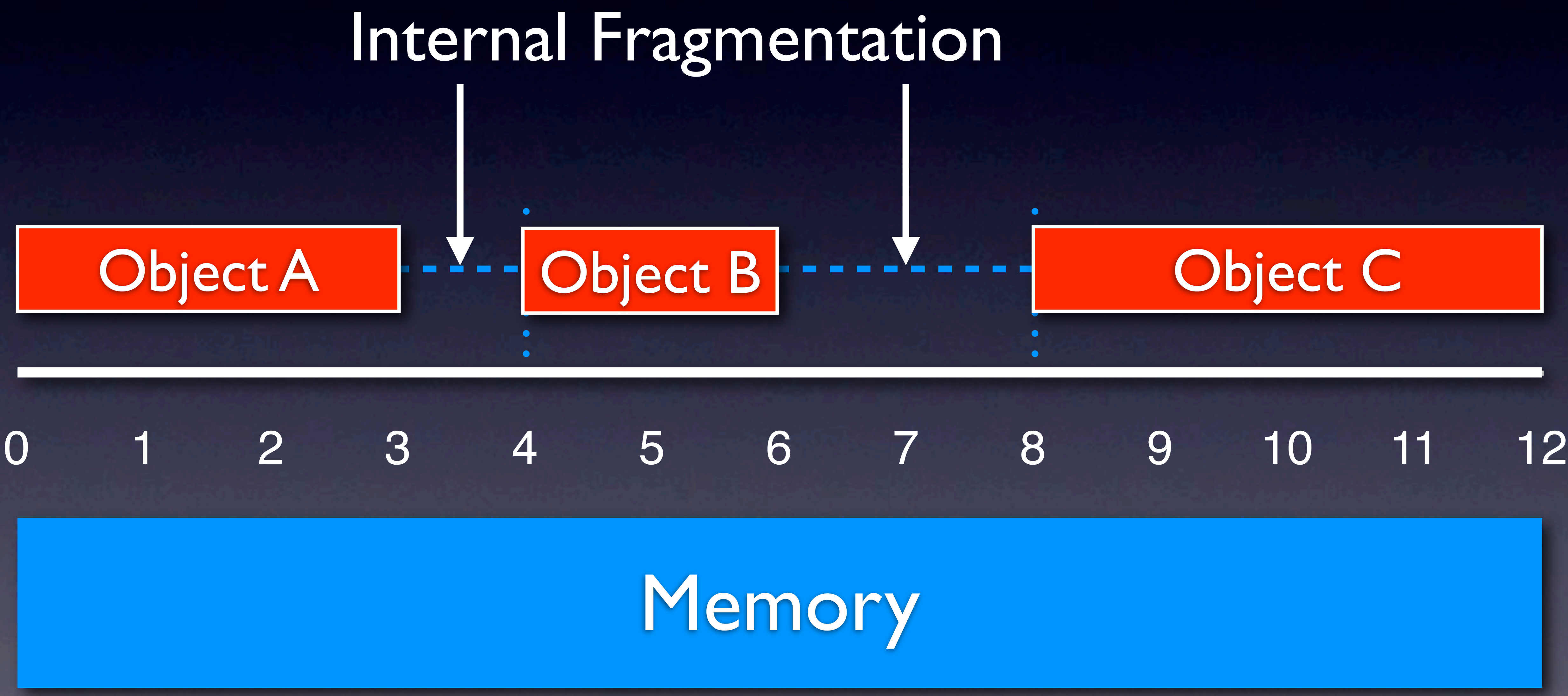
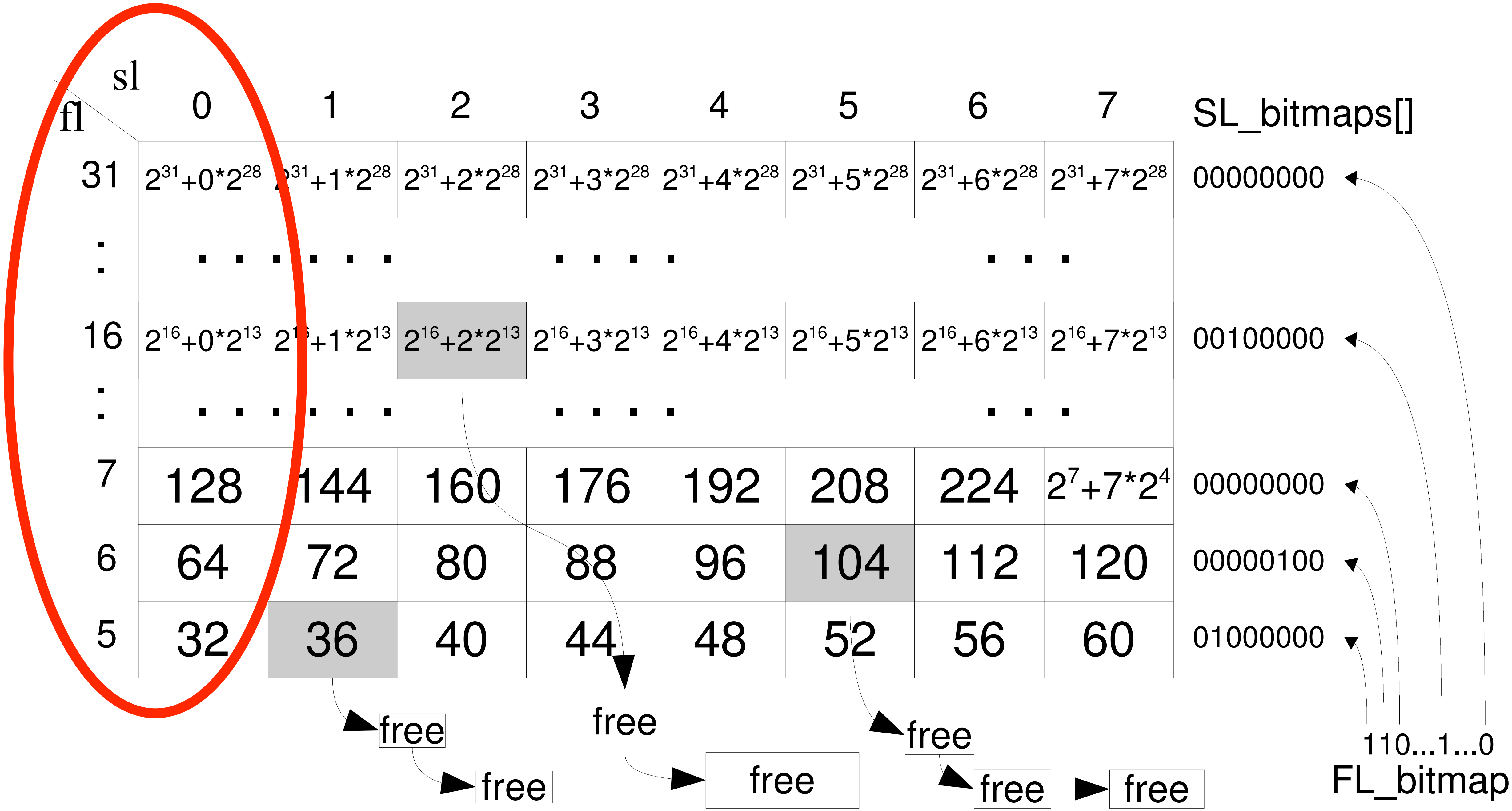


Partitioning



There is a <u>trade-off</u> between external and internal fragmentation

Half-fit

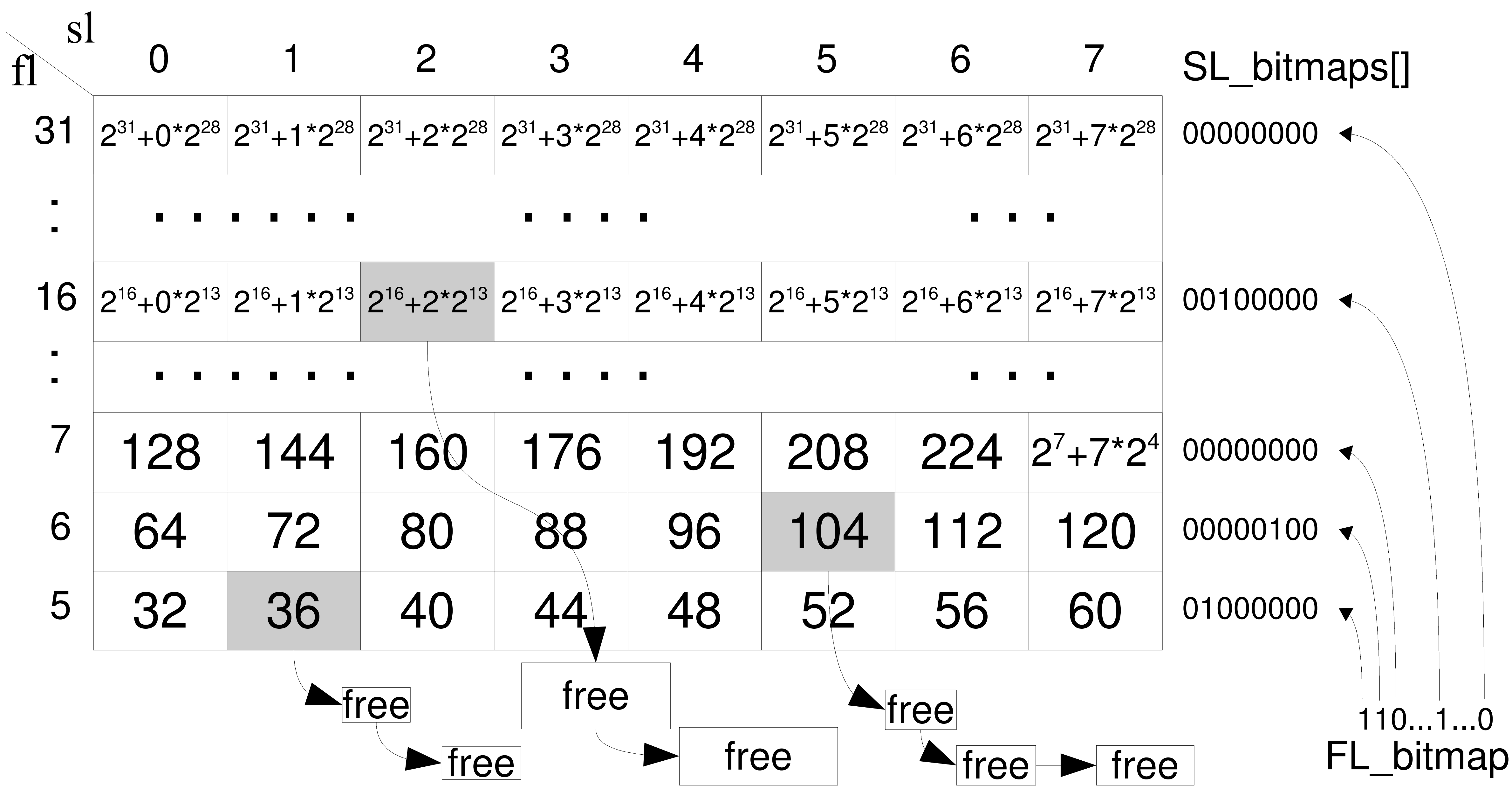


[Masmano et al., In J. of Real-Time Systems, 2008]

Half-fit Complexity

- Allocation:
  - `malloc` takes constant time
- Deallocation:
  - `free` takes constant time
- Access:
  - `read` and `write` take constant time
- Unpredictable fragmentation

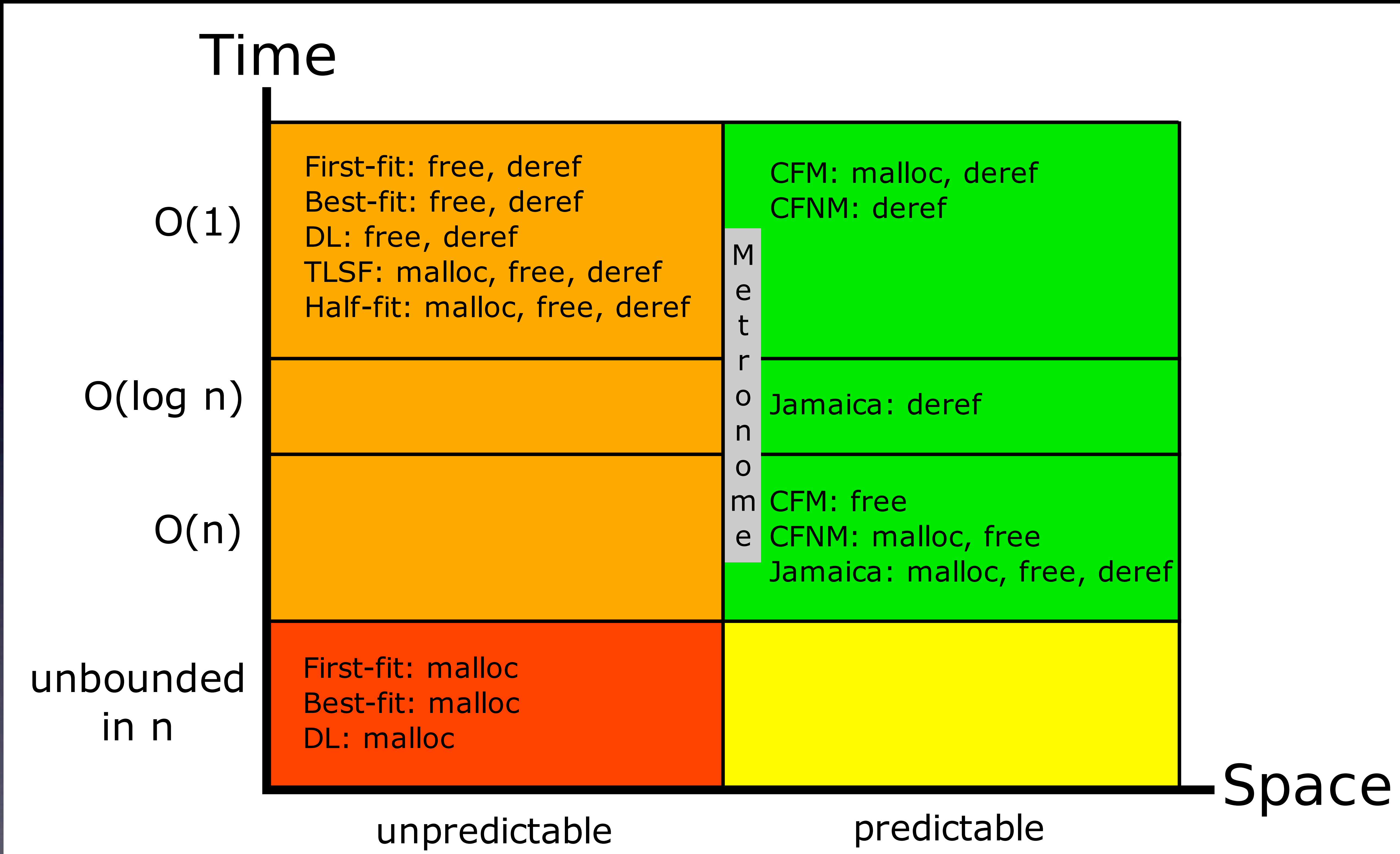
Two-level Segregated Fit (TLSF)



[Masmano et al., In J. of Real-Time Systems, 2008]

TLSF Complexity

- Allocation:
  - `malloc` takes constant time
- Deallocation:
  - `free` takes constant time
- Access:
  - `read` and `write` take constant time
- Unpredictable fragmentation (yet better than HF)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Space
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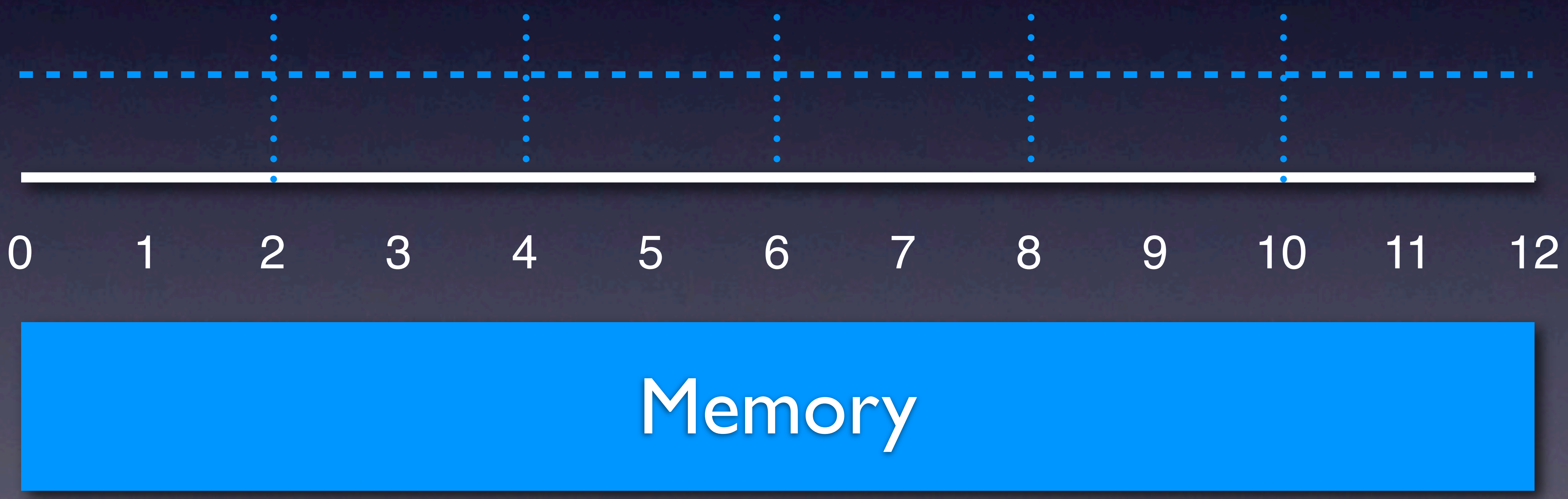
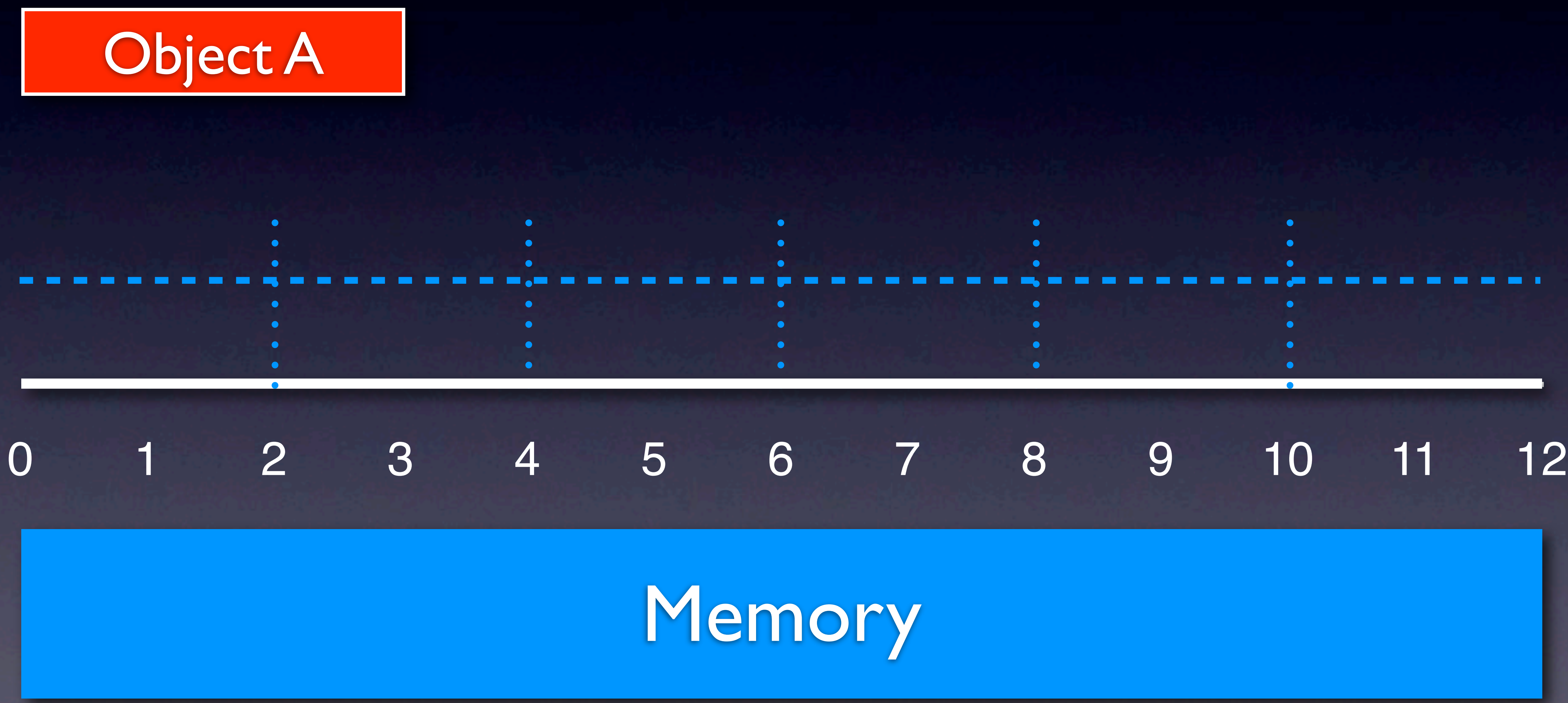
predictable

