

Terra: A Virtual Machine-Based Platform for Trusted Computing

**by Tal Garfinkel, Ben Pfaff, Jim Chow, Mendel
Rosenblum and Dan Boneh**



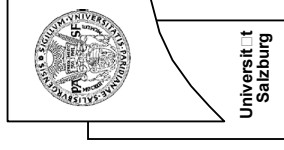
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Content

- Introduction
- What is Terra?
- How does Terra work?
- Prototype implementation
- Conclusion



Facts

- The variety of applications:



- Secure bank application

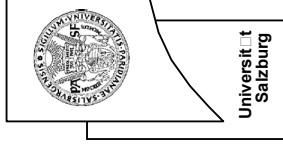


- Online game

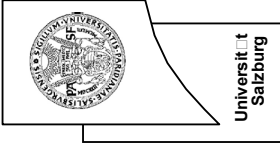
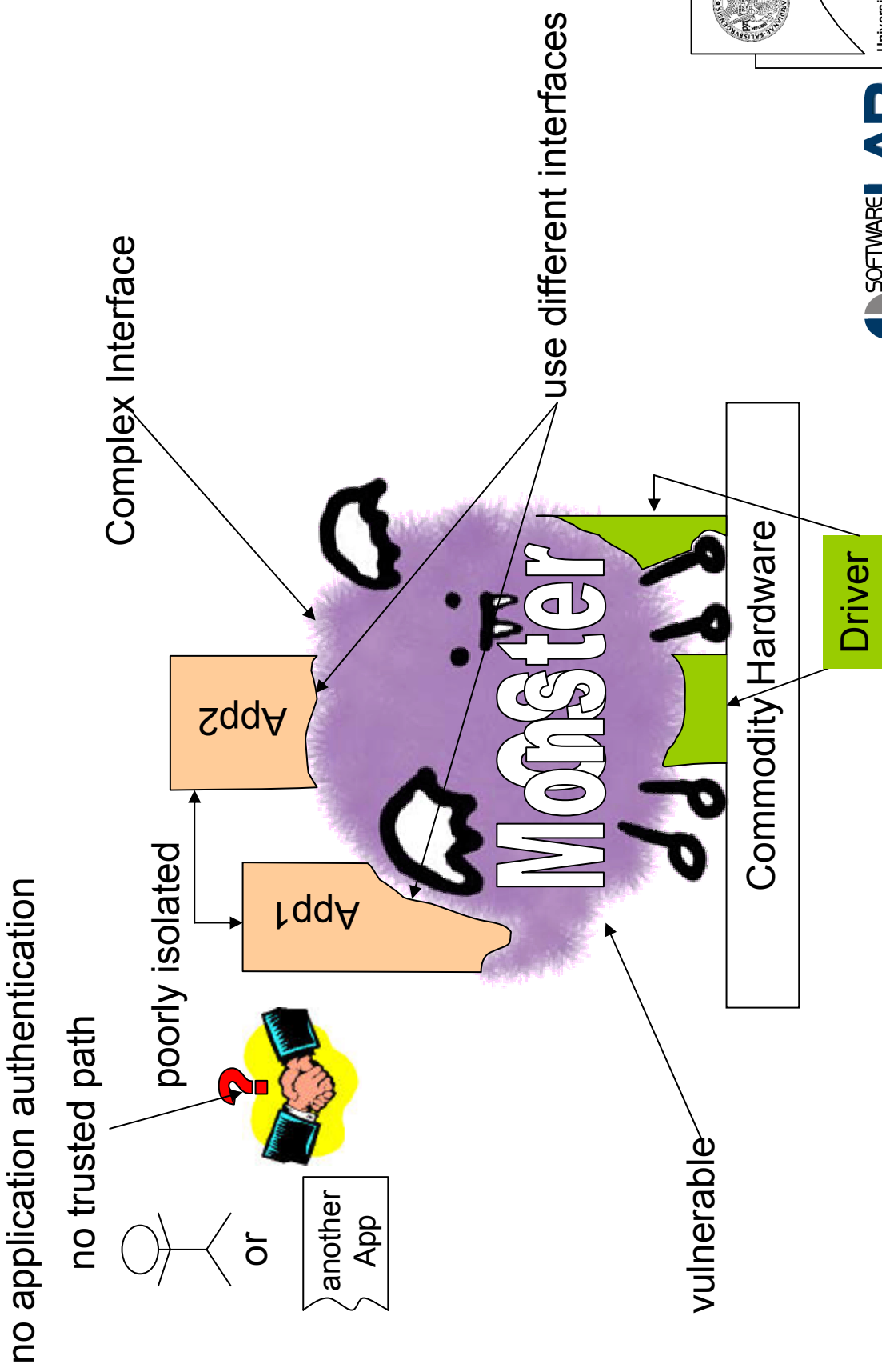


