

# Computational Systems Engineering

Christoph Kirsch  
University of Salzburg

3 Unit Graduate Course, Winter 2004/2005  
Chapter 4: Benchmarking

# Heisenberg

- Uncertainty Relation:

The more precisely the position is determined, the less precisely the momentum is known in this instant, and vice versa.

- Uncertainty Principle

# Throughput

- the rate at which a system processes objects end-to-end
- unit: objects/unit of time
- example: bits/sec, cars/day, people/year
- it is not speed, which is distance/unit of time
- good for absolute performance
- bad for perceived performance

# Latency

- the time a system needs to process an object end-to-end
- unit: unit of time
- example: milliseconds
- again, it is not speed
- jumbo vs. supersonic, truck vs. fiber
- throughput and latency are opposed goals

# It's the Latency, Stupid

- Stuart Cheshire, May 1996
- Making more throughput is easy
- Get 10 phone lines instead of one
- Making less latency is not
- It is better to have a small share of a high-throughput, low-latency line than to have all of a low-throughput, high-latency line

# Stanford - MIT, 1996

- The distance from Stanford to Boston is 4320km.
- The speed of light in vacuum is  $300 \times 10^6$  m/s.
- The speed of light in fibre is roughly 66% of the speed of light in vacuum.
- The speed of light in fibre is  $300 \times 10^6$  m/s \* 0.66 =  $200 \times 10^6$  m/s.
- The one-way delay to Boston is  $4320 \text{ km} / 200 \times 10^6 \text{ m/s} = 21.6\text{ms}$ .
- The round-trip time to Boston and back is 43.2ms.
- The current ping time from Stanford to Boston over today's Internet is about 85ms:  

```
[cheshire@nitro]$ ping -c 1 lcs.mit.edu  
PING lcs.mit.edu (18.26.0.36): 56 data bytes  
64 bytes from 18.26.0.36: icmp_seq=0 ttl=238 time=84.5 ms
```
- So: the hardware of the Internet can currently achieve within a factor of two of the speed of light.

# Salzburg - Berkeley, 2004

- Wahine: ck\$ traceroute www.eecs.berkeley.edu  
traceroute to web1.eecs.berkeley.edu (169.229.60.94), 30 hops max, 40 byte packets
- |    |  |            |            |            |
|----|--|------------|------------|------------|
| 1  | 10.0.0.254 (10.0.0.254)                                    | 24.201 ms  | 3.239 ms   | 3.303 ms   |
| 2  | 81-223-189-113.itzling.xdsl-line.inode.at (81.223.189.113) | 157.561 ms | 13.088 ms  | 31.037 ms  |
| 3  | voe2-vl-00-010.voesend.vien.inode.at (62.99.170.221)       | 18.832 ms  | 21.808 ms  | 80.836 ms  |
| 4  | vie2-vl-00-020.shuttle.vien.inode.at (62.99.170.205)       | 33.394 ms  | 23.258 ms  | 16.192 ms  |
| 5  | otta-gb-03-002.shuttle.vien.inode.at (62.99.170.17)        | 17.28 ms   | 26.757 ms  | *          |
| 6  | ff-m-po-02-001.frankfm.germ.inode.de (62.99.170.106)       | 155.312 ms | 32.134 ms  | 32.479 ms  |
| 7  | po1-0.core01.fra03.atlas.cogentco.com (80.81.192.63)       | 75.607 ms  | 42.227 ms  | 34.975 ms  |
| 8  | p12-0.core01.dca01.atlas.cogentco.com (154.54.1.17)        | 131.743 ms | 125.974 ms | 119.981 ms |
| 9  | p15-0.core01.dca01.atlas.cogentco.com (66.28.4.21)         | 130.403 ms | *          | 127.108 ms |
| 10 | p10-0.core02.sfo01.atlas.cogentco.com (66.28.4.209)        | 303.818 ms | 203.276 ms | 201.534 ms |
| 11 | cenic.demarc.cogentco.com (38.112.6.226)                   | 193.605 ms | 208.736 ms | 206.211 ms |
| 12 | inet-ucb--lax-isp.cenic.net (137.164.24.142)               | 304.41 ms  | 264.708 ms | 190.957 ms |
| 13 | vlan194.inr-202-doecev.berkeley.edu (128.32.0.251)         | 194.394 ms | 196.212 ms | 234.06 ms  |
| 14 | doecev-soda-br-6-4.eecs.berkeley.edu (128.32.255.170)      | 192.544 ms | 214.33 ms  | 194.122 ms |
| 15 | sbd2a.eecs.berkeley.edu (169.229.59.226)                   | 203.337 ms | 437.41 ms  | 427.557 ms |
| 16 | web1.eecs.berkeley.edu (169.229.60.94)                     | 354.457 ms | 197.736 ms | 210.725 ms |

# Thrashing

- a condition in which a computational system spends more time administrating resources than utilizing resources
- little or no progress of application functionality
- typically patterns of requests for resources followed by inadequate access emerge
- examples: virtual memory page faults, network collisions



# What to Benchmark?

- throughput: bits/seconds
- latency: milliseconds
- error rate: %
- maximum number of connections
- thrashing
- system time: disk, network, reactor, scheduler
- user time: parser, error handler, generator

# How to Benchmark

- Hardware: server plus multiple client machines plus network switch
- Software: httpperf, apachebench, autobench
- Problems: bottlenecks (network, clients)
- Solution: vary network/client capacity
- Check out: SPECweb99

# Static Benchmarking

- request the same file many times
- simple but not real world (caching)
- tool: `httperf`
- ```
httperf --server host.mydomain.com \  
        --uri /index.html \  
        --num-conn 5000 \  
        --num-call 10 \  
        --rate 200 \  
        --timeout 5
```

# Automatic Benchmarking

- request the same file many times at increasing rates
- tool: autobench
- generates output in CSV/TSV format

# Dynamic Benchmarking

- request different files many times
- use sequences of requests obtained from web logs
- tool: `httperf`
- ```
httperf --server host.mydomain.com \  
--wssesslog 1000,2,session.log \  
--max-piped-calls 5 \  
--rate 20 \  
--timeout 5
```

# Jobs

- System Administration
- Website
- Benchmarking
- Webserver development
- Library development