Jamaica



Jamaica



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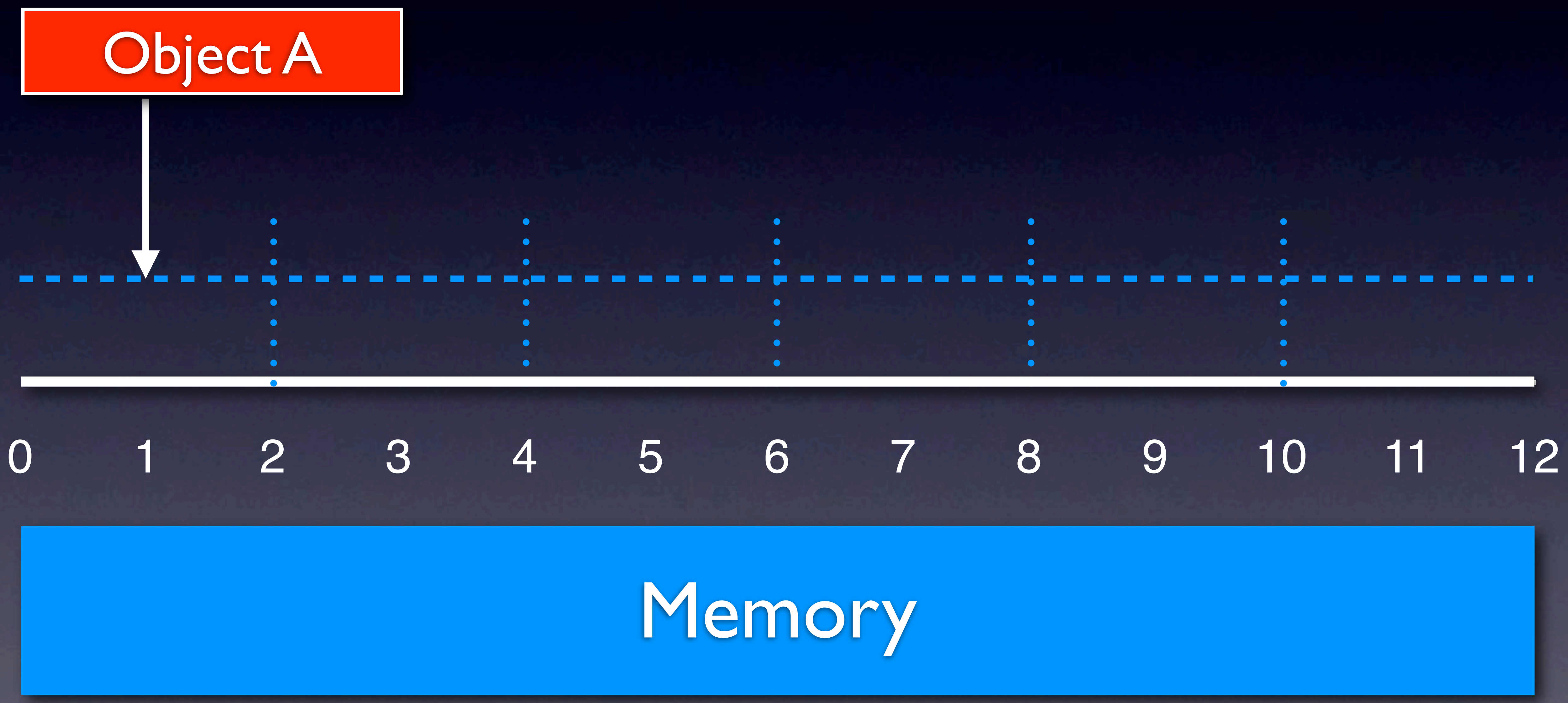

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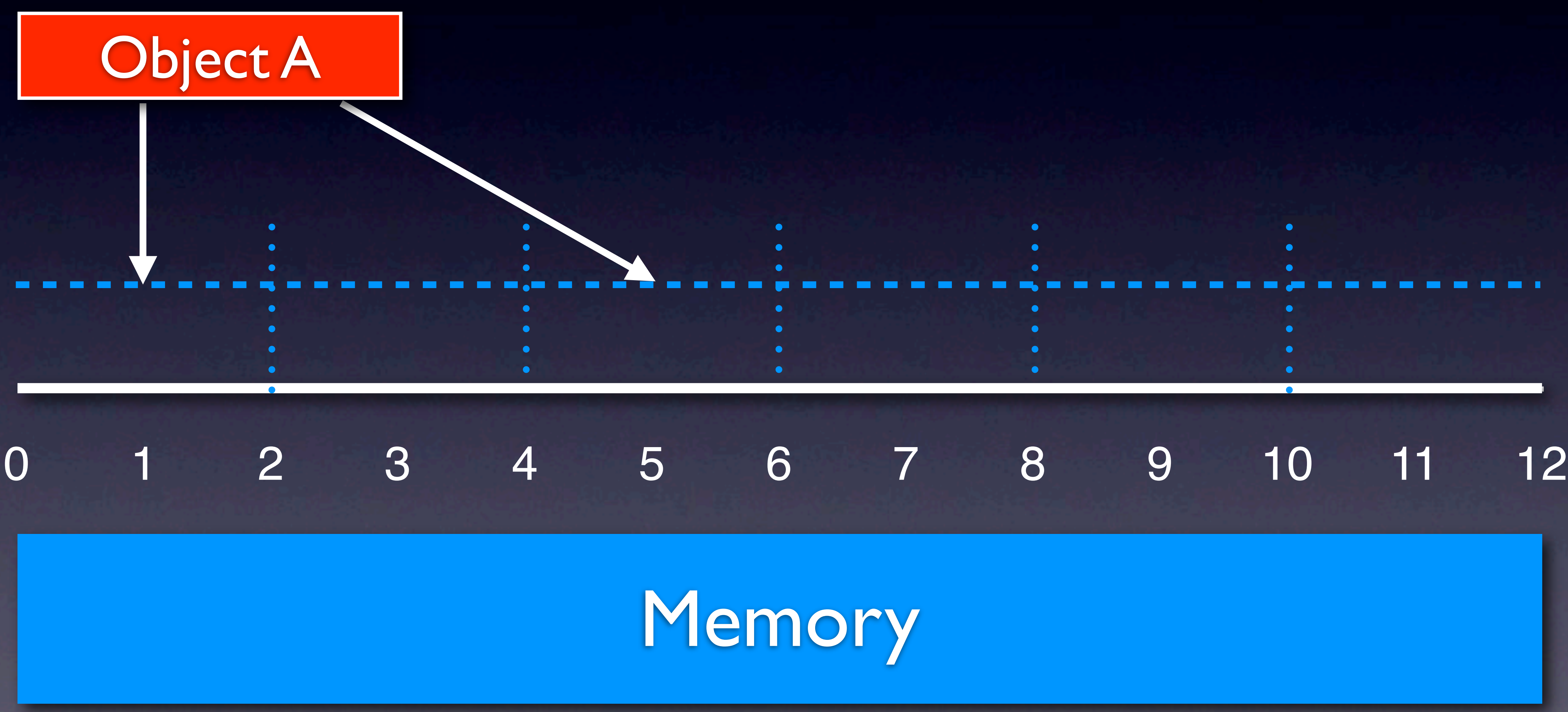

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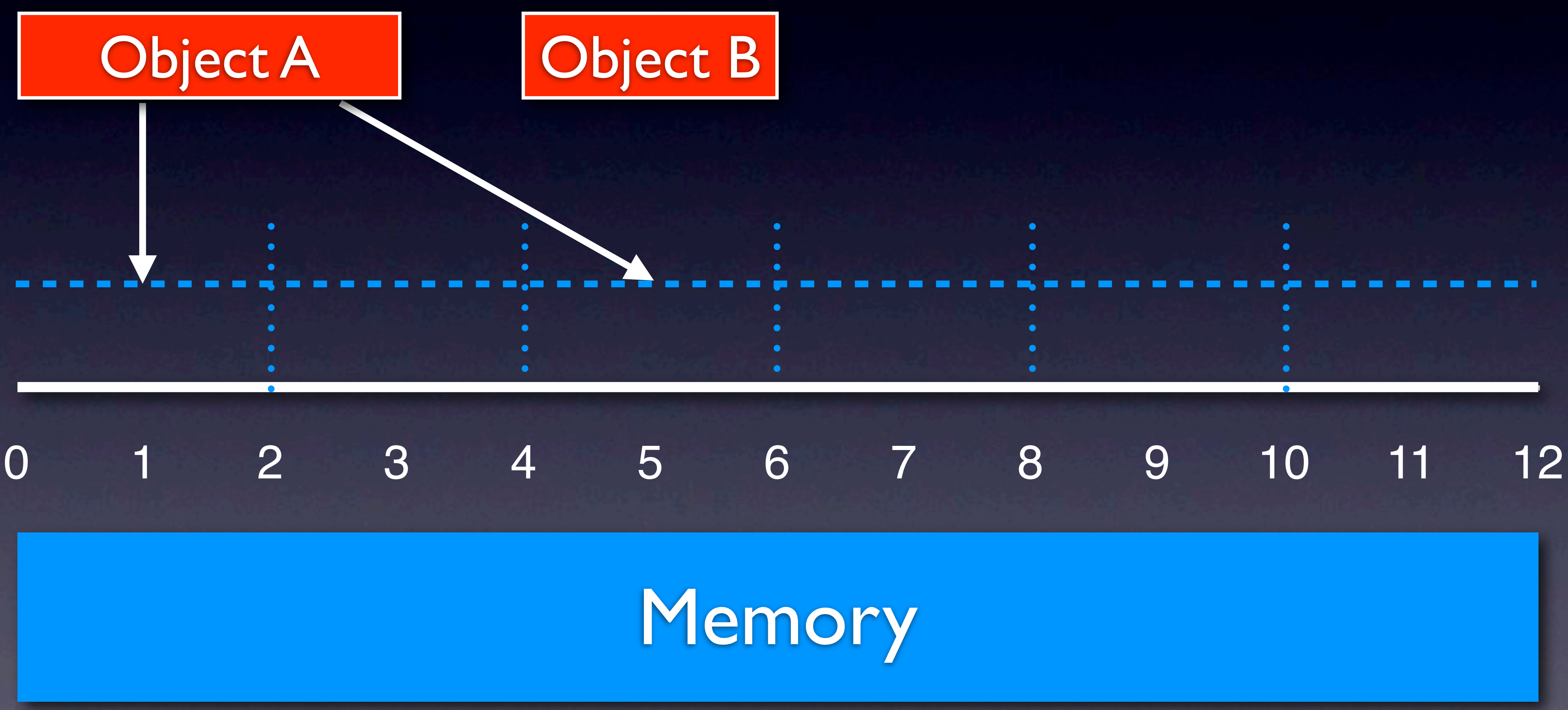

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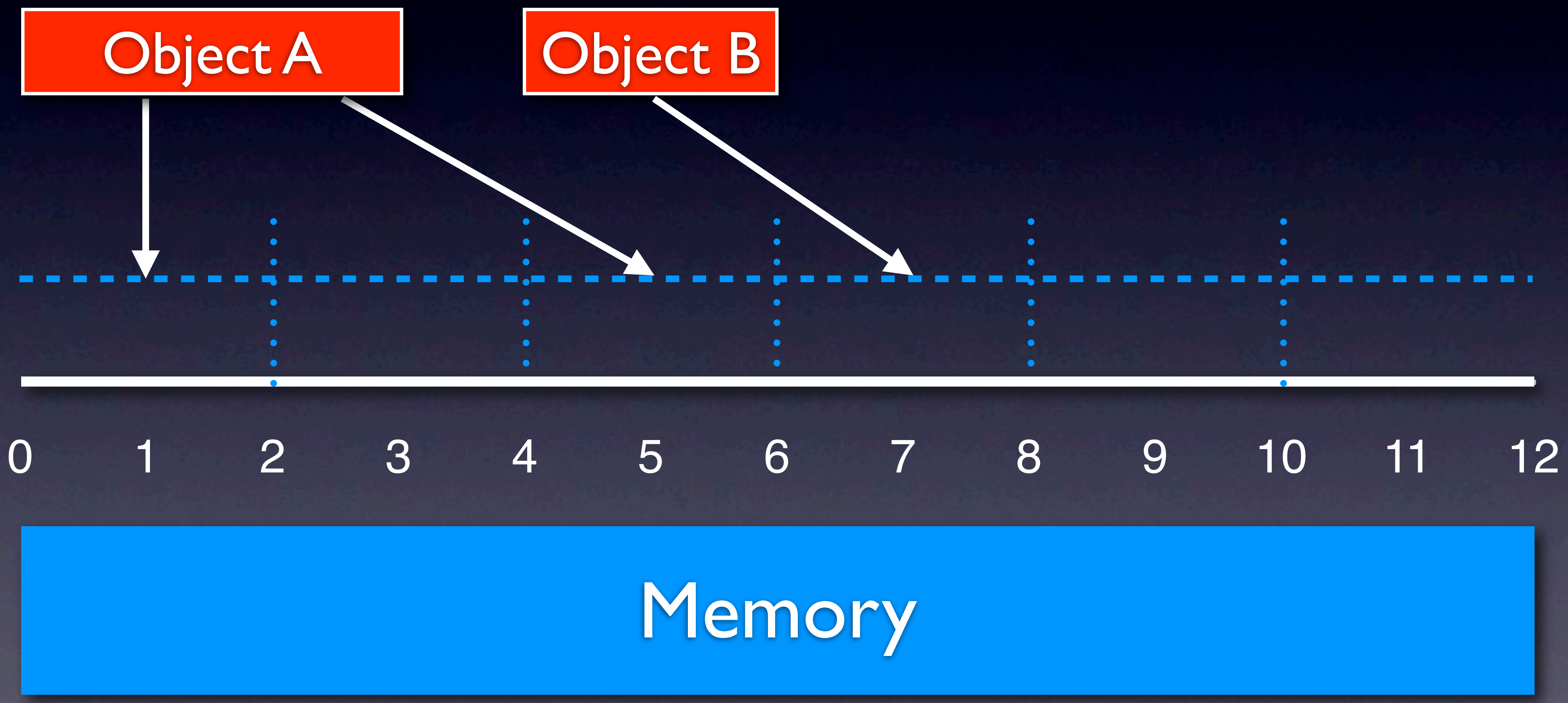

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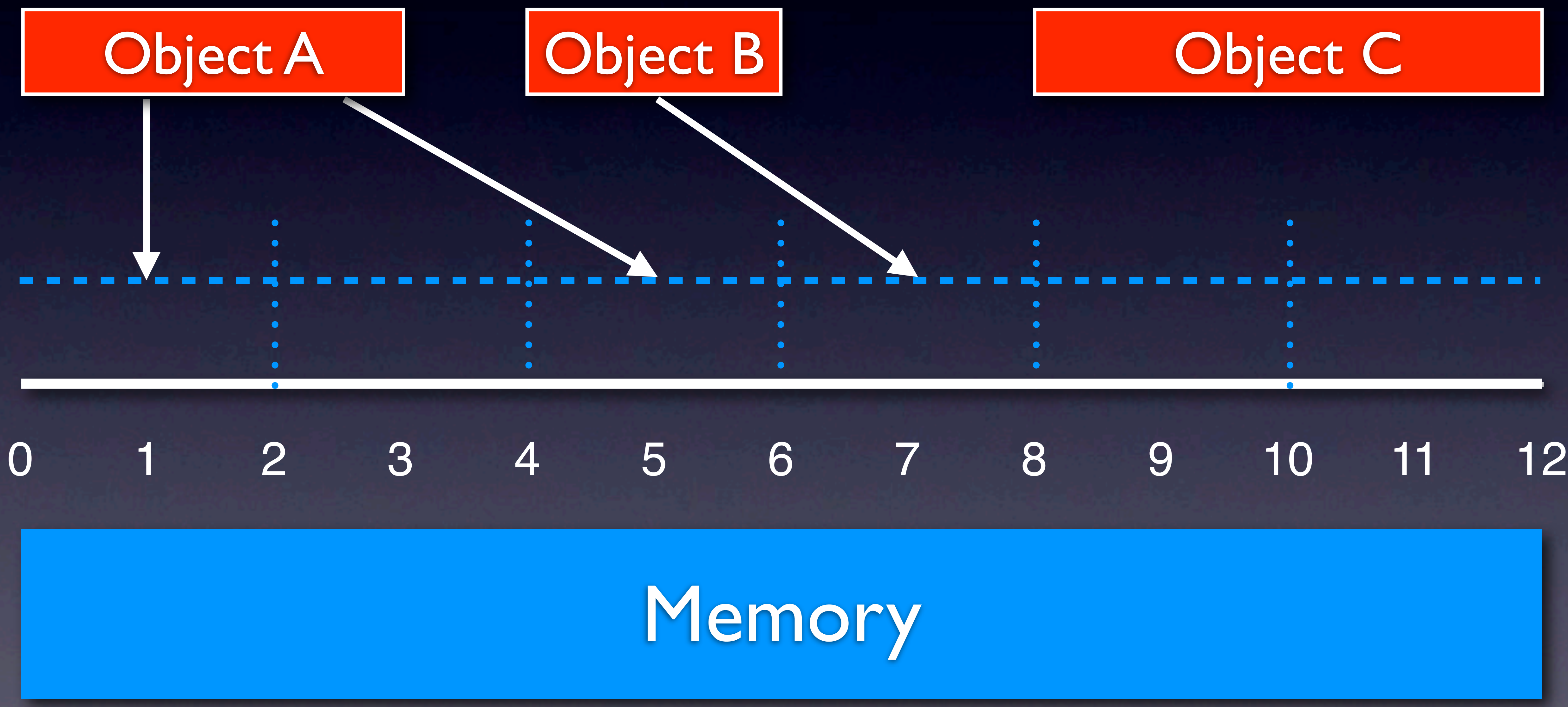

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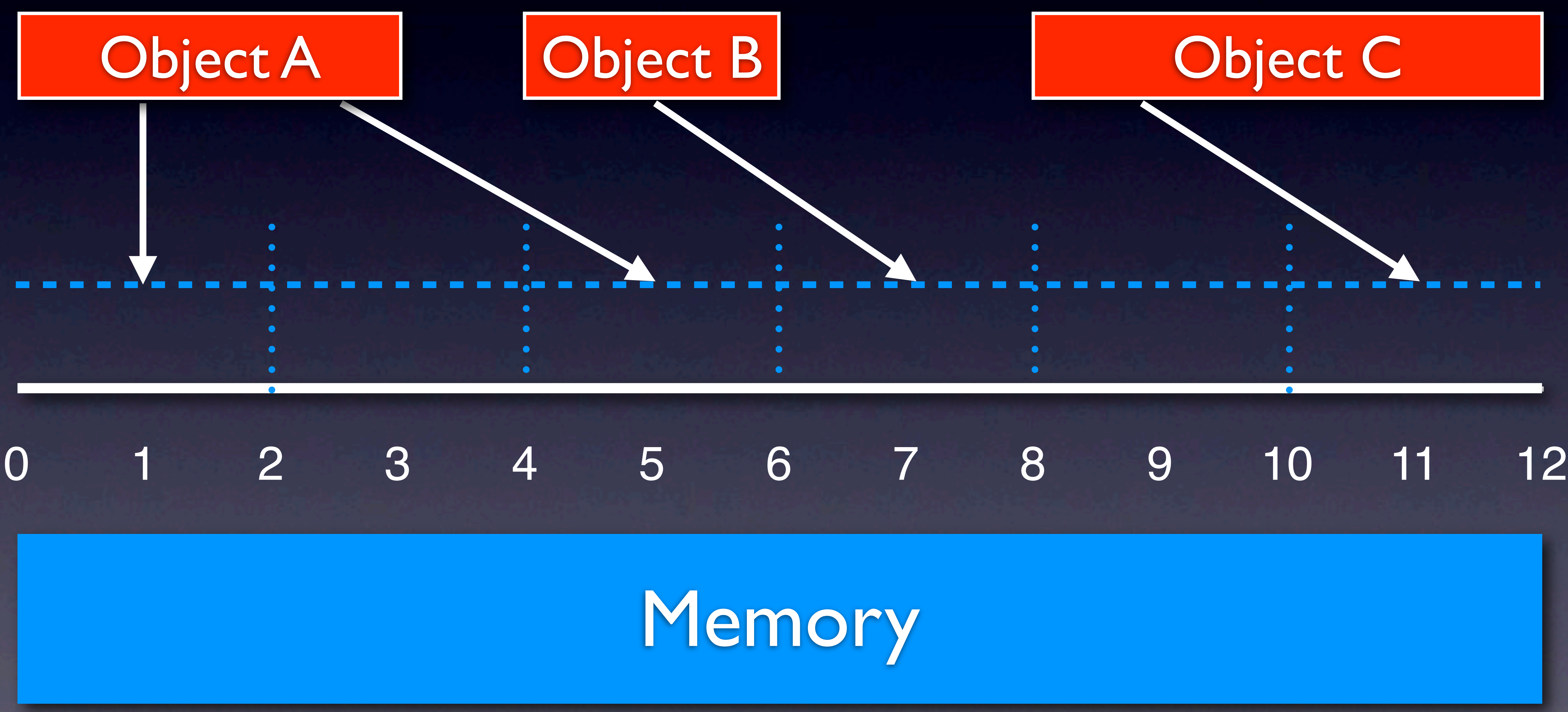