- Web application

} different security requirements



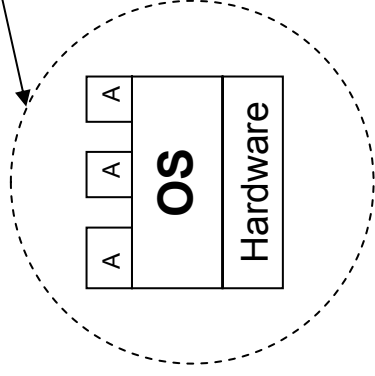
Common situation



Open vs. Closed Systems (I)



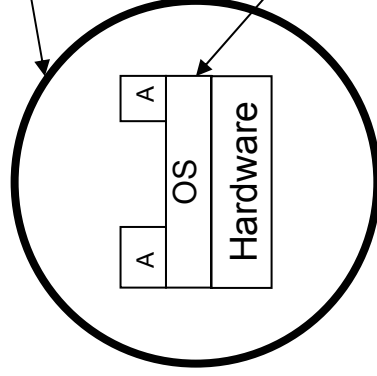
Open system, e.g.



→ Platform owner has full control

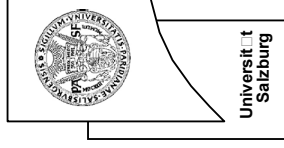


Closed System, e.g.



,thinner', more application tailored OS

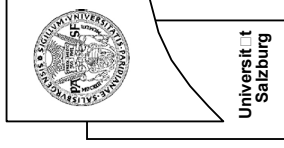
→ the developer has full control



Open vs. Closed Systems (II)

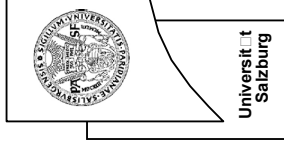
- Open Systems
 - Support for many application
 - Huge amount of existing code
 - Take advantage of commodity hardware
- Close Systems
 - Slimmer, more specialized OS
 - Applications are tailored to their security requirements
 - Provide hardware tamper resistance

What is Terra?



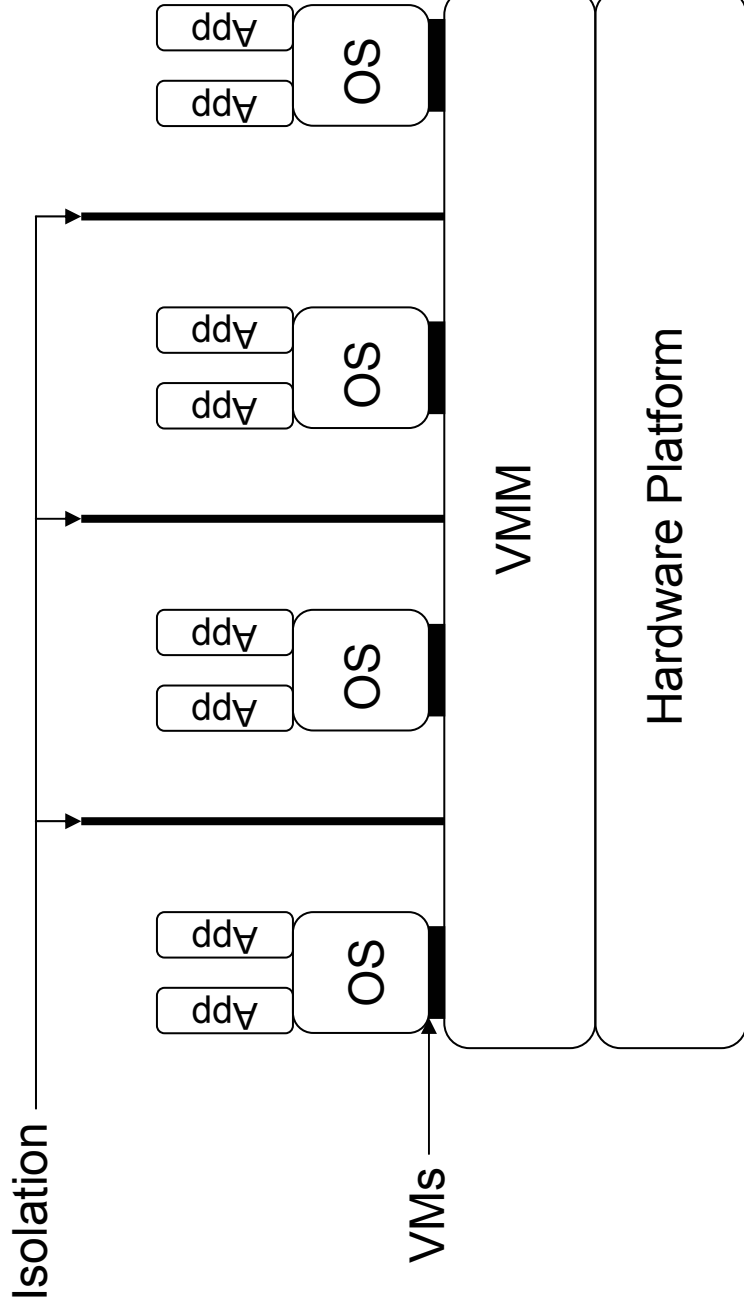
What does Terra try to achieve?