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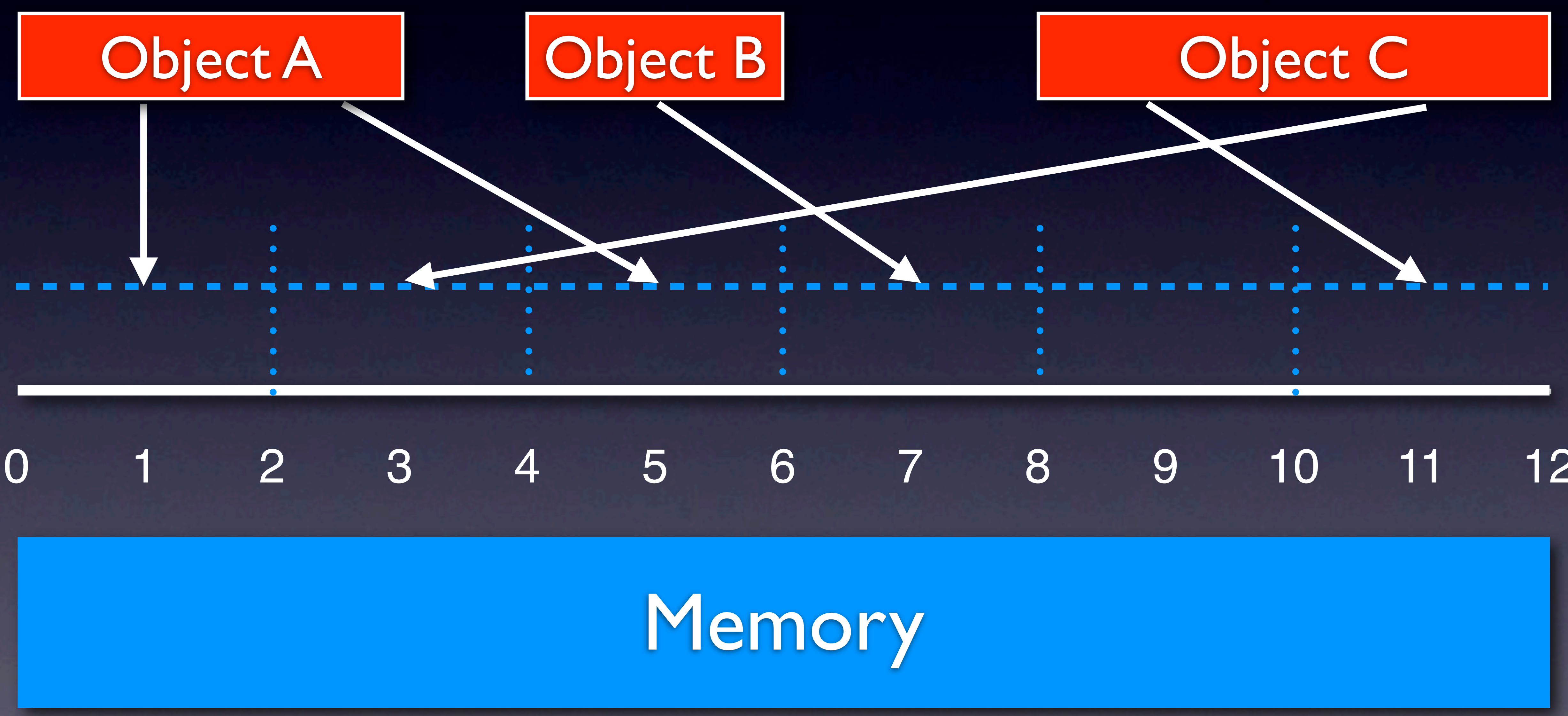
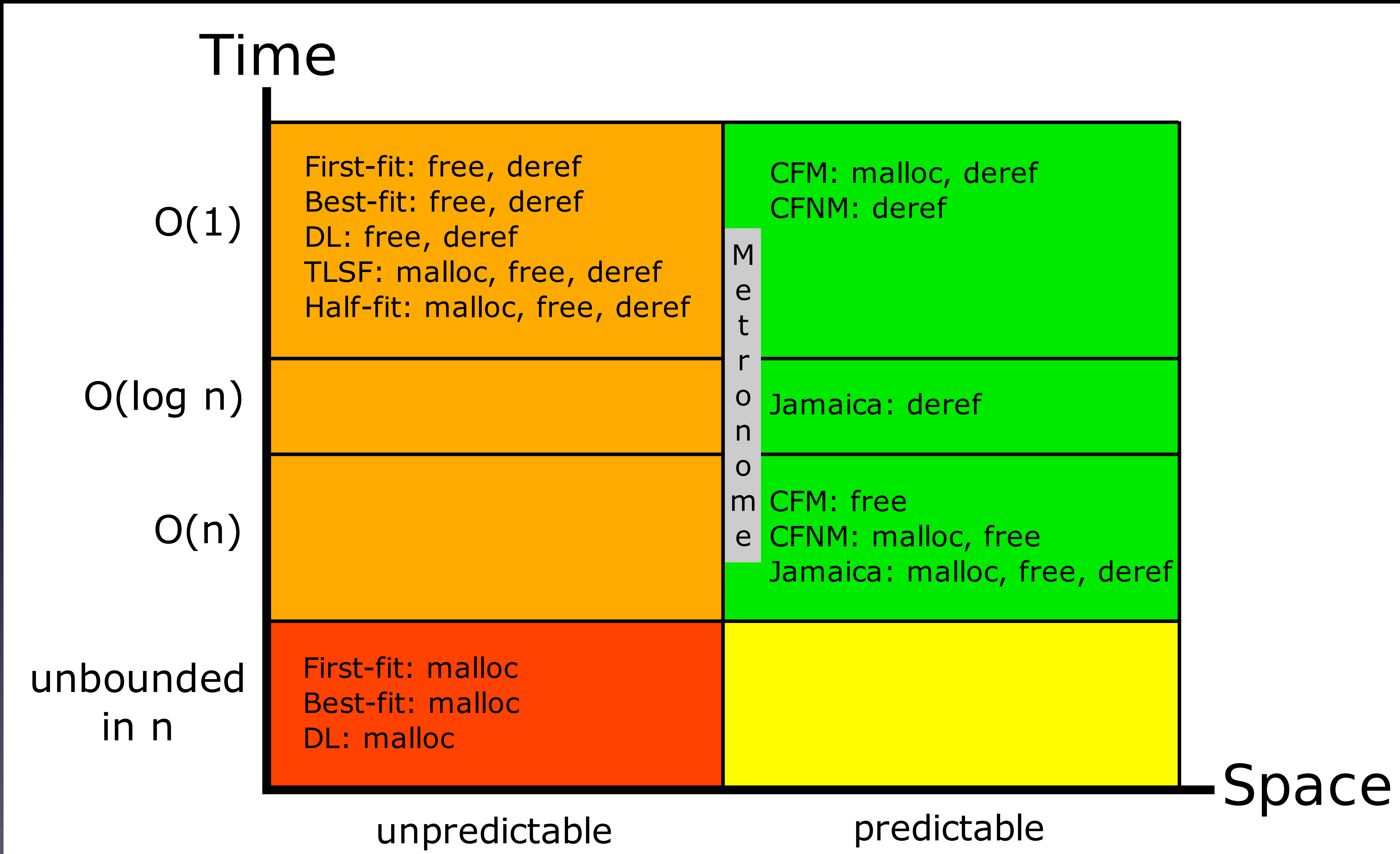


Jamaica Complexity

- Allocation:
  - `malloc(n)` takes time proportional to `n`
- Deallocation:
  - `free(n)` takes time proportional to `n`
- Access:
  - `read` and `write` take time <u>proportional</u> to `n`
- Predictable fragmentation



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unpredictable
predictable
Space

Introduction to Compact-fit

Concurrent Compact-fit

Concurrency & Scalability
versus
Fragmentation & Compaction

Questions

- Does allocation/deallocation throughput scale with multiple processors?

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- Which aspects influence scalability?

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- Does compaction of large objects harm system latency?

Questions

- Does allocation/deallocation throughput scale with multiple processors?
- Which aspects influence scalability?
- Does compaction of large objects harm system latency?
- Does concurrency and incrementality affect memory consumption?

Partial Compaction

- Per-size-class partial compaction bound κ bounds size-class fragmentation:
  - $\kappa = 1$: fully compacting
  - $1 < \kappa < \infty$: partially compacting
  - $\kappa = \infty$: non-compacting

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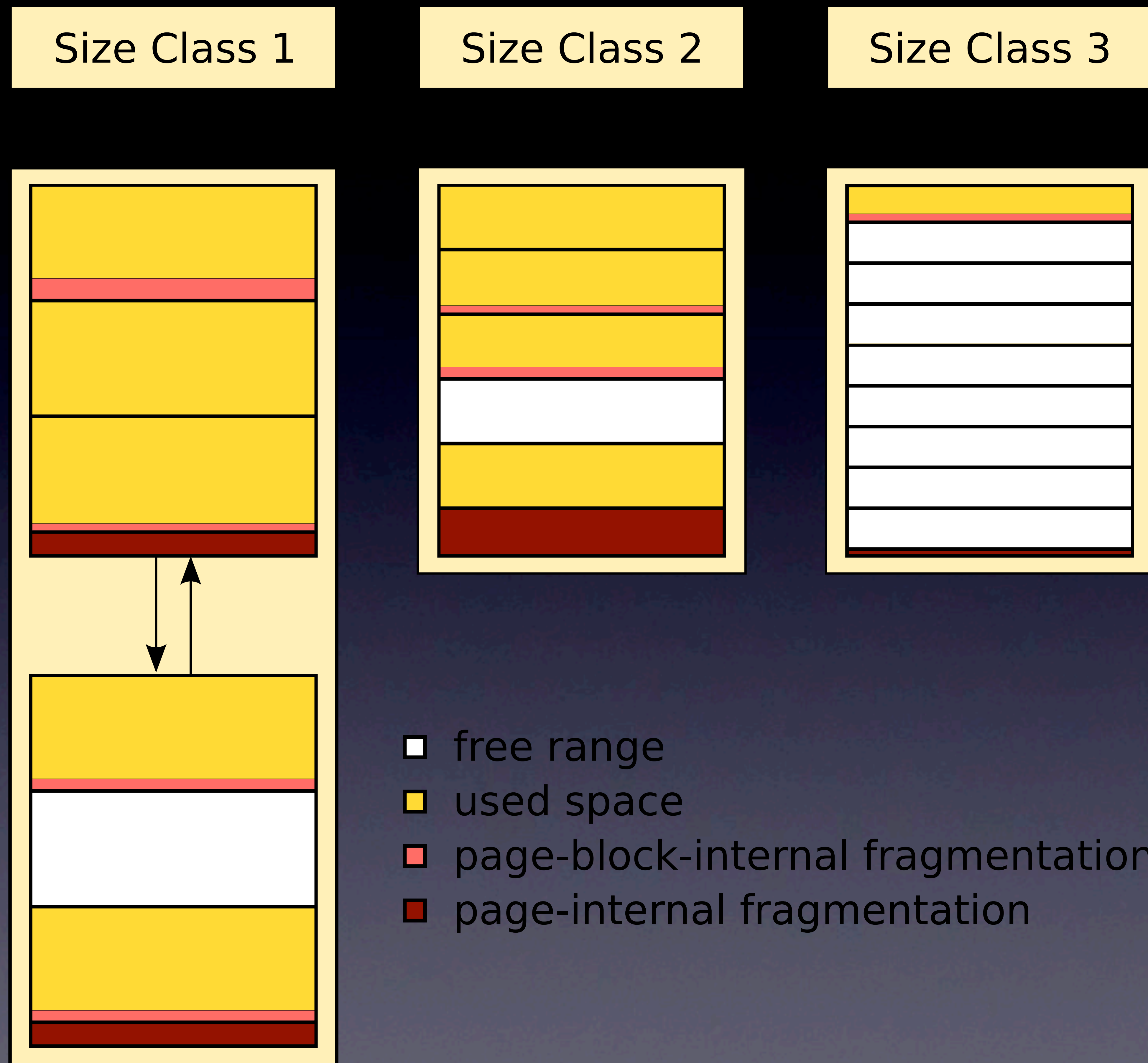
  - $\kappa = \infty$: non-compacting
- Non-compacting CF can be optimized by not using abstract addresses

Fragmentation through Partitioning

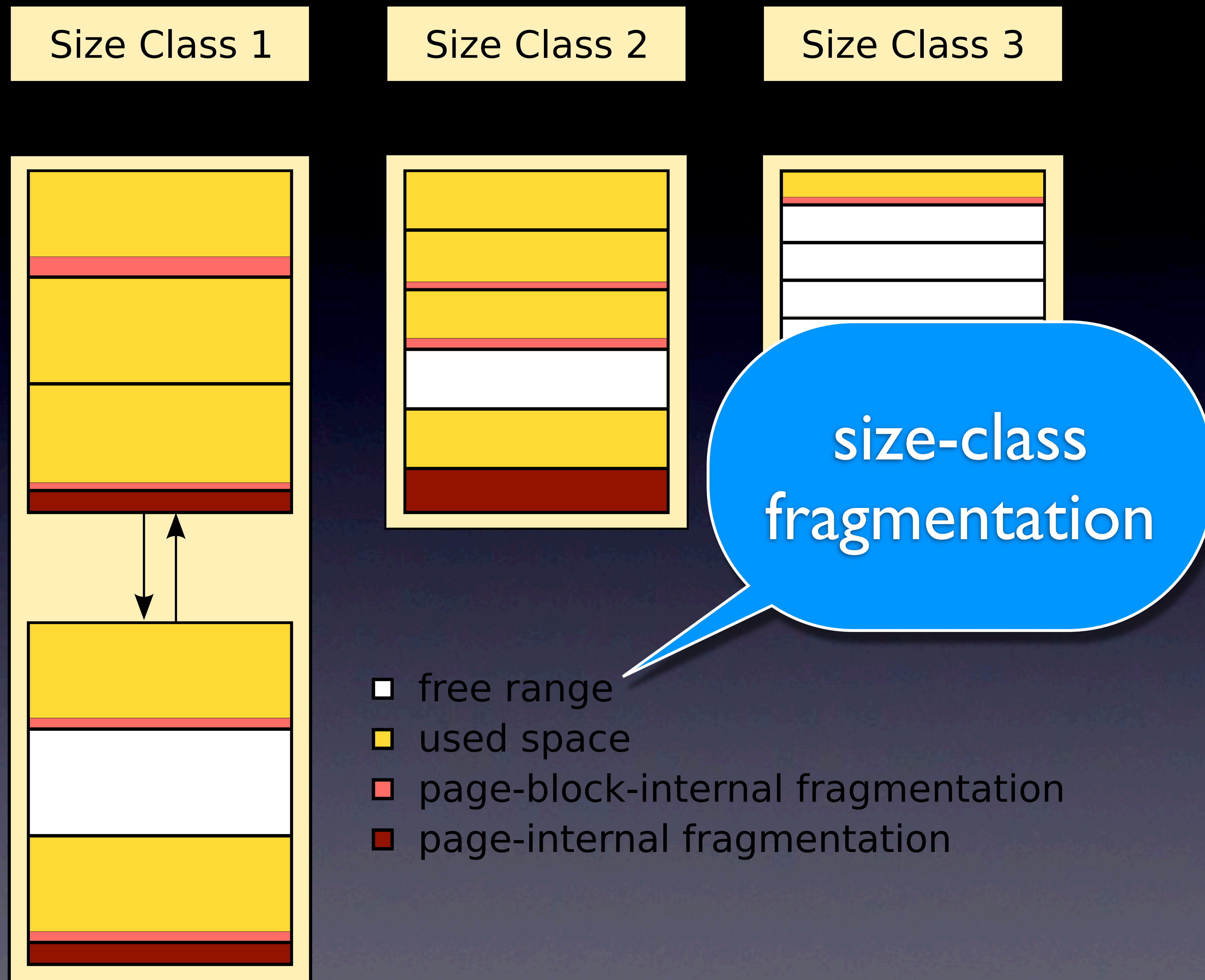
- Fragmentation through partitioning is fixed at compile time and is not controlled by partial compaction:
  - Page-block-internal fragmentation
  - Page-internal fragmentation

Fragmentation through Partitioning

- Fragmentation through partitioning is fixed at compile time and is not controlled by partial compaction:
  - Page-block-internal fragmentation
  - Page-internal fragmentation
- May dominate overall fragmentation



Size Class 1
Size Class 2
Size Class 3
free range
used space
page-block-internal fragmentation
page-internal fragmentation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Incremental Compaction

- Global compaction increment ι bounds size of memory involved in any atomic compaction operation:
  - $1 < \iota < \infty$: incremental compaction of objects larger than ι
  - $\iota = \infty$: non-incremental compaction

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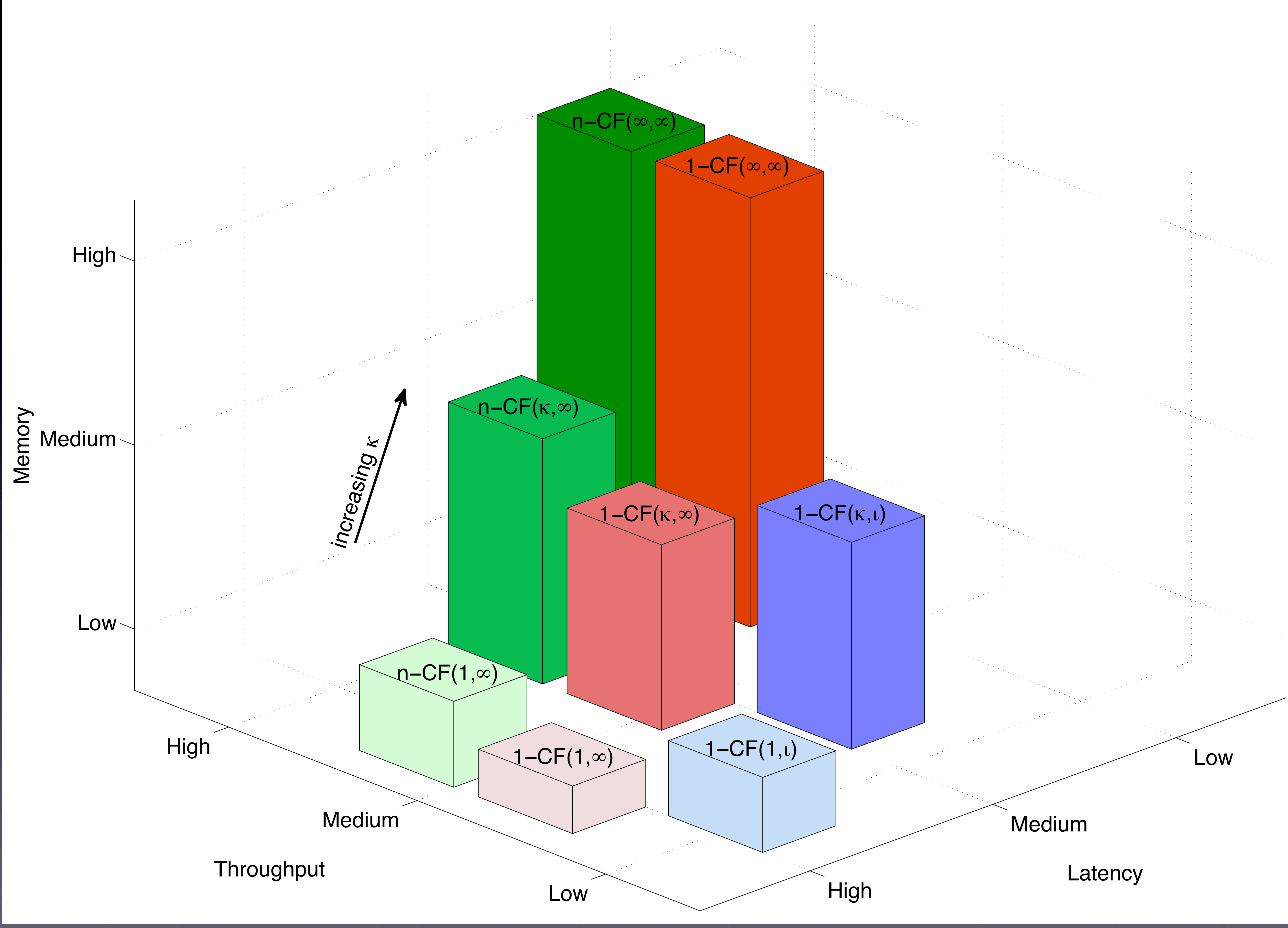
- Global compaction increment ι bounds size of memory involved in any atomic compaction operation:
  - $1 < \iota < \infty$: incremental compaction of objects larger than ι
  - $\iota = \infty$: non-incremental compaction
- Incremental compaction creates transient size-class fragmentation

CF Configurations

- 1-CF(κ, ι)
  - one CF instance for multiple threads
  - partial compaction bound κ
  - compaction increment ι

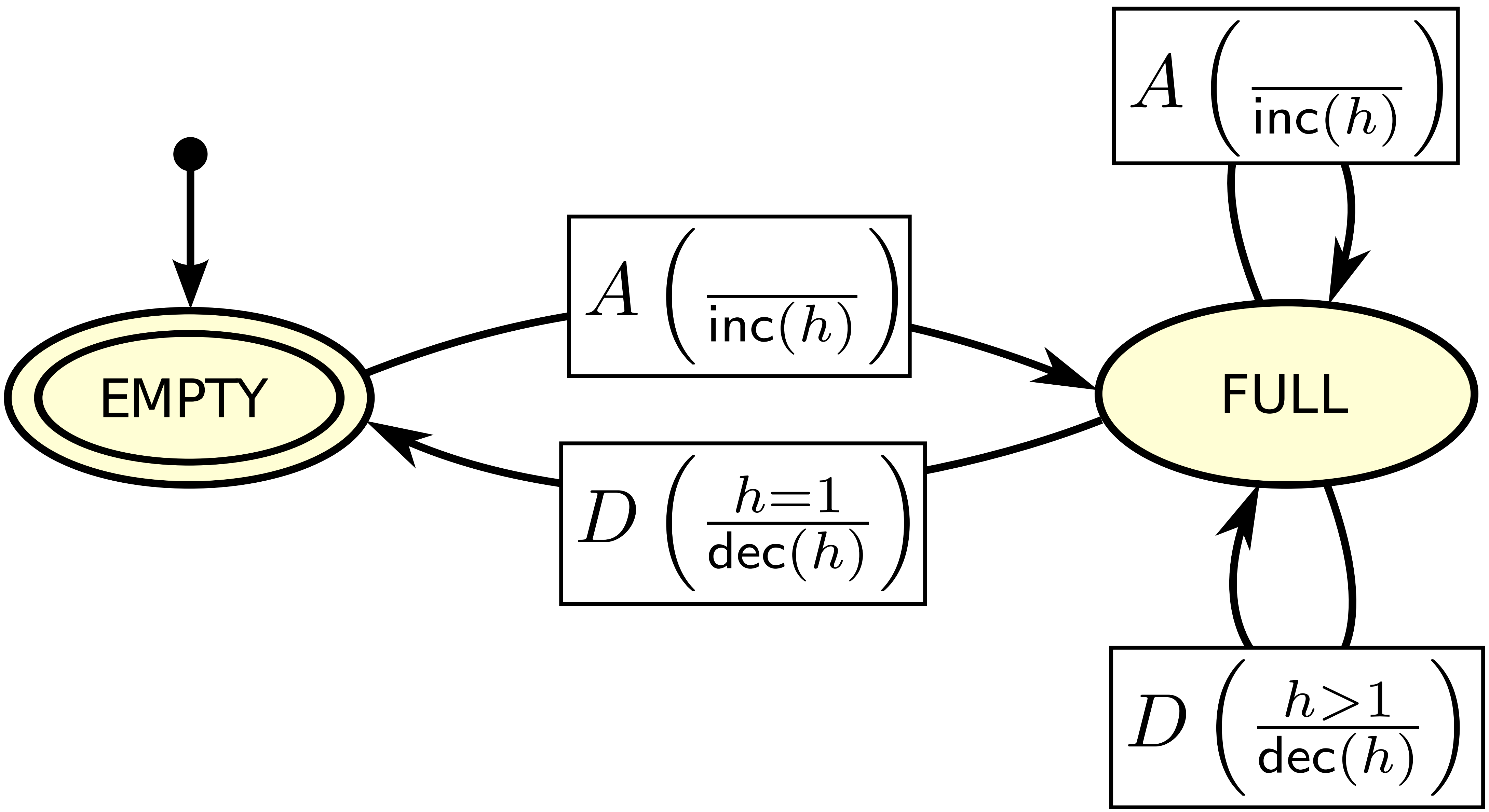
CF Configurations

- 1-CF(κ, ι)
  - one CF instance for multiple threads
  - partial compaction bound κ
  - compaction increment ι
- n-CF(κ, ι)
  - n CF instances for n threads
  - allows to control degree of sharing