- Terra tries to combine the advantages of an Open and a Closed box.
 - Terra implements two principles:
 - *Isolation*
 - *Attestation*
- of applications



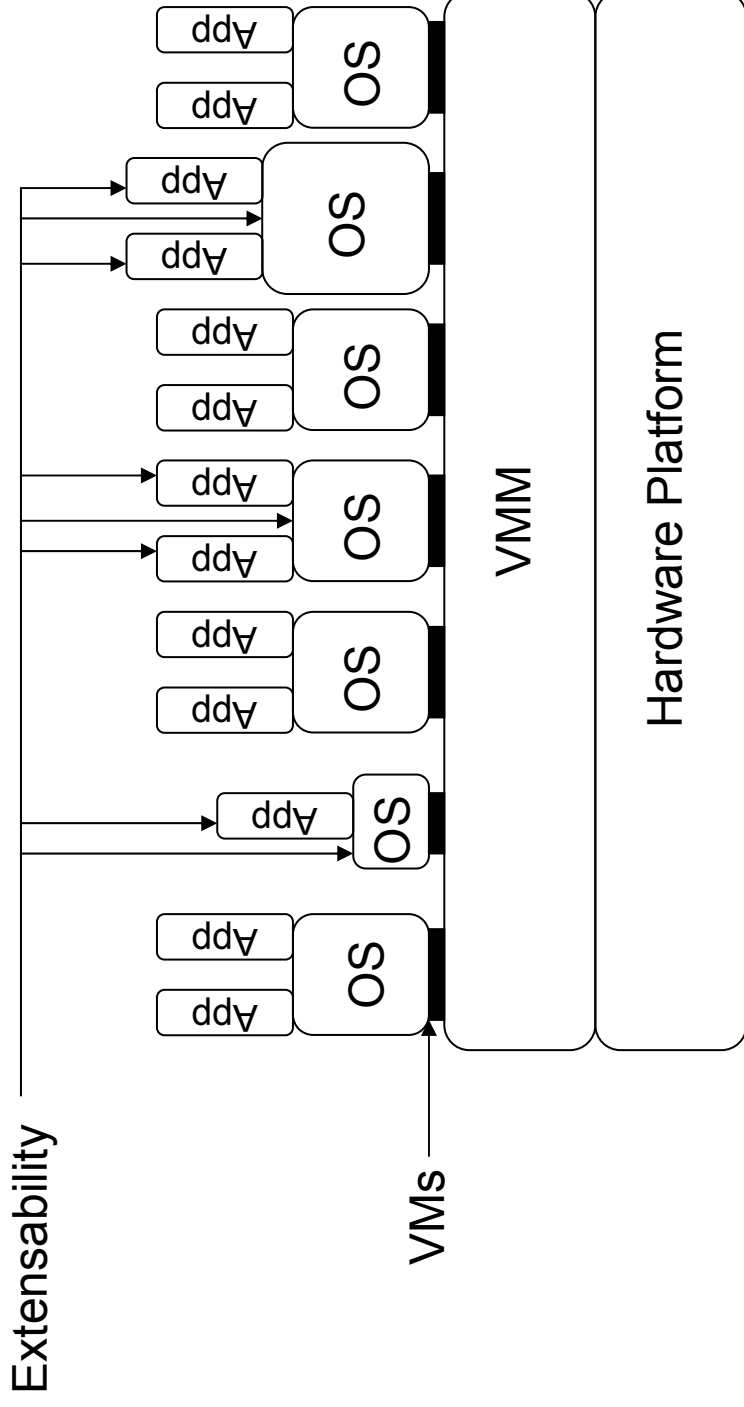
What is Terra?

- The heart of Terra is a *Virtual Machine Monitor (VMM)*



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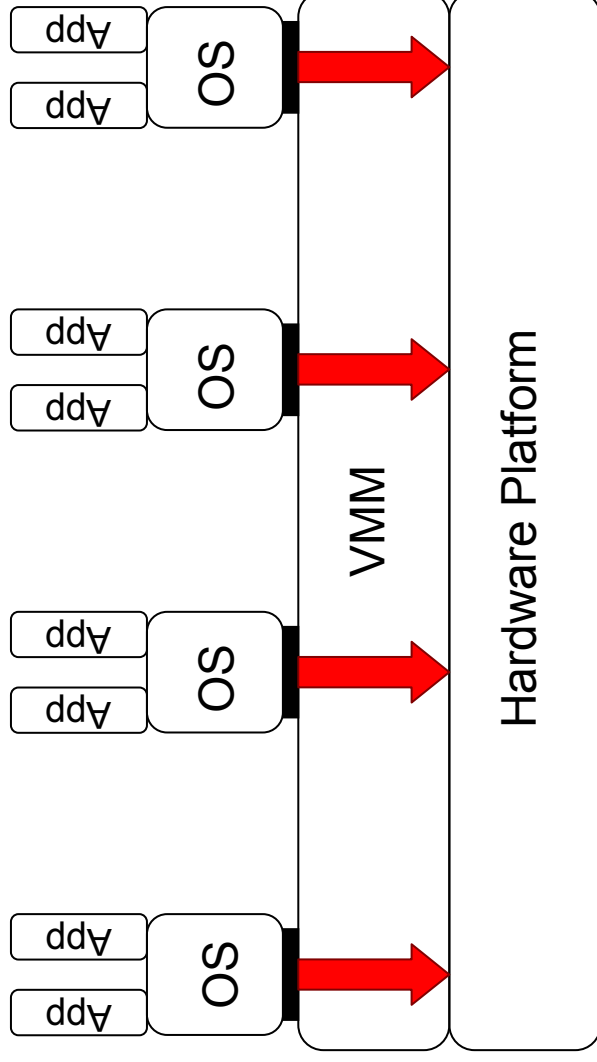