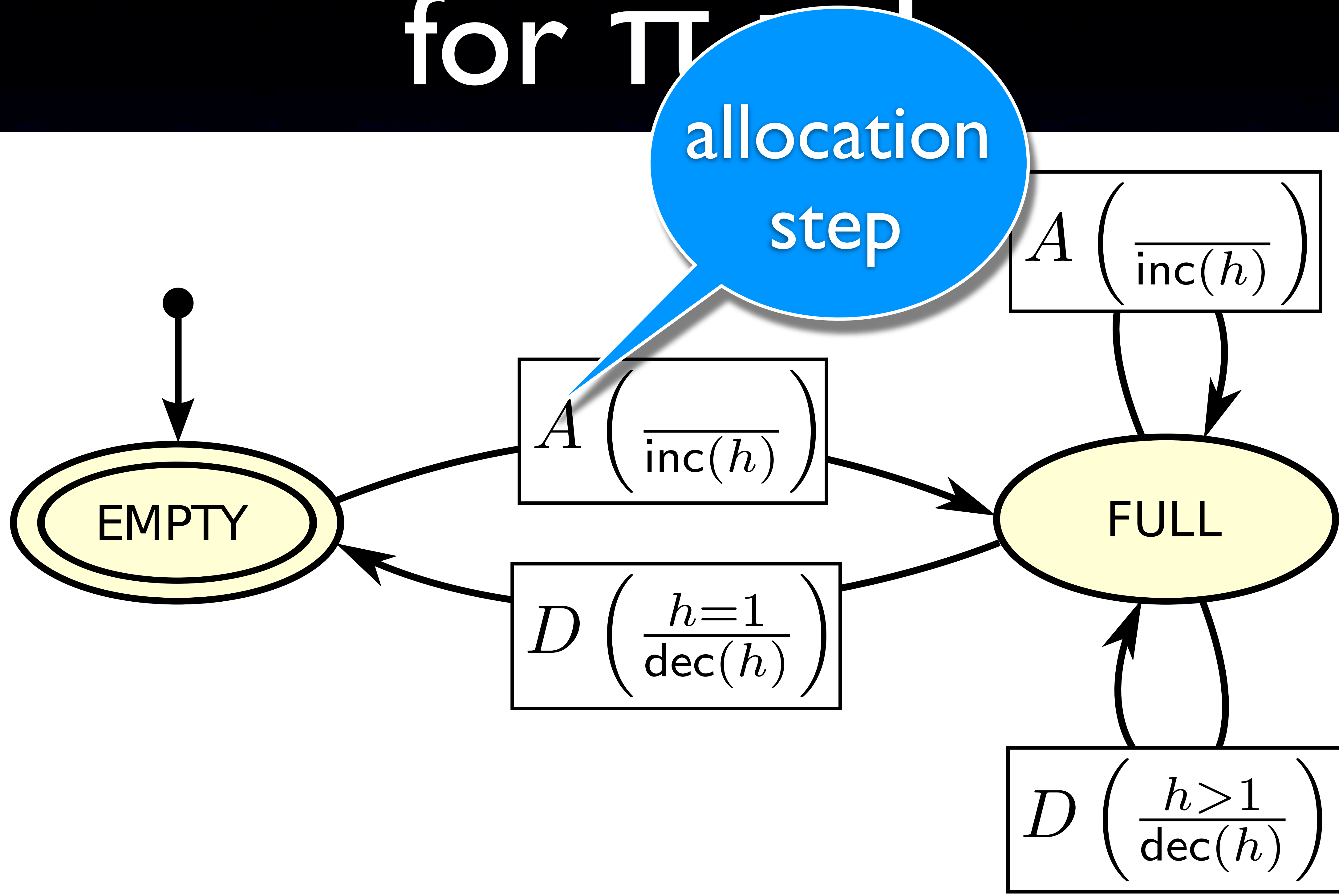

n–CF(∞,∞)
1–CF(∞,∞)
n–CF(κ,∞)
1–CF(κ,∞)
1–CF(κ,ι)
n–CF(1,∞)
1–CF(1,∞)
1–CF(1,ι)
increasing κ
Memory
High
Medium
Low
Throughput
High
Medium
Low
Latency
Low
Medium
High

To make CF
concurrent and incremental
we model the algorithm
as a
finite state machine
whose transitions
must be atomic!

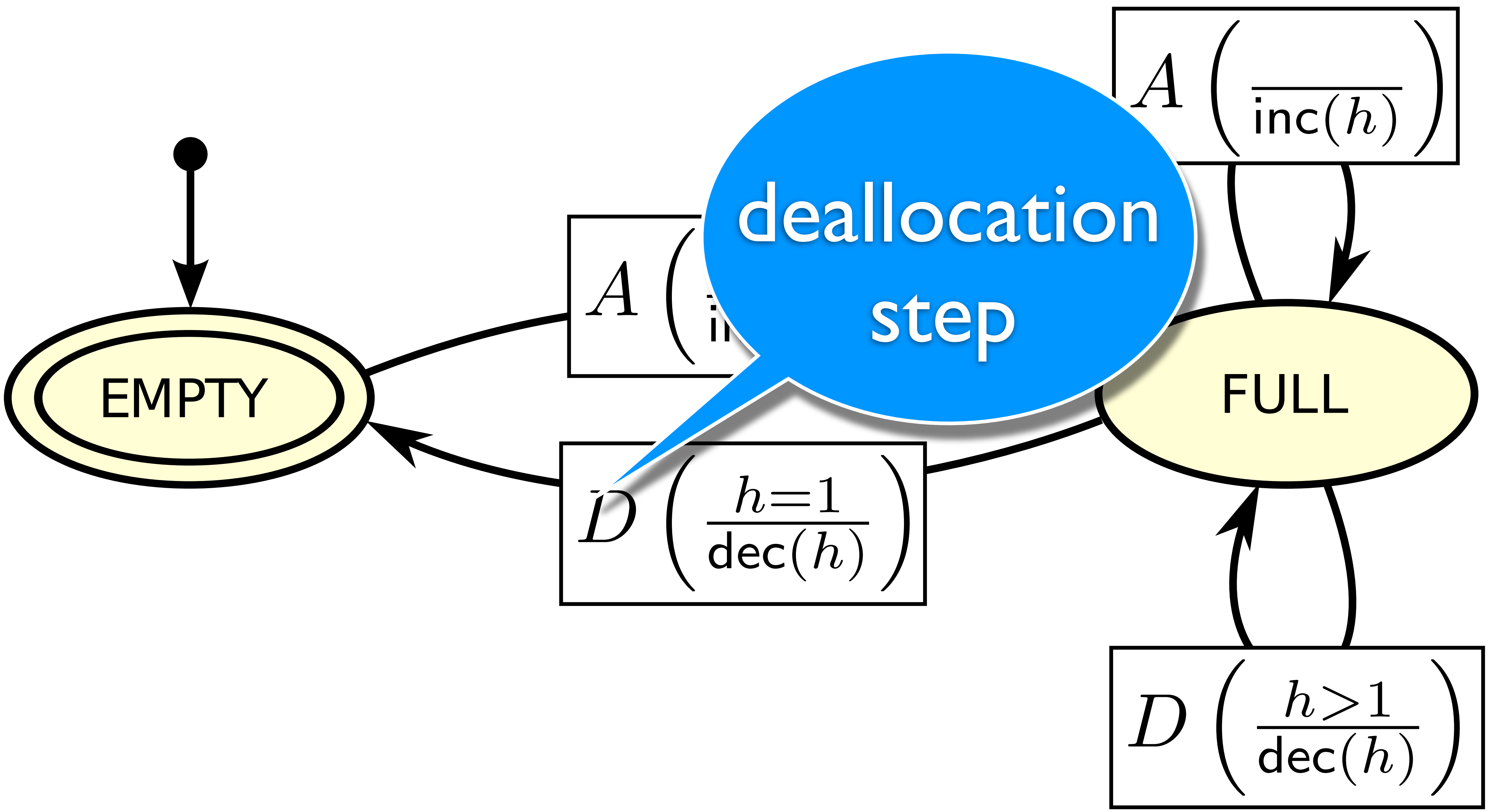
Size-Class Automaton for π = 1



Size-Class Automaton for Π



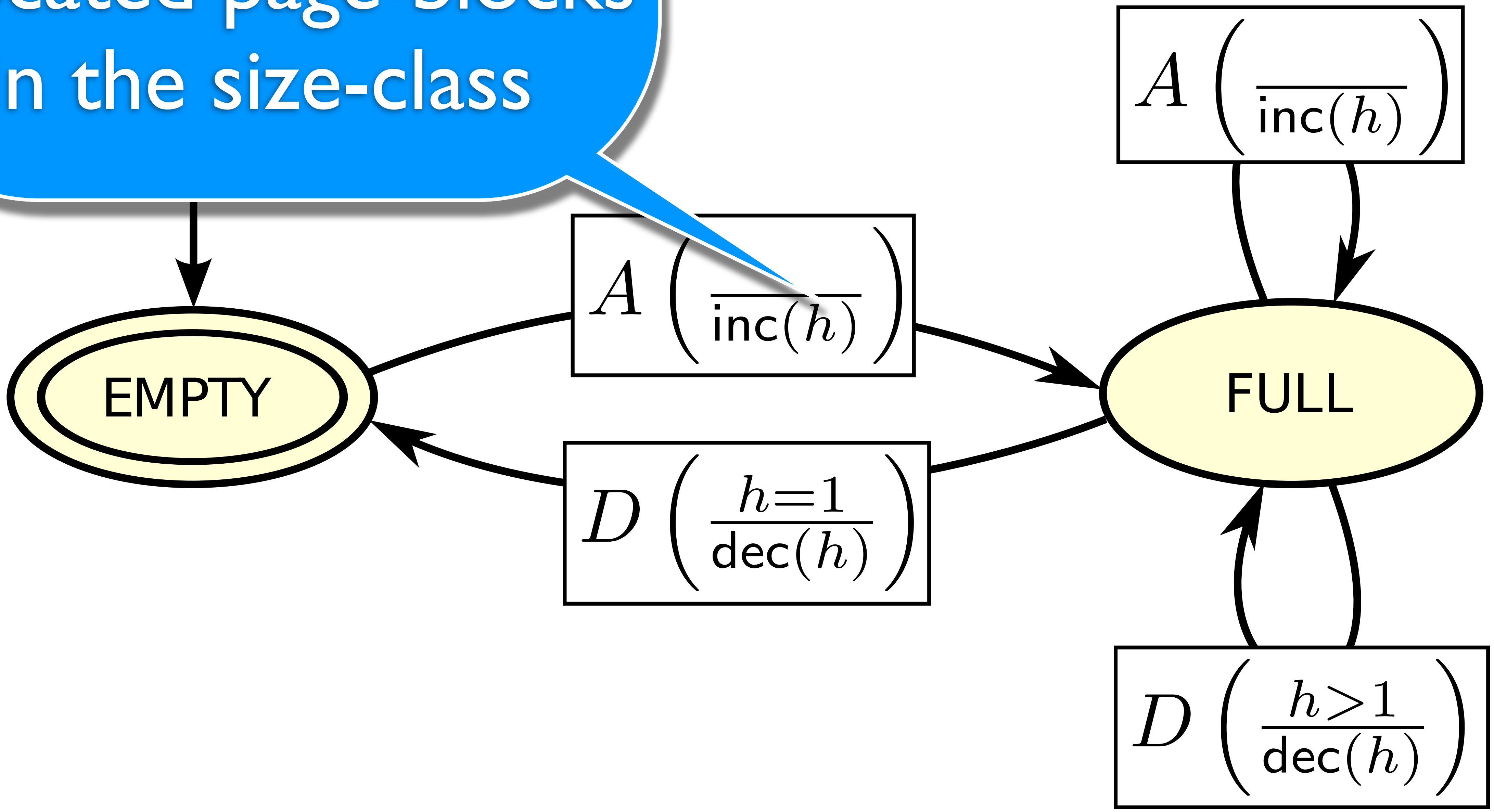
Size-Class Automaton for π = 1



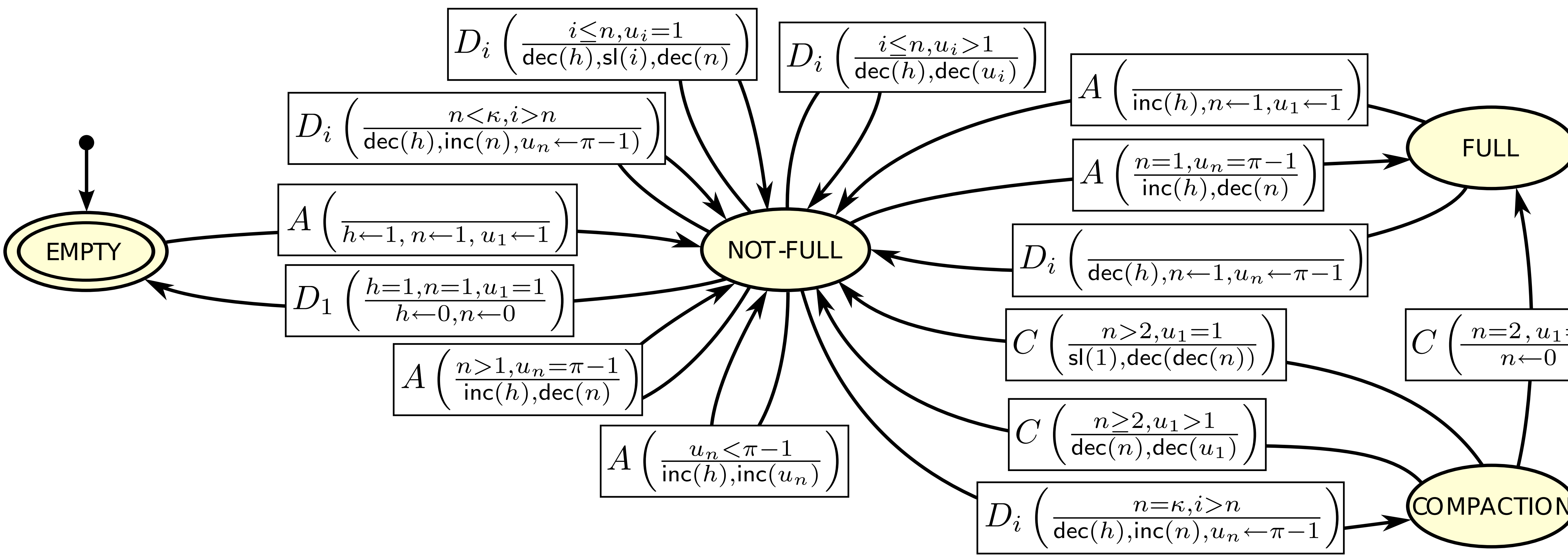
Size-Class Automaton

$\pi = 1$

h is the total # of allocated page-blocks in the size-class



Size-Class Automaton for π > 1

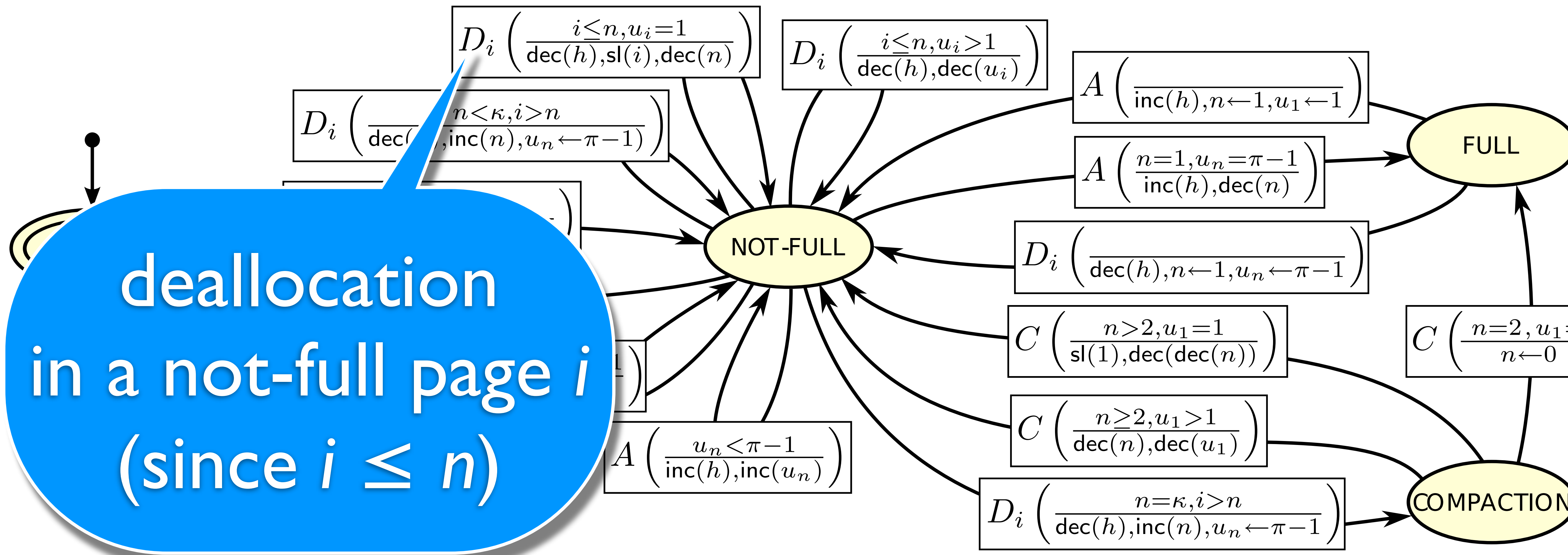


h is the total # of allocated page-blocks in the size-class

n is the # of not-full pages

u_i is the # of used page-blocks in a not-full page *i*

Size-Class Automaton for π > 1



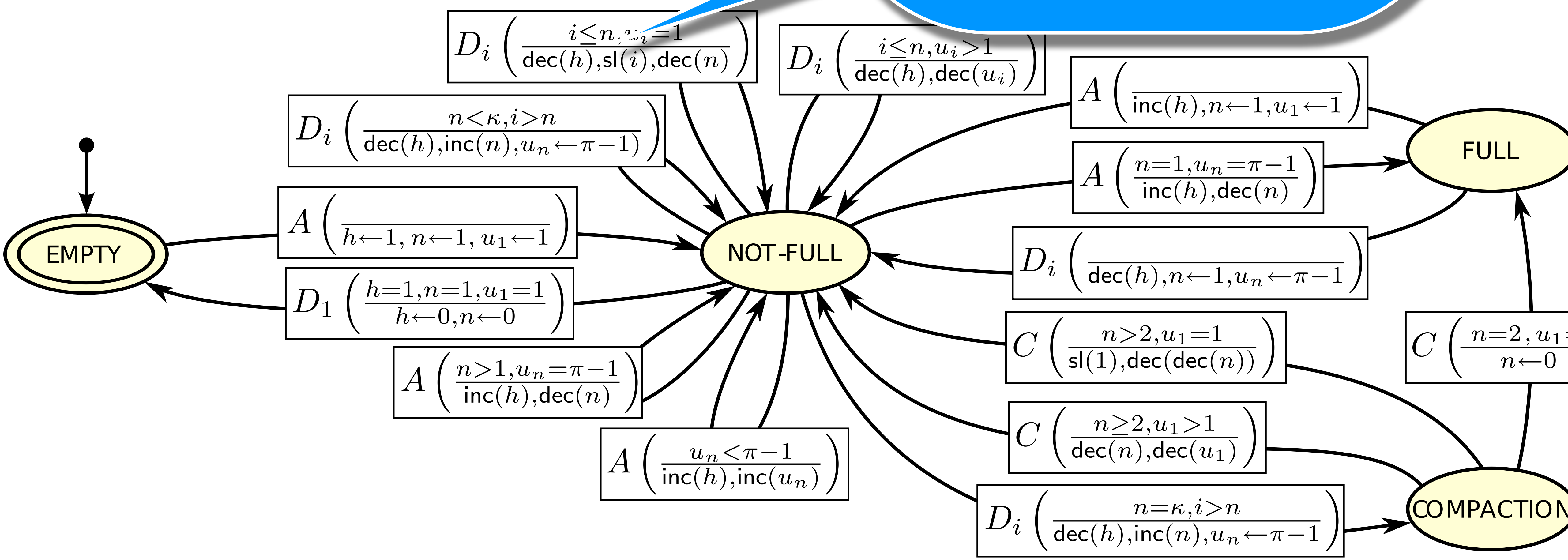
h is the total # of allocated page-blocks in the size-class

n is the # of not-full pages

u_i is the # of used page-blocks in a not-full page i

Size-Class A
for T

remove page since it is now empty



h is the total # of allocated page-blocks in the size-class
n is the # of not-full pages
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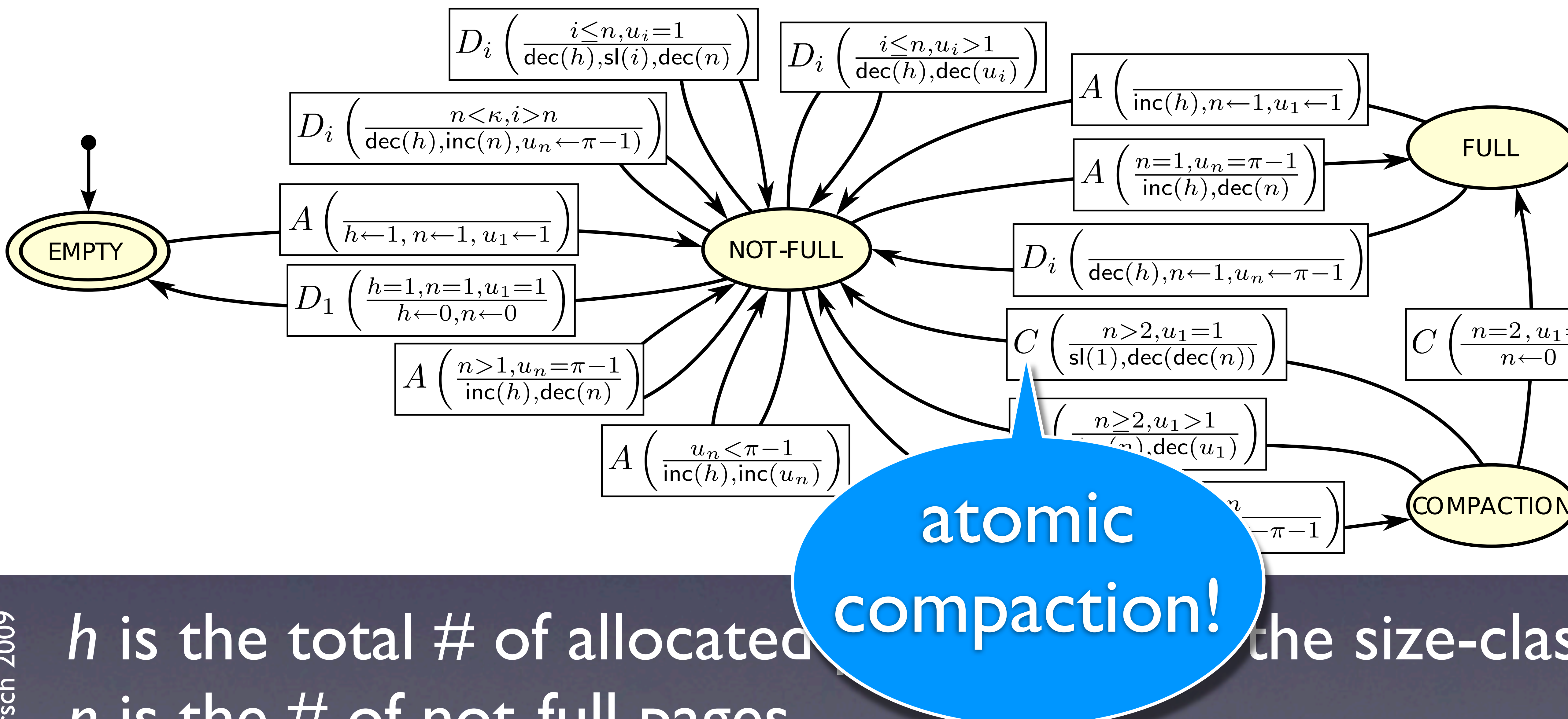
Size-Class Automaton for π > 1



h is the total # of allocated … the size-class
n is the # of not-full pages
u_i is the # of used page-blocks in a not-full page *i*