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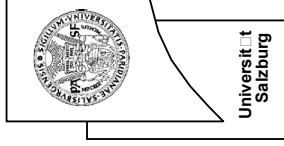
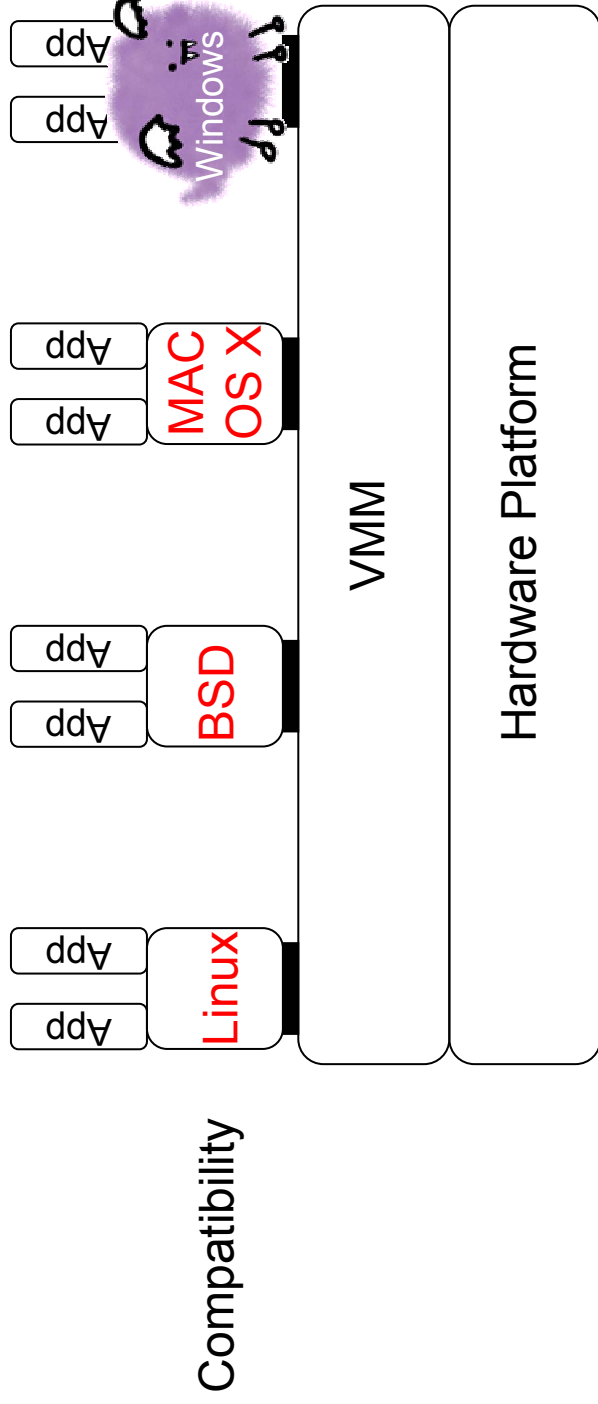


Efficiency



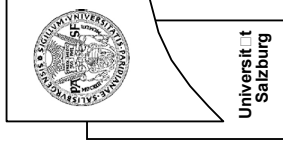
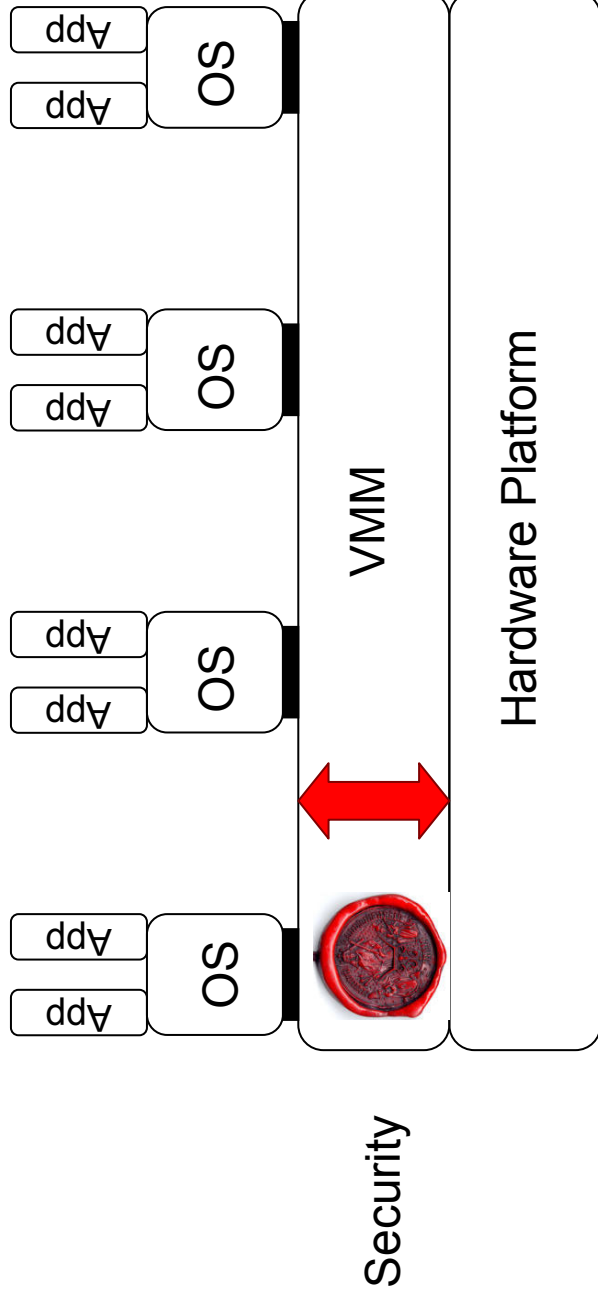
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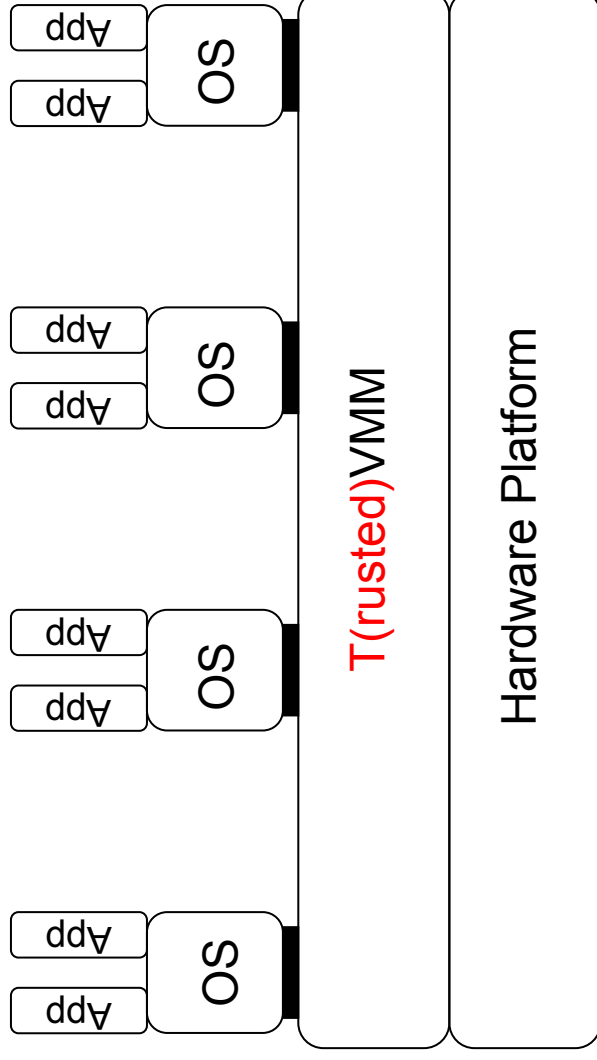
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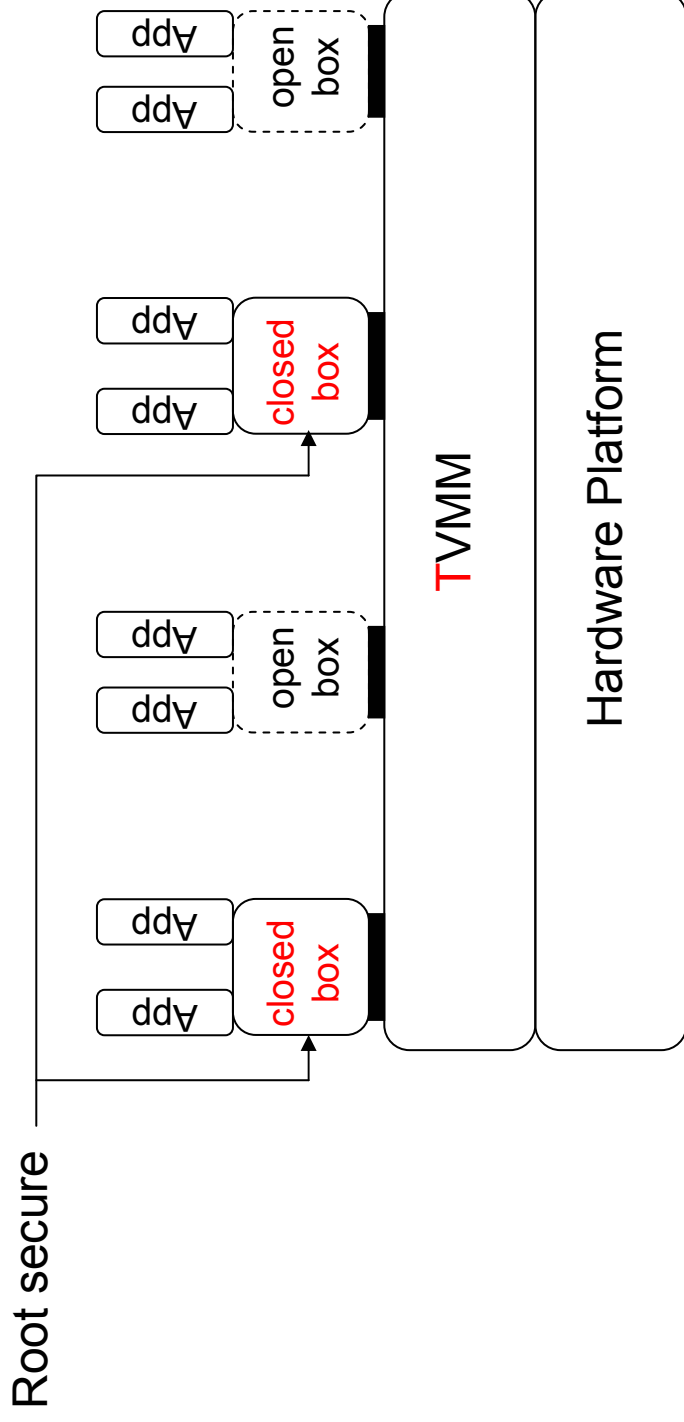
What is Terra?

- The heart of Terra is a **Trusted Virtual Machine Monitor (TVMM)**



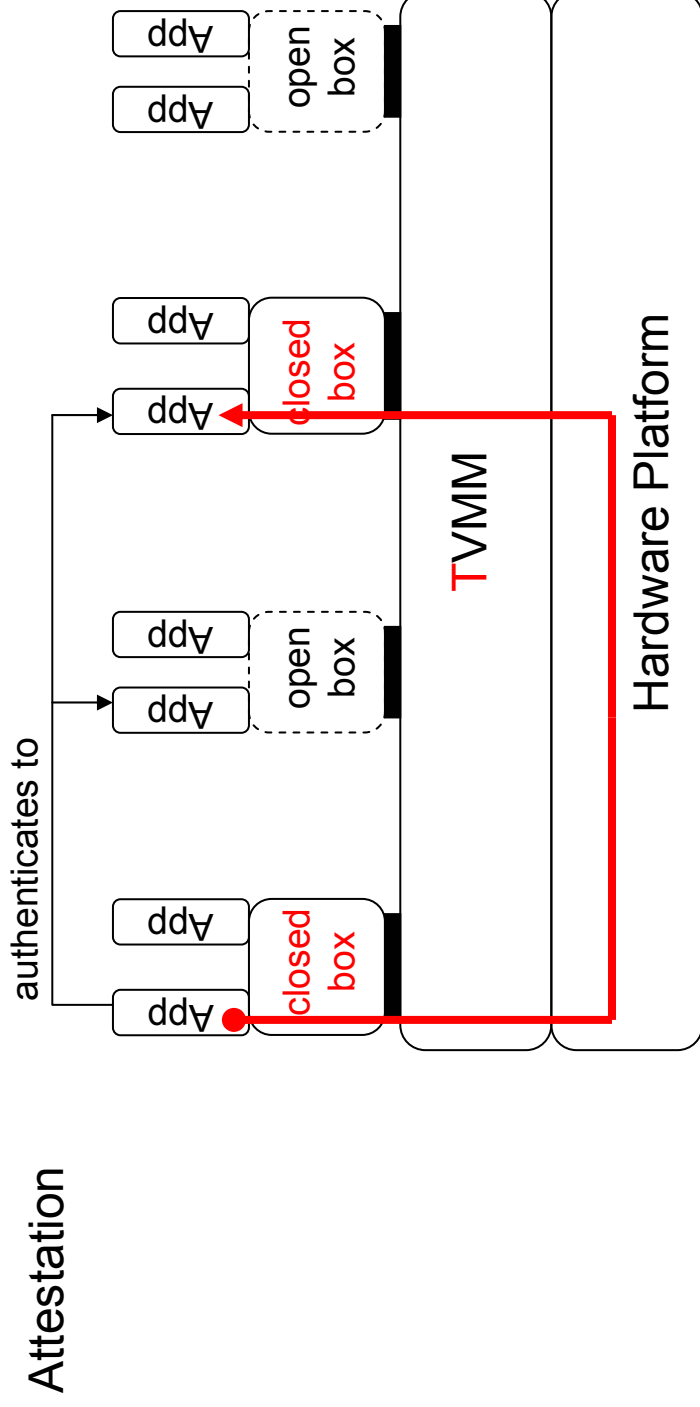
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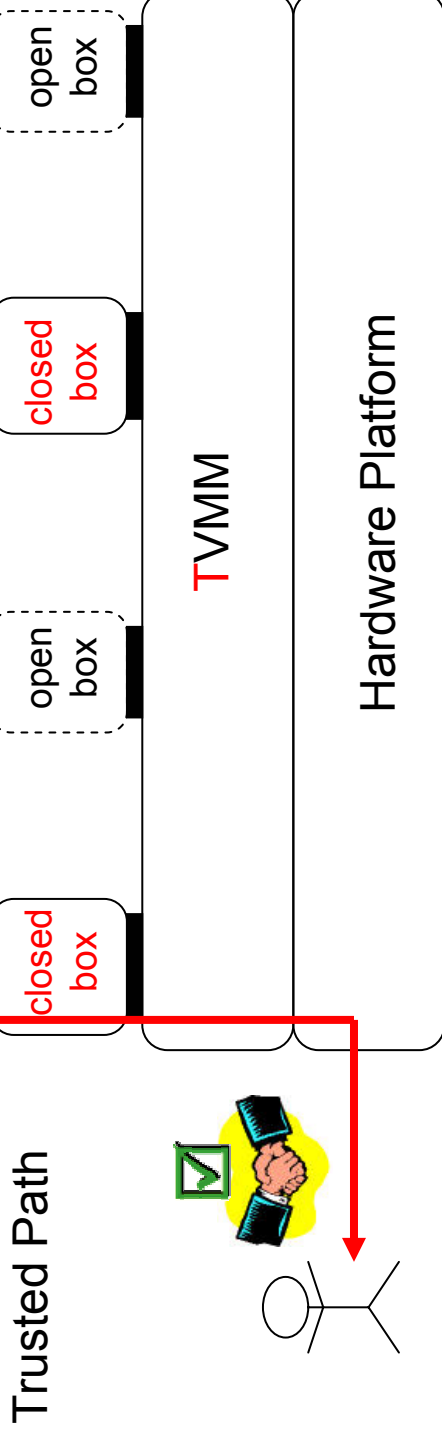
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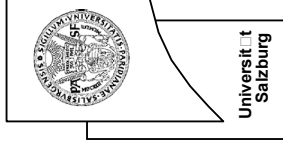