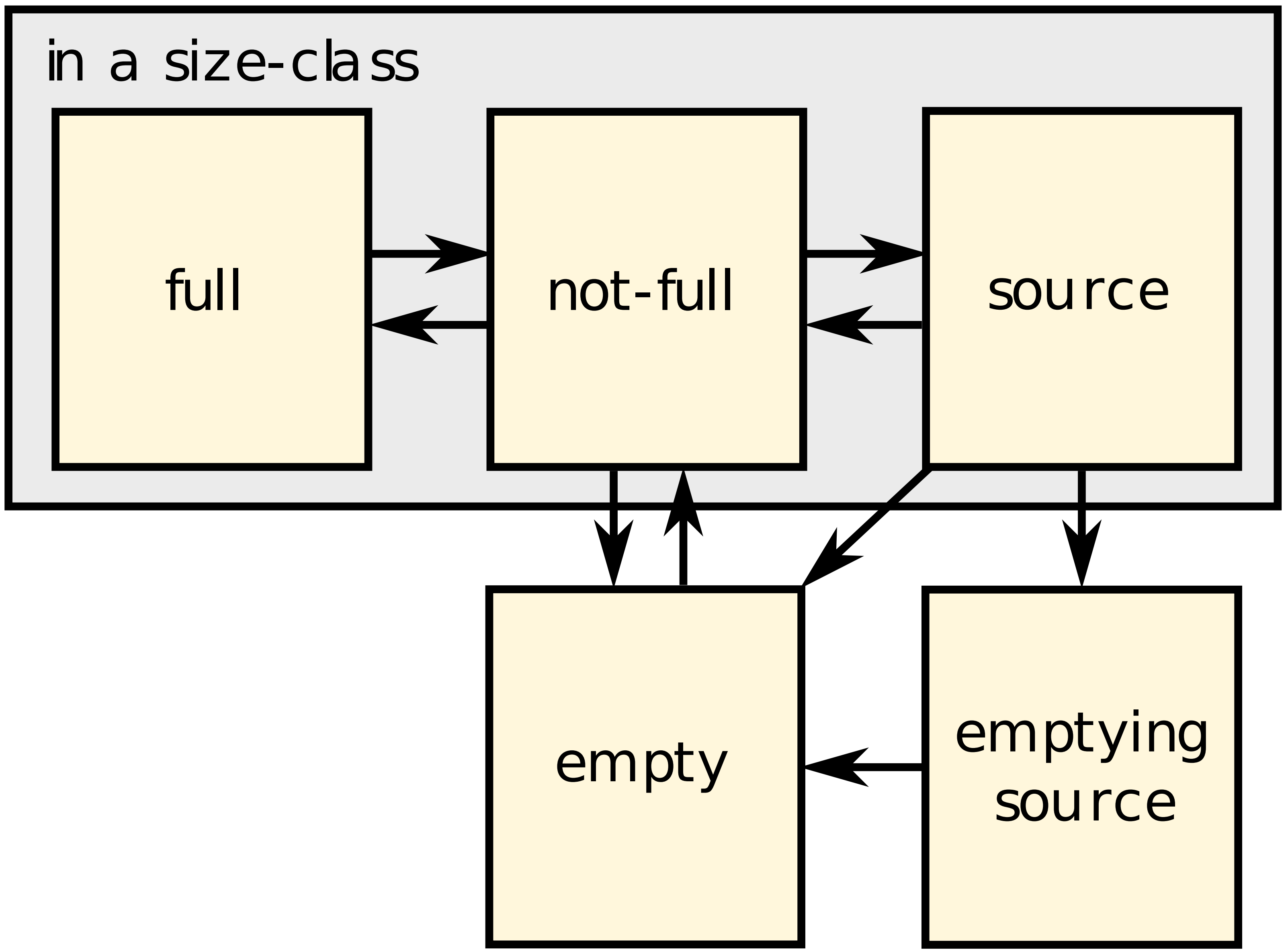
Incremental Compaction

- A page-block that is <u>incrementally</u> moved actually occupies two page-blocks:
  - source page-block
  - target page-block

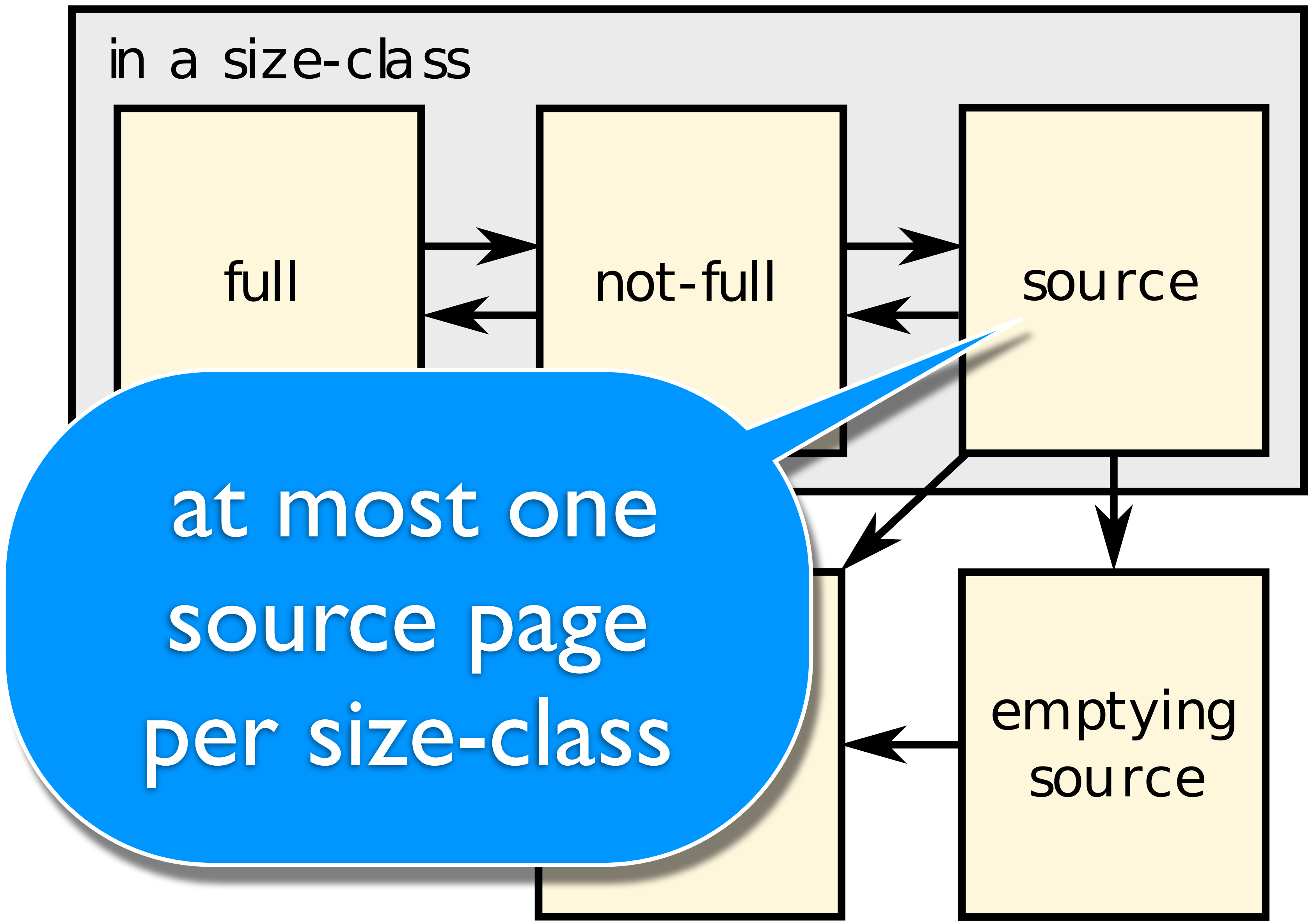
Incremental Compaction

- A page-block that is <u>incrementally</u> moved actually occupies two page-blocks:
  - source page-block
  - target page-block
- A page containing source page-blocks is called source page
  - may also contain used and free page-blocks

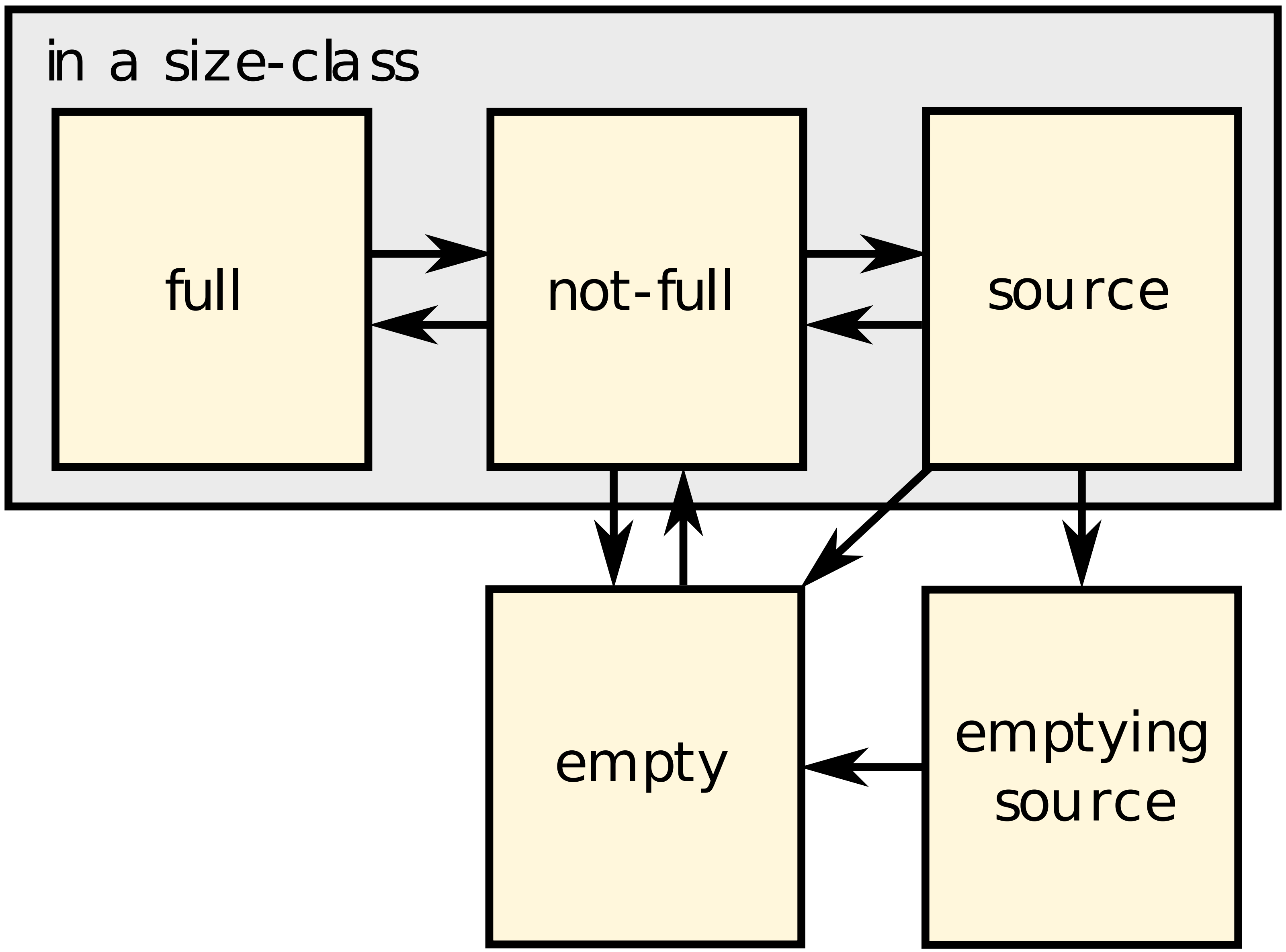
The Lifetime of a Page



The Lifetime of a Page

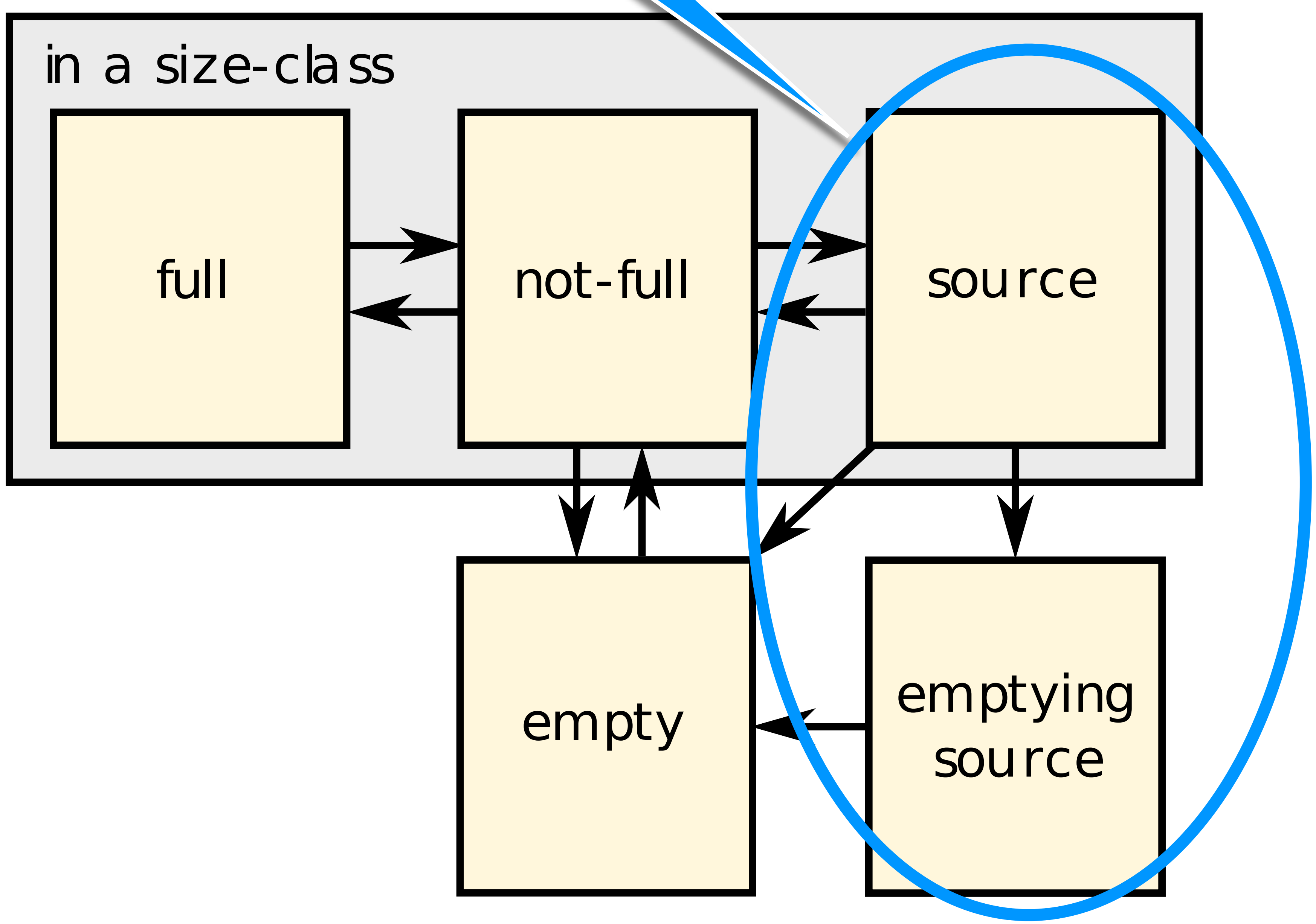


The Lifetime of a Page

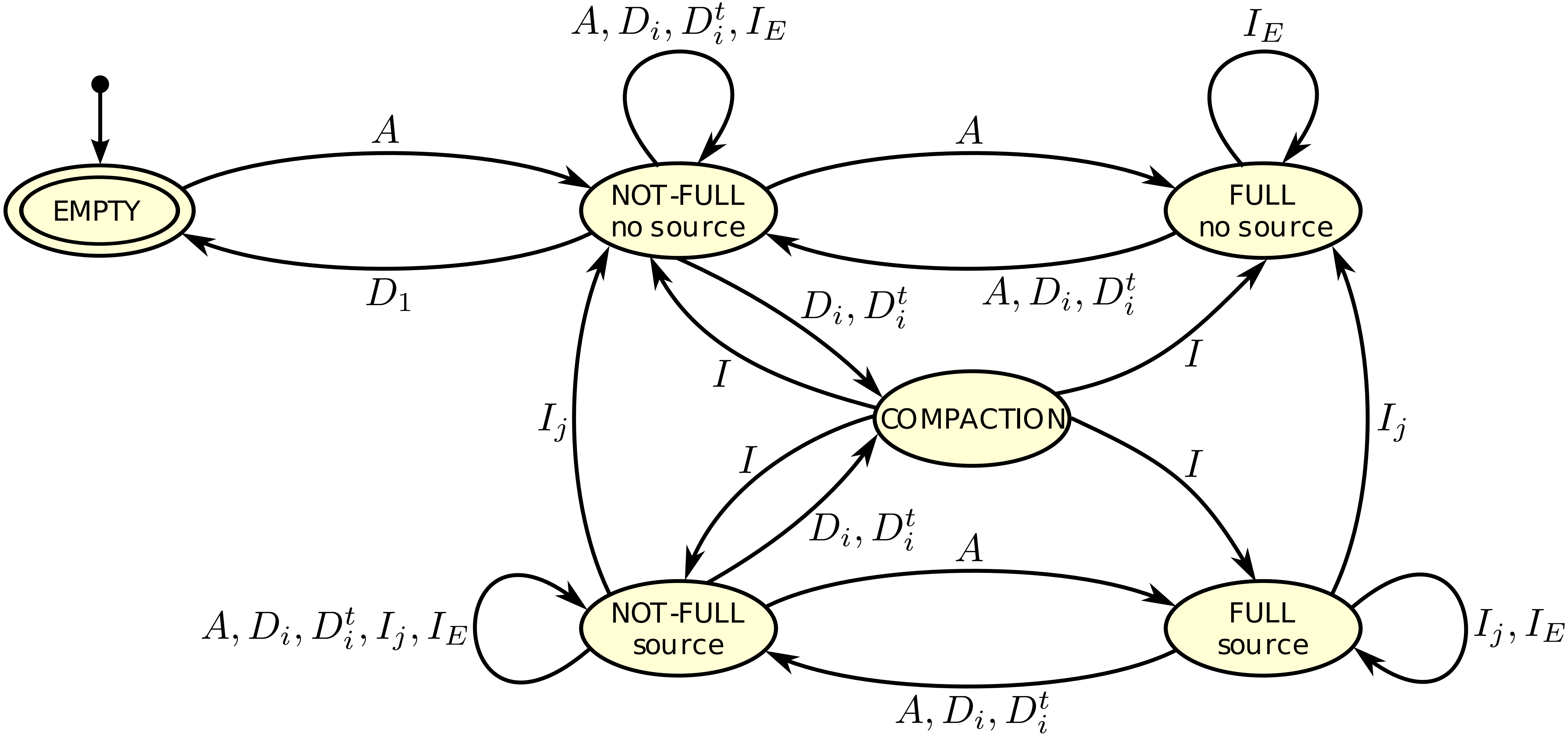


transient size-class fragmentation

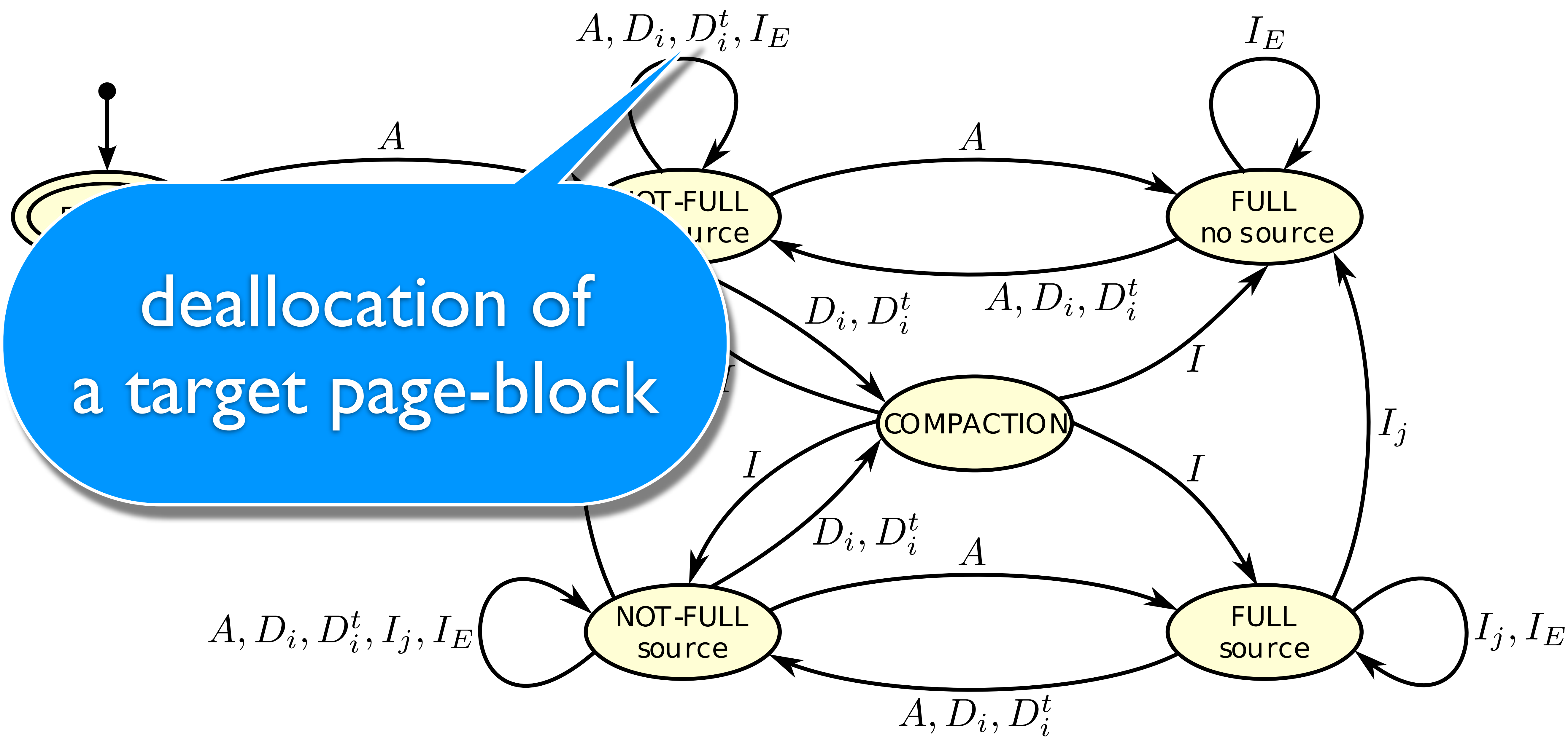
# of a Page



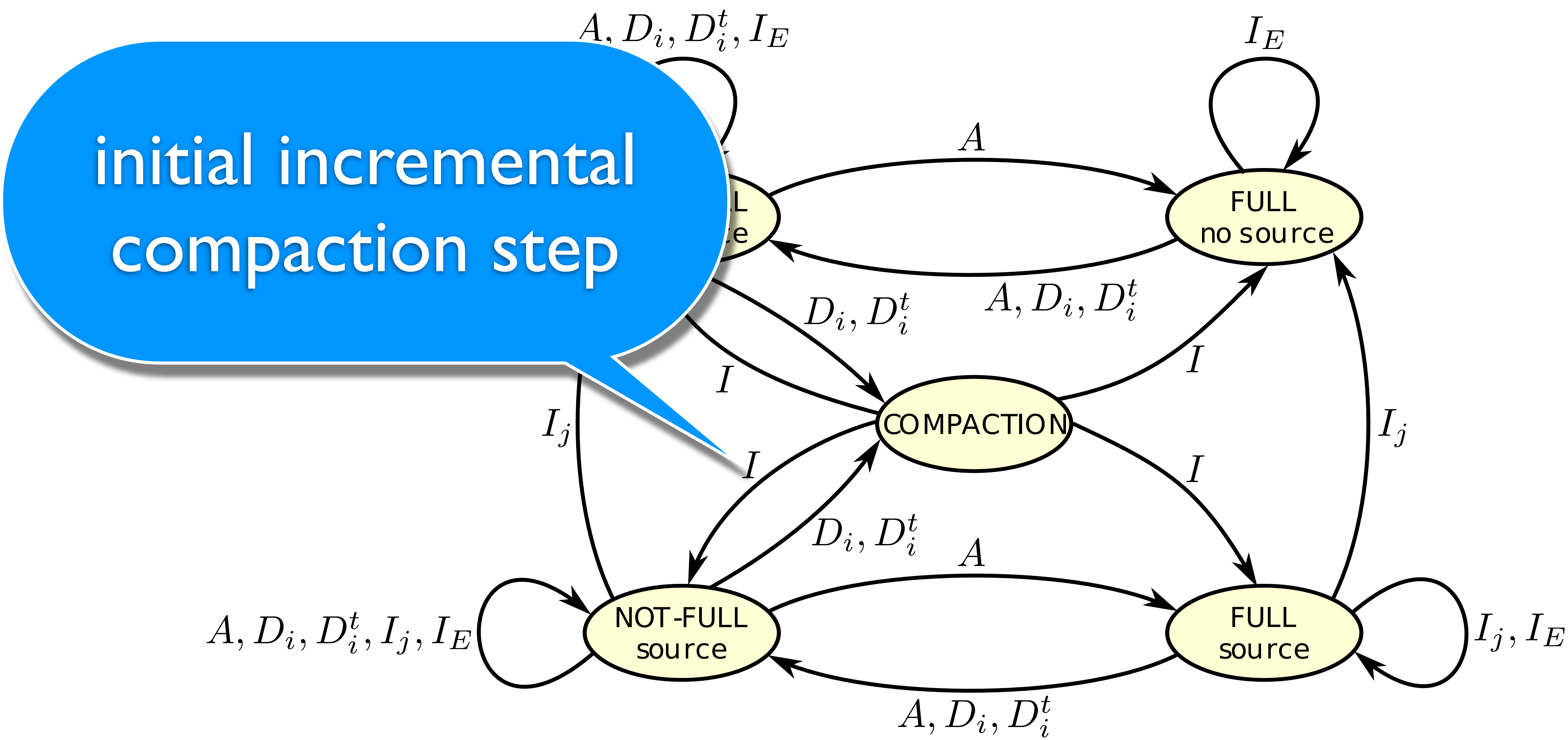
Incremental Size-Class Automaton for π > 1



Incremental Size-Class Automaton for π > 1

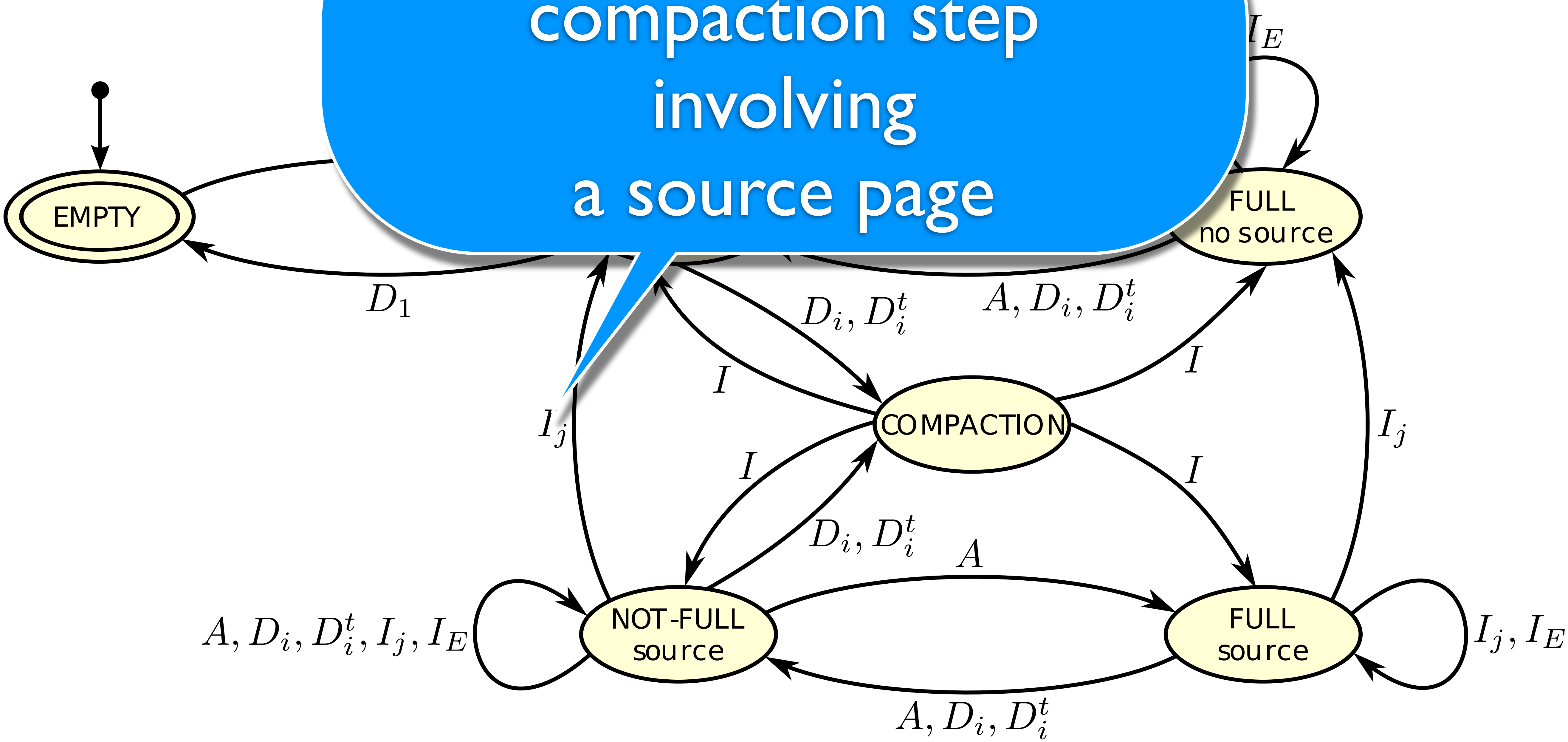


Incremental Size-Class Automaton for π > 1

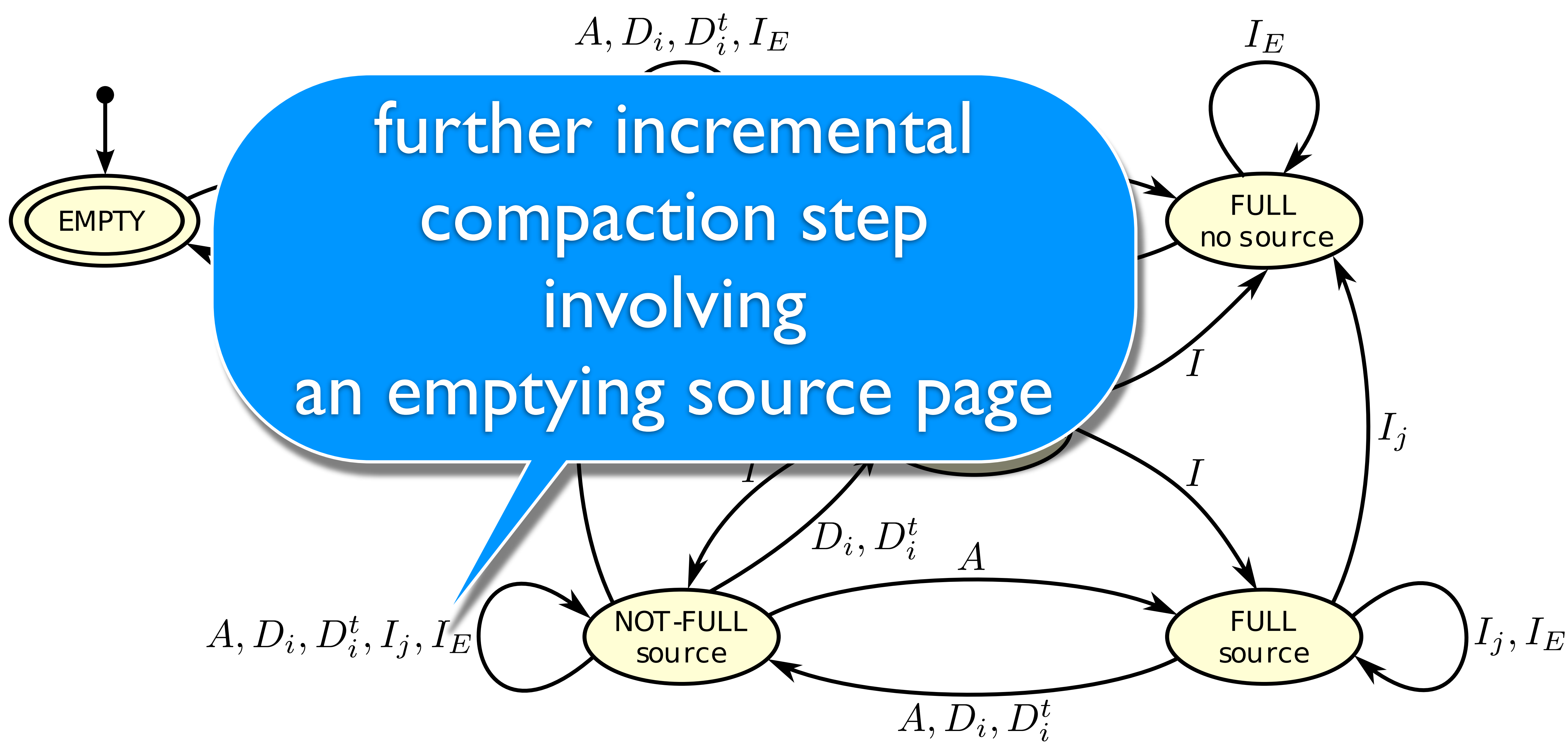


Incremental Size-Class Automaton for $\Pi > 1$

further incremental compaction step involving a source page



Incremental Size-Class Automaton for π > 1



Temporal and Spatial Complexity per CF Configuration and Size-Class

	malloc	free	latency
$1\text{-CF}(\infty,\infty)$	$O(n)$	$O(n)$	$O(1)$
$1\text{-CF}(\kappa,\infty)$	$O(n)$	$O(n+\beta)$	$O(\beta)$
$n\text{-CF}(\infty,\infty)$	$O(1)$	$O(1)$	$O(1)$
$n\text{-CF}(\kappa,\infty)$	$O(1)$	$O(\beta)$	$O(\beta)$
$1\text{-CF}(\kappa,\iota)$	$O(n)$	$O(n+\beta+\lfloor\frac{\beta}{\iota}\rfloor)$	$O(\min(\beta,\iota))$

	memory size	size-class fragmentation
$1\text{-CF}(\infty,\infty)$	$O(n*m*\pi*\beta)$	$O(n*m*(\pi-1)*\beta)$
$1\text{-CF}(\kappa,\infty)$	$O((n*m+\kappa*(\pi-1))*\beta)$	$O(\kappa*(\pi-1)*\beta)$
$n\text{-CF}(\infty,\infty)$	$O(n*m*\pi*\beta)$	$O(n*m*(\pi-1)*\beta)$
$n\text{-CF}(\kappa,\infty)$	$O(n*(m+\kappa*(\pi-1))*\beta)$	$O(n*\kappa*(\pi-1)*\beta)$
$1\text{-CF}(\kappa,\iota)$	$O((n*m+n*\pi+\kappa*(\pi-1))*\beta)$	$O((n*\pi+\kappa*(\pi-1))*\beta)$

Temporal and Spatial Complexity per CF Configuration and Size-Class

	malloc	free	latency
1-CF(∞, ∞)	$O(n)$	$O(n)$	$O(1)$
1-CF(κ, ∞)	$O(n)$	$O(n + \beta)$	$O(\beta)$
n-CF(∞, ∞)	$O(1)$	$O(1)$	$O(1)$
[illegible]	$O(1)$	$O(\beta)$	$O(\beta)$
[illegible]	[illegible]	[illegible]$(n + \beta + \lfloor \frac{\beta}{\iota} \rfloor)$	$O(\min(\beta, \iota))$

n is the # of threads

	[illegible] size	size-class fragmentation
1-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
1-CF(κ, ∞)	$O((n * m + \kappa * (\pi - 1)) * \beta)$	$O(\kappa * (\pi - 1) * \beta)$
n-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
n-CF(κ, ∞)	$O(n * (m + \kappa * (\pi - 1)) * \beta)$	$O(n * \kappa * (\pi - 1) * \beta)$
1-CF(κ, ι)	$O((n * m + n * \pi + \kappa * (\pi - 1)) * \beta)$	$O((n * \pi + \kappa * (\pi - 1)) * \beta)$

Temporal and Spatial Complexity per CF Configuration and Size-Class

	malloc	free	latency
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n-CF(∞, ∞)	$O(1)$	$O(1)$	$O(1)$
n-CF(κ, ∞)	$O(1)$	$O(\beta)$	$O(\beta)$
1-CF(κ, ι)	$O(n)$	$O(n + \beta + \lfloor \frac{\beta}{\iota} \rfloor)$	$O(\min(\beta, \iota))$

	memory size	size-class fragmentation
1-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
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	malloc	free	latency
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1-CF(κ, ∞)	$O(n)$	$O(n+\beta)$	$O(\beta)$
n-CF(∞, ∞)	$O(1)$	$O(1)$	$O(1)$
n-CF(κ, ∞)	$O(1)$	$O(\beta)$	[illegible]
1-CF(κ, ι)	$O(n)$	$O($ [illegible]	[illegible]

β is the page-block size

	memory size	[illegible]
1-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
1-CF(κ, ∞)	$O((n * m + \kappa * (\pi - 1)) * \beta)$	$O(\kappa * (\pi - 1) * \beta)$
n-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
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Temporal and Spatial Complexity per CF Configuration and Size-Class

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n-CF(∞, ∞)	$O(1)$	$O(1)$	$O(1)$
n-CF(κ, ∞)	$O(1)$	$O(\beta)$	$O(\beta)$
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	memory size	size-class fragmentation
1-CF(∞, ∞)	$O(n*m*\pi*\beta)$	$O(n*m*(\pi-1)*\beta)$
1-CF(κ, ∞)	$O((n*m+\kappa*(\pi-1))*\beta)$	$O(\kappa*(\pi-1)*\beta)$
n-CF(∞, ∞)	$O(n*m*\pi*\beta)$	$O(n*m*(\pi-1)*\beta)$
n-CF(κ, ∞)	$O(n*(m+\kappa*(\pi-1))*\beta)$	$O(n*\kappa*(\pi-1)*\beta)$
1-CF(κ, ι)	$O((n*m+n*\pi+\kappa*(\pi-1))*\beta)$	$O((n*\pi+\kappa*(\pi-1))*\beta)$

Temporal and Spatial Complexity per CF Configuration and Size-Class

		latency
1-CF		$O(1)$
1-C		$O(\beta)$
n-C		$O(1)$
n-CF		$O(\beta)$
1-CF(κ, ι)		$O(\min(\beta, \iota))$

m is the # of per-thread-allocated page-blocks

	memory size	size-class fragmentation
1-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
1-CF(κ, ∞)	$O((n * m + \kappa * (\pi - 1)) * \beta)$	$O(\kappa * (\pi - 1) * \beta)$
n-CF(∞, ∞)	$O(n * m * \pi * \beta)$	$O(n * m * (\pi - 1) * \beta)$
n-CF(κ, ∞)	$O(n * (m + \kappa * (\pi - 1)) * \beta)$	$O(n * \kappa * (\pi - 1) * \beta)$
1-CF(κ, ι)	$O((n * m + n * \pi + \kappa * (\pi - 1)) * \beta)$	$O((n * \pi + \kappa * (\pi - 1)) * \beta)$

Temporal and Spatial Complexity per CF Configuration and Size-Class

	malloc	free	latency
1-CF(∞, ∞)	$O(n)$	$O(n)$	$O(1)$
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n-CF(∞, ∞)	$O(1)$	$O(1)$	$O(1)$
n-CF(κ, ∞)	$O(1)$	$O(\beta)$	$O(\beta)$
1-CF(κ, ι)	$O(n)$	$O(n + \beta + \lfloor \frac{\beta}{\iota} \rfloor)$	$O(\min(\beta, \iota))$