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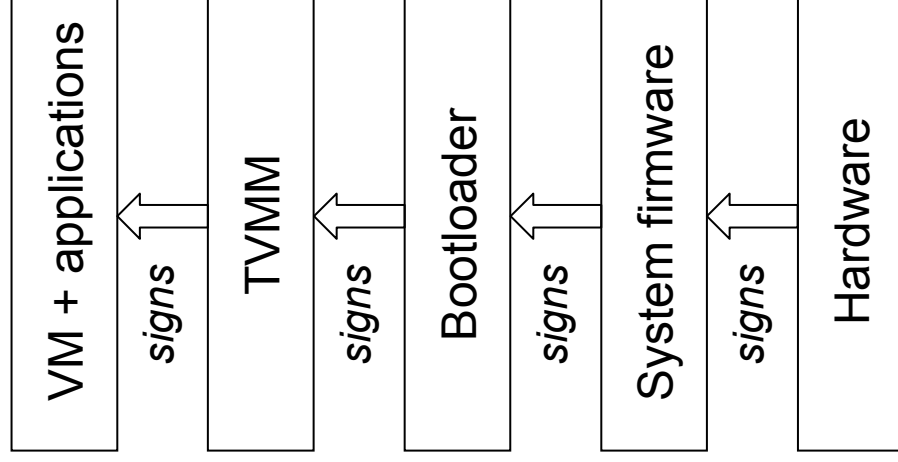


How does Terra work?

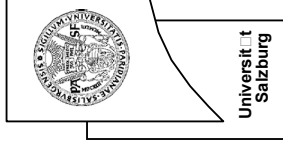


Attestation (I)

- Building a certificate chain

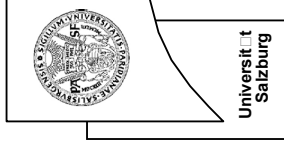