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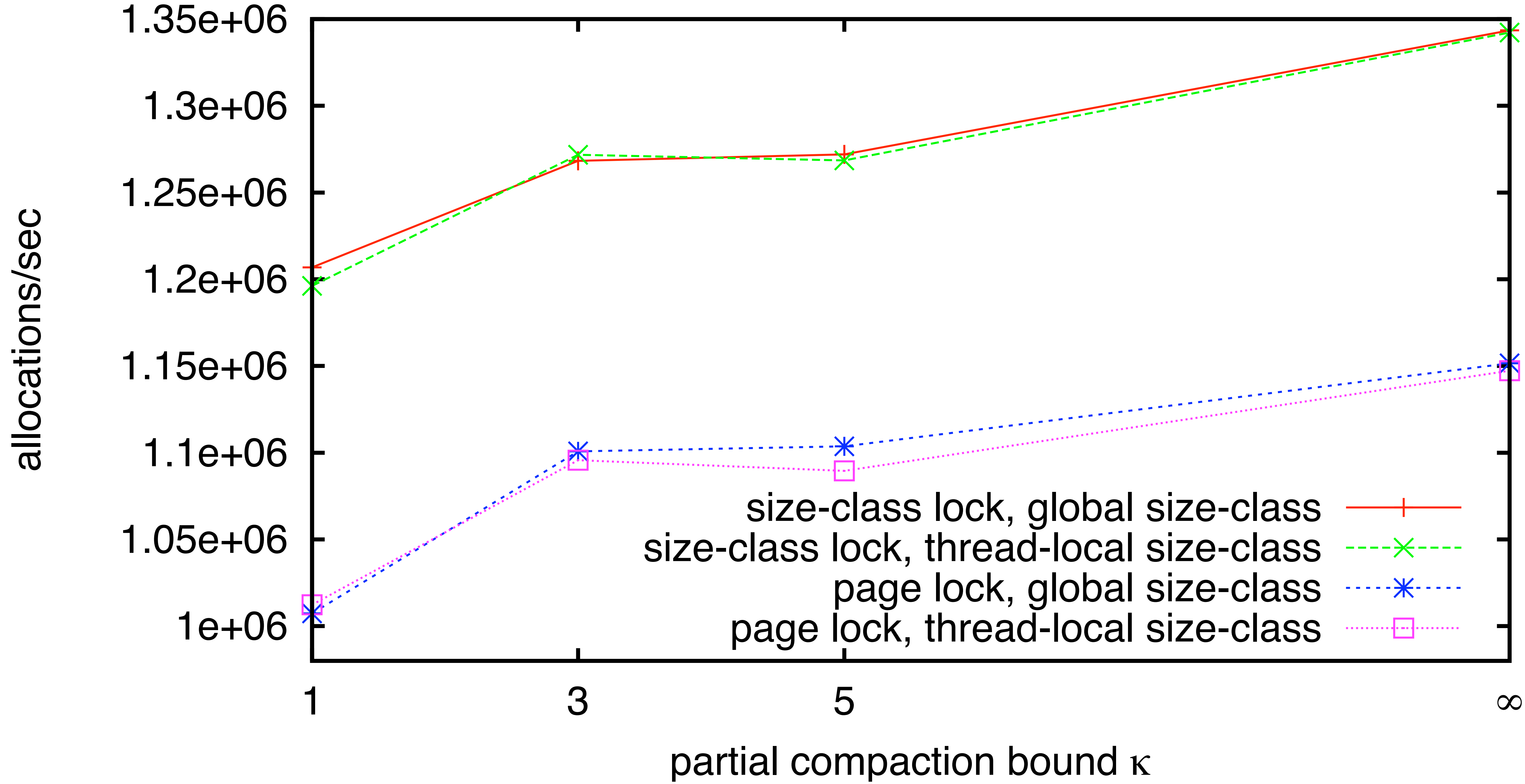
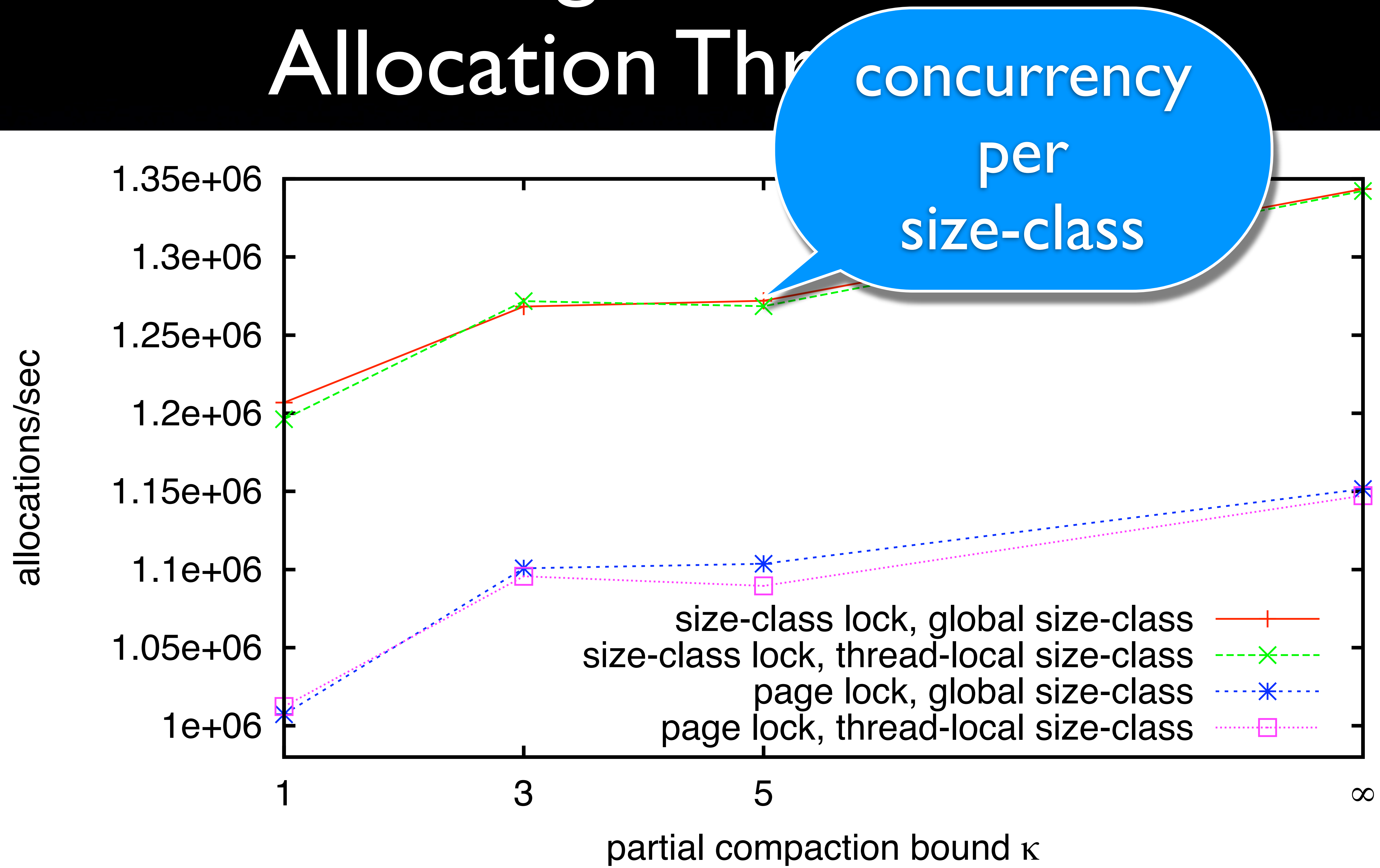
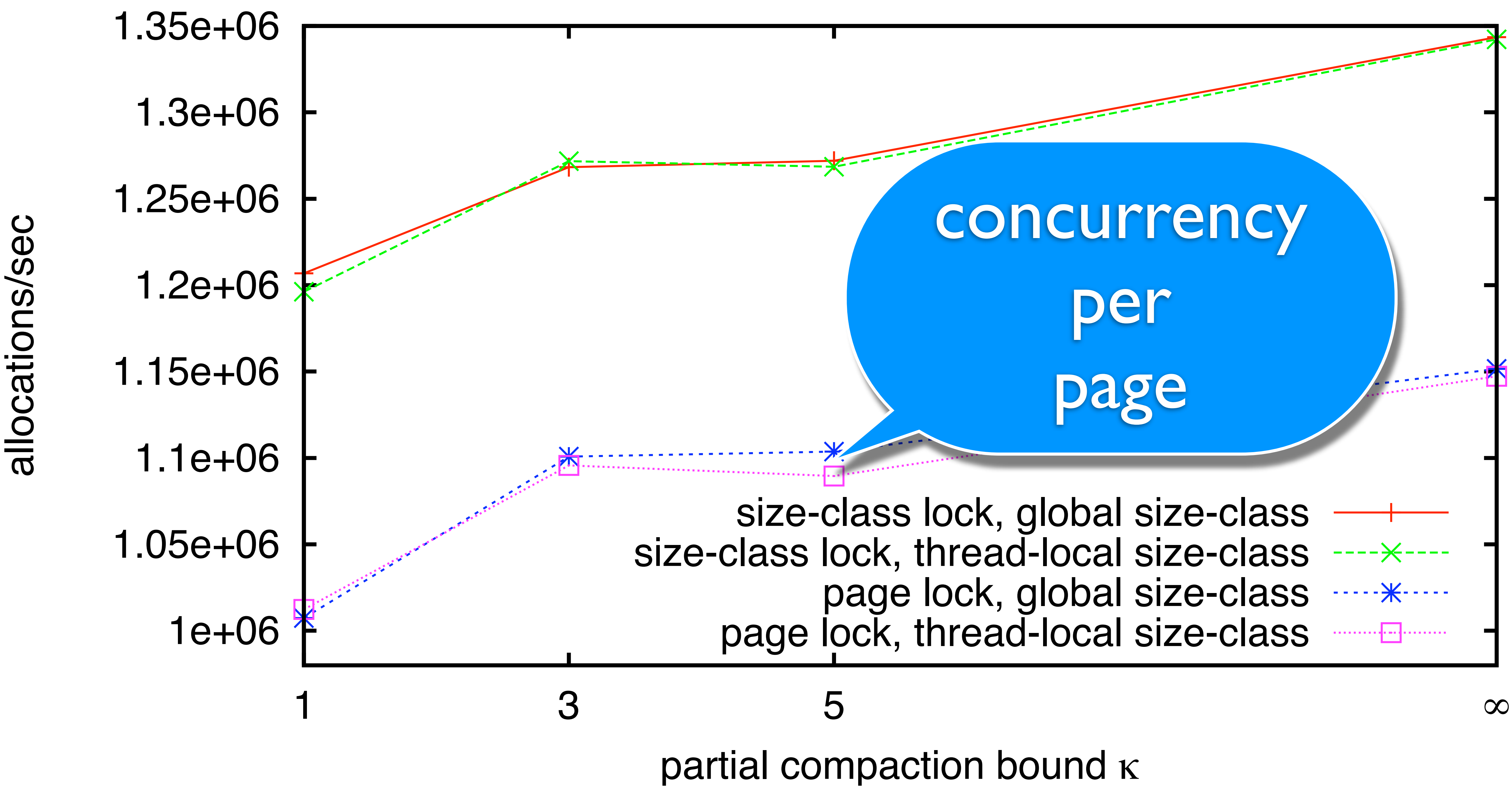
Single Thread Allocation Throughput



Single Thread Allocation Thr

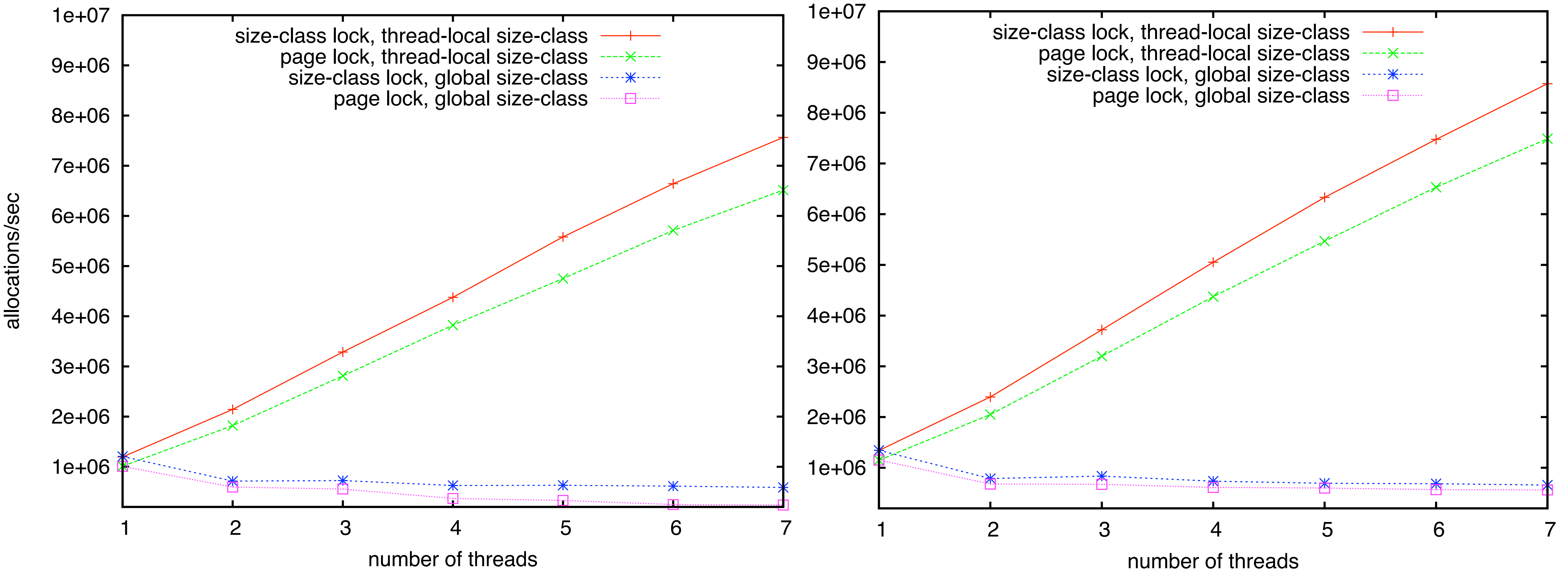


Single Thread Allocation Throughput



- less compaction may result in better allocation throughput
- size-class locks better than page locks

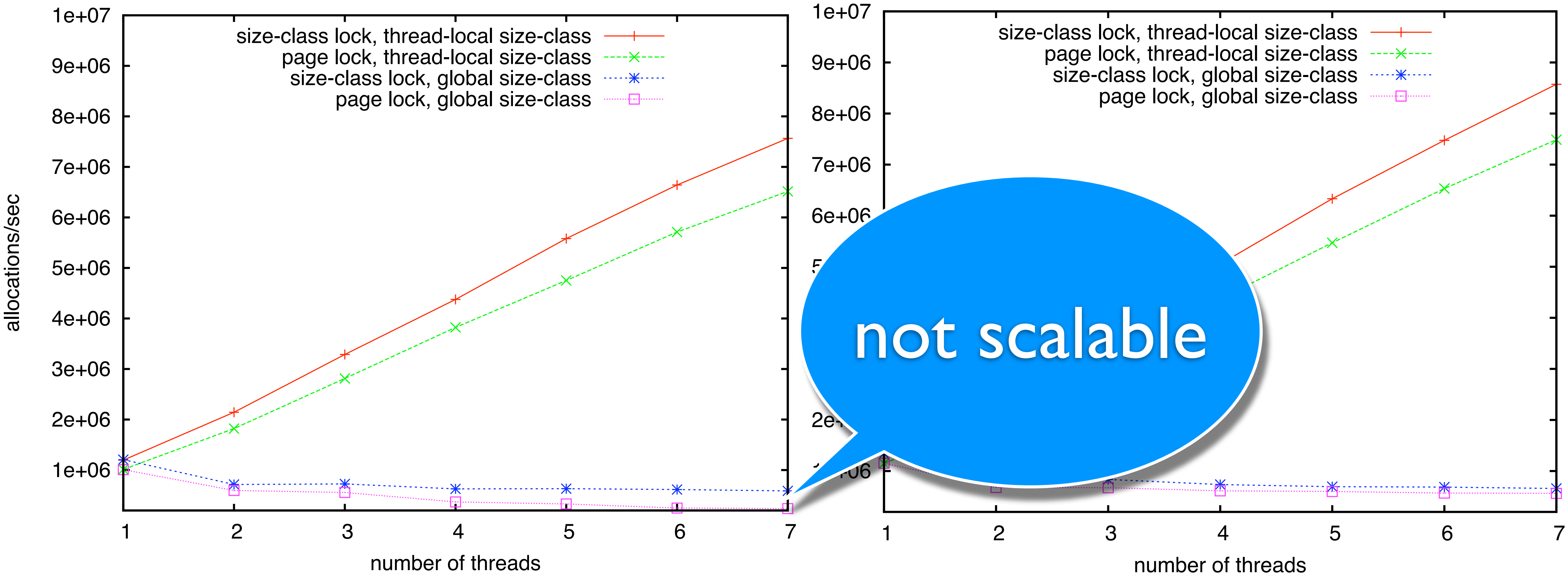
Scalability of Allocation Throughput



(a) full compaction

(b) optimized, non-compacting

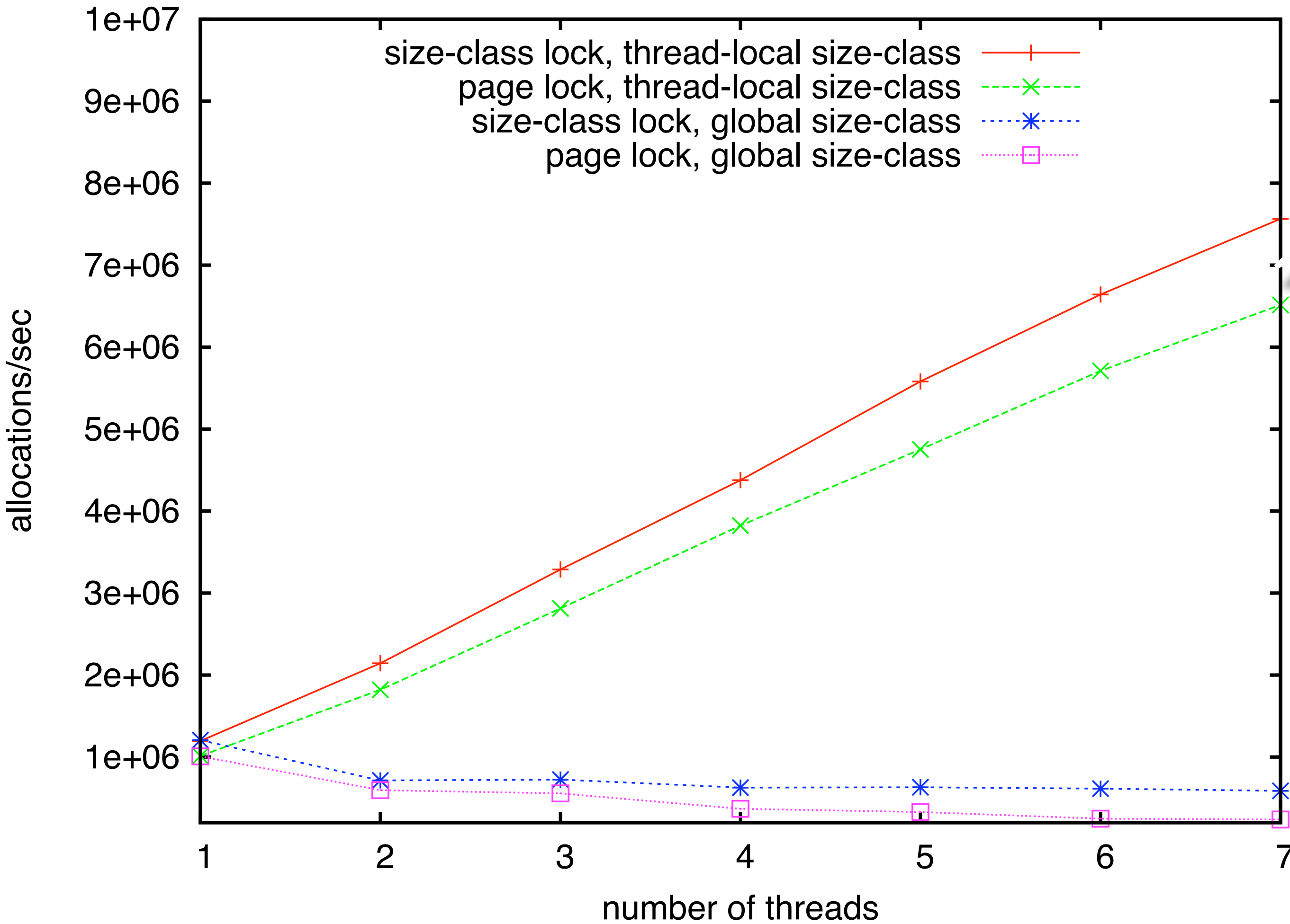
Scalability of Allocation Throughput



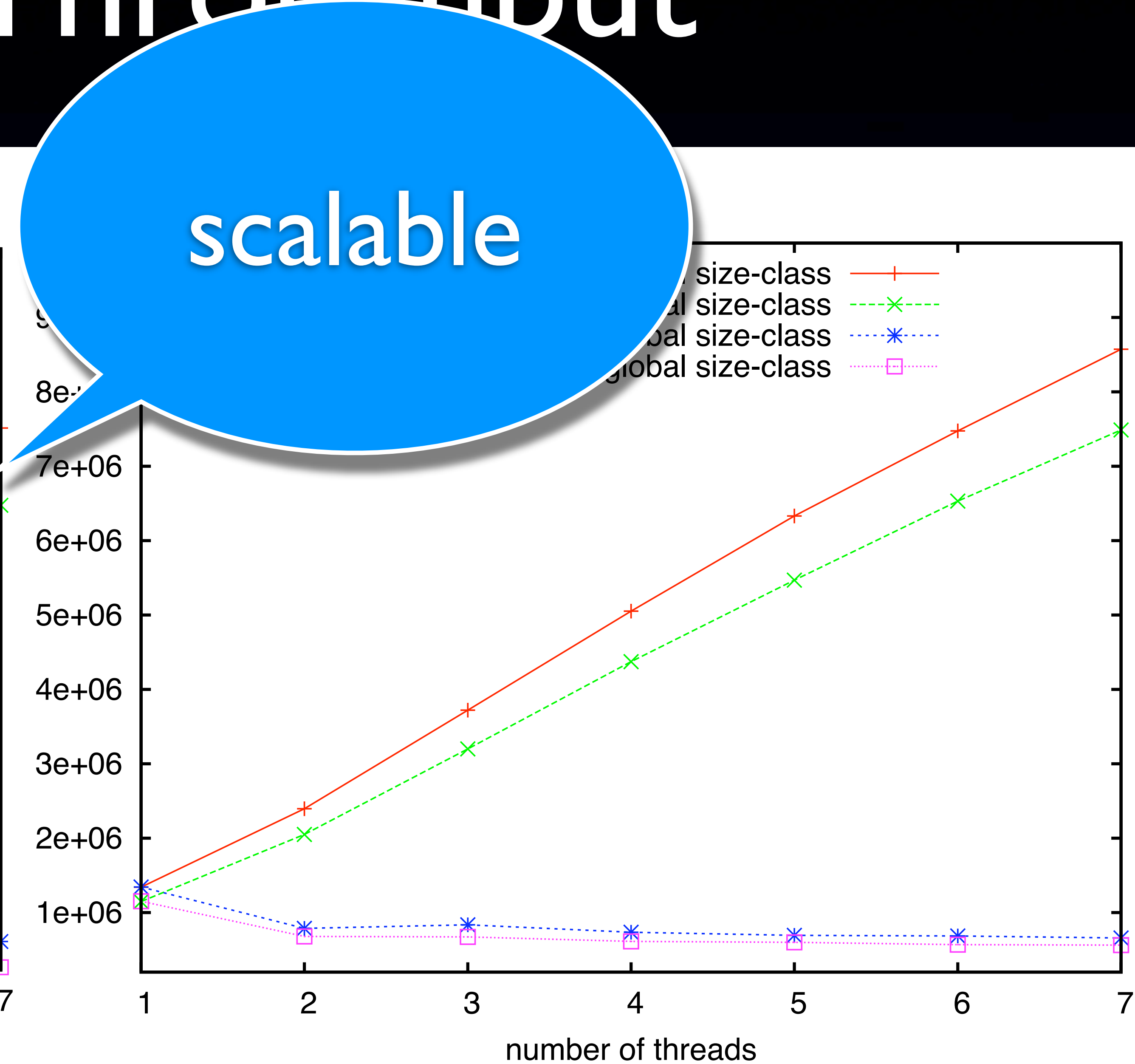
(a) full compaction

(b) optimized, non-compacting

Scalability of Allocation Throughput

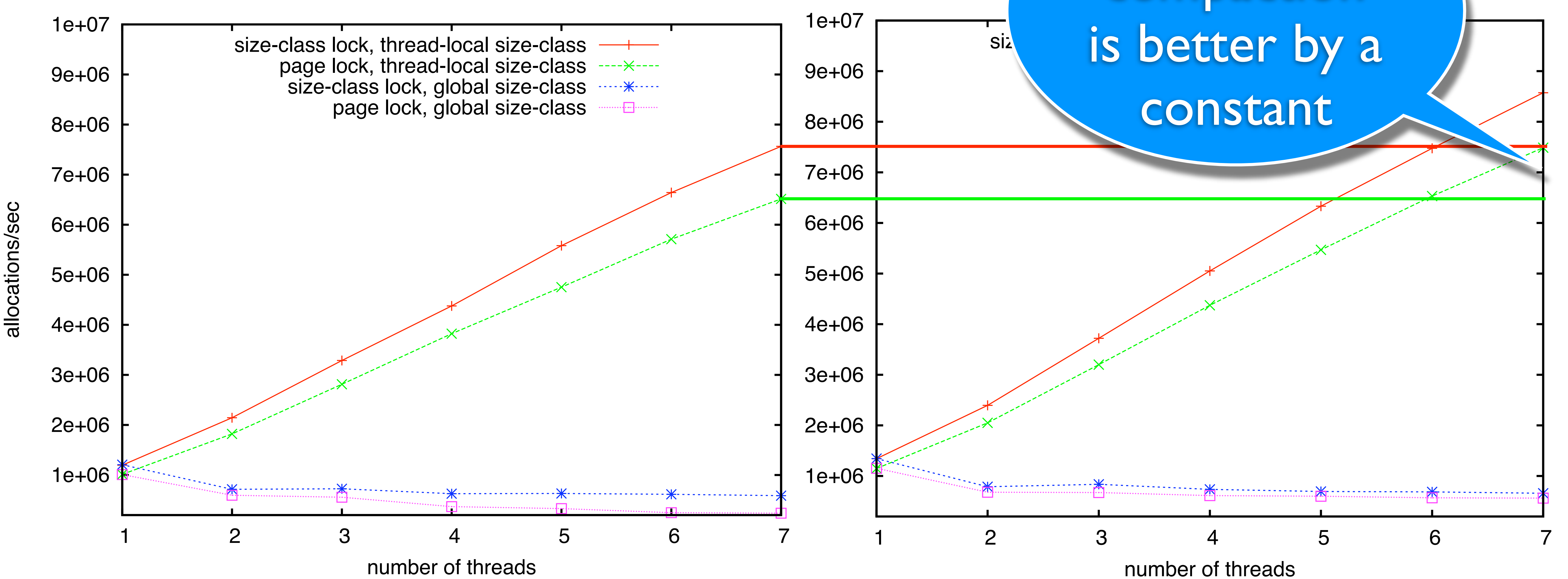


(a) full compaction



(b) optimized, non-compacting

Scalability of Allocation Throughput

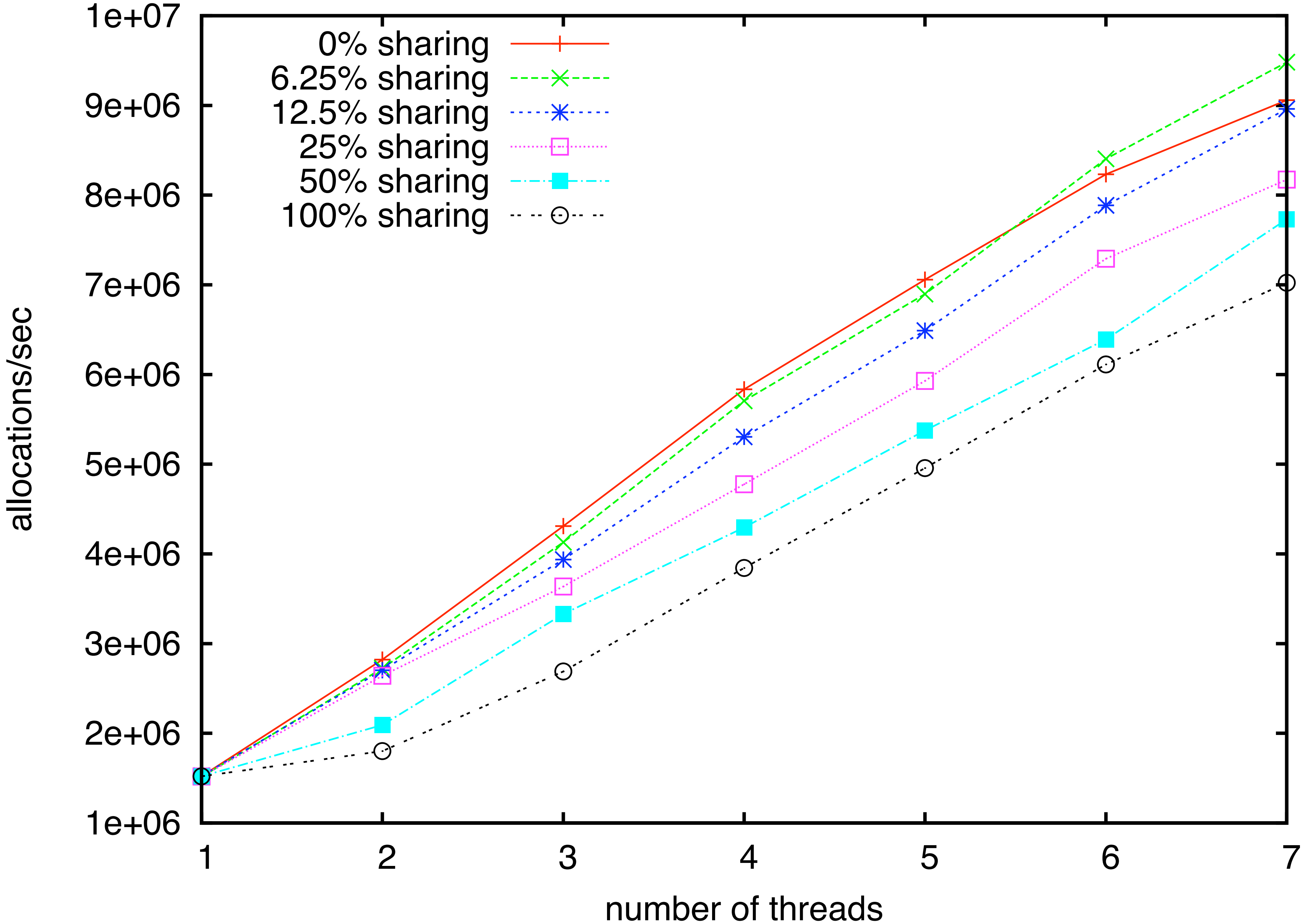


(a) full compaction

(b) optimized, non-compacting

- global size-class locks do not scale
- full compaction only requires constant factor

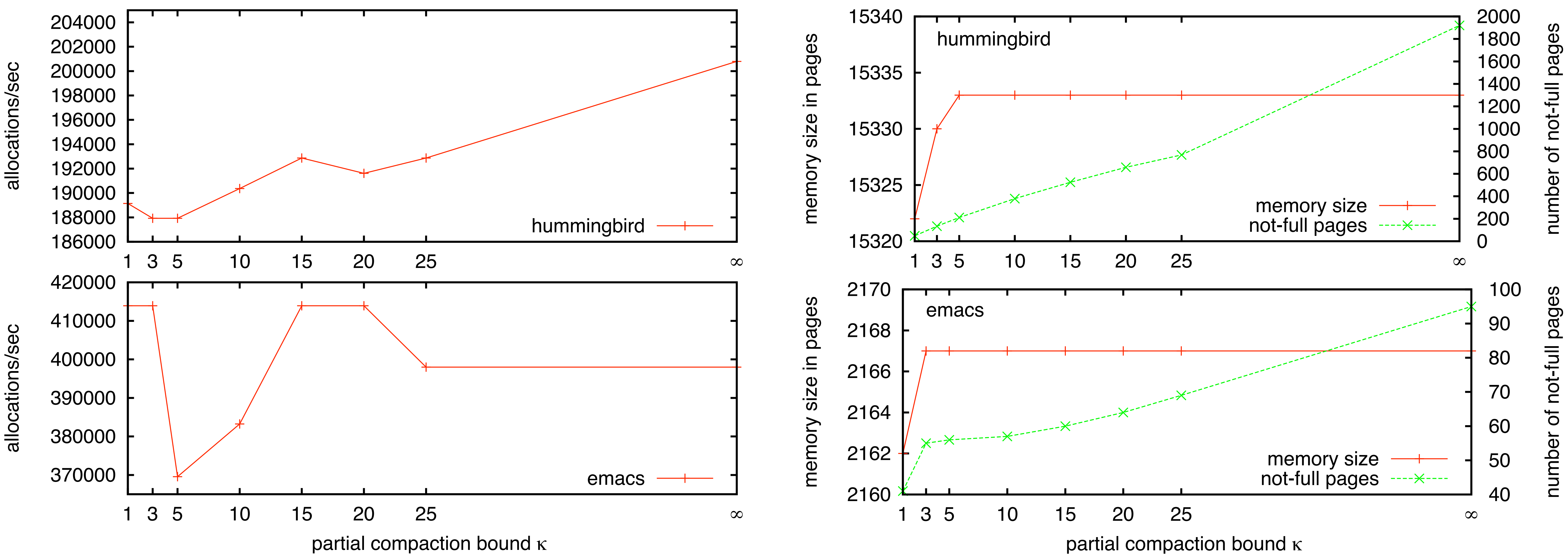
Scalability of Allocation Throughput



(c) opt., non-comp. with sharing

- level of sharing determines scalability

Real Application Performance



- less compaction may result in better allocation throughput
- size-class fragmentation increases with less compaction but total memory consumption may not

TLSF vs. opt., non-comp. CF Performance

	memory (in MB)				
	TLSF	CF (16B blocks)		CF (32B blocks)	
	memory size	memory size	size-class fragmentation	memory size	size-class fragmentation
Emacs	25.7	34.6	0.46	34.5	0.38
Hummingbird	203.7	245.3	8.3	245.9	11.4

	malloc (in clock ticks)				free (in clock ticks)			
	TLSF		CF		TLSF		CF	
	avg time	max time	avg time	max time	avg time	max time	avg time	max time
Emacs	228	93359	260	81662	153	71159	279	74798
Hummingbird	411	109079	529	98820	500	69192	574	79914

TLSF vs. opt., non-comp. CF Performance

	memo				
	TLSF	CF (16B bloc			cks)
	memory size	memory size	size-cl fragment		-class gmentation
Emacs	25.7	34.6	0.46	34.5	0.38
Hummingbird	203.7	245.3	8.3	245.9	11.4

only 1.3%
of the 35% more
memory

	malloc (in clock ticks)				free (in clock ticks)			
	TLSF		CF		TLSF		CF	
	avg time	max time	avg time	max time	avg time	max time	avg time	max time
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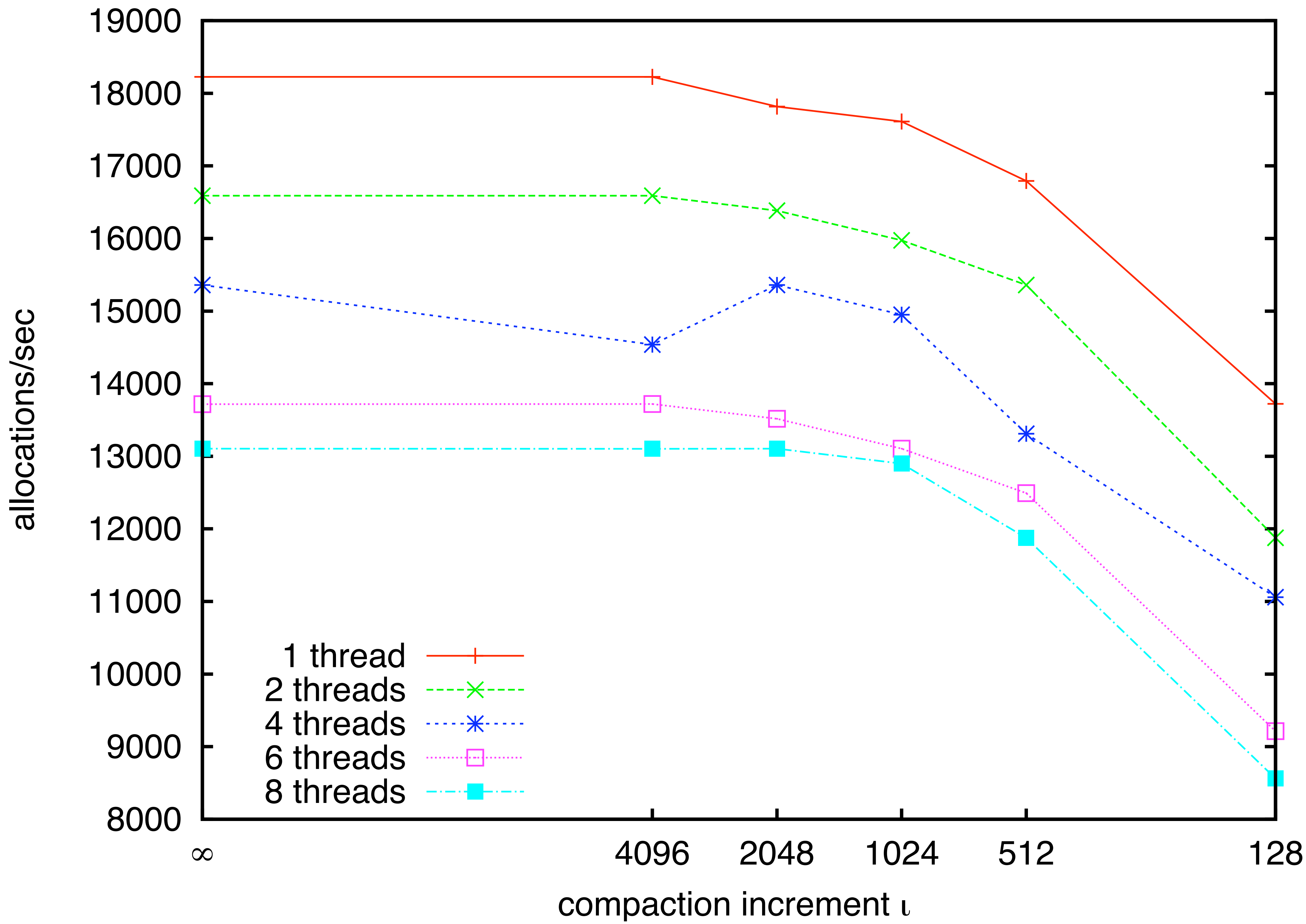
	malloc (in clock ticks)				[illegible] ks)			
	TLSF		CF		[illegible]		[illegible]	
	avg time	max time	avg time	max time	[illegible]	[illegible]	[illegible] me	max time
Emacs	228	93359	260	81662	153	71159	279	74798
Hummingbird	411	109079	529	98820	500	69192	574	79914

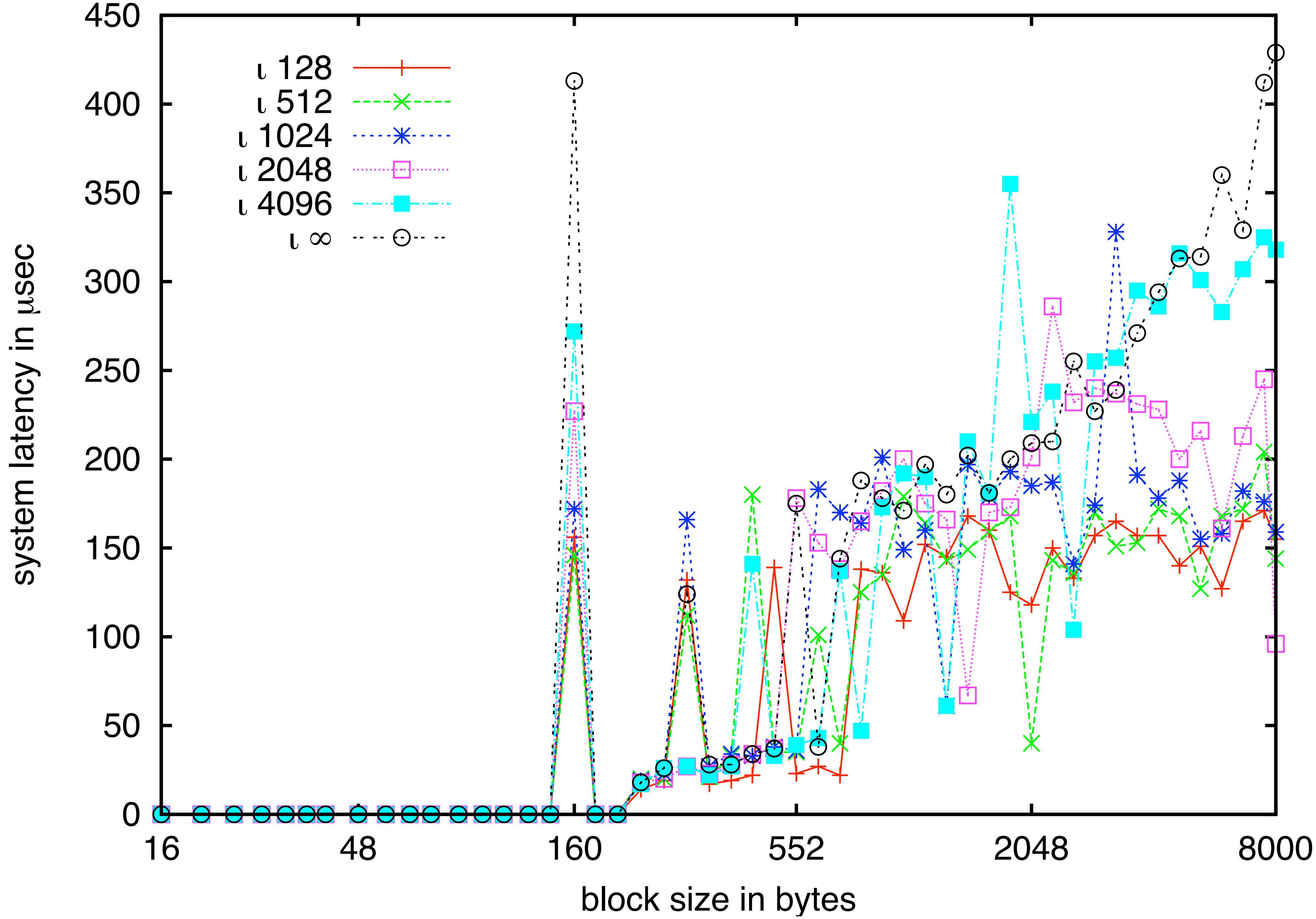
sometimes even better than TLSF

- fragmentation through partitioning dominates CF memory consumption
- opt., non-comp. CF only slightly slower than TLSF

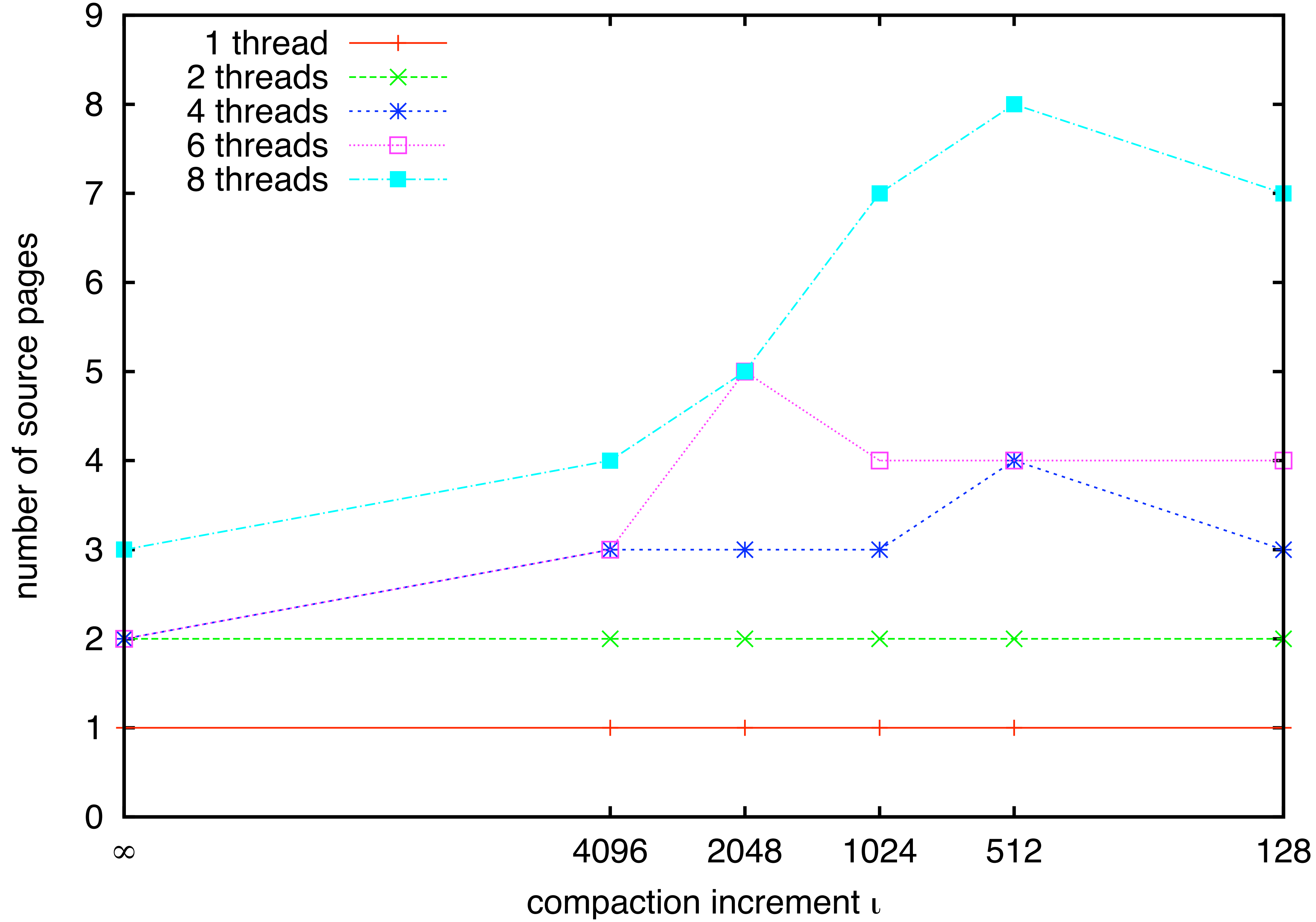
Allocation Throughput with Decreasing Compaction Increment



System Latency with 8 Threads and Increasing Block Size



Transient Size-Class Fragmentation with Decreasing Compaction Increment




Thank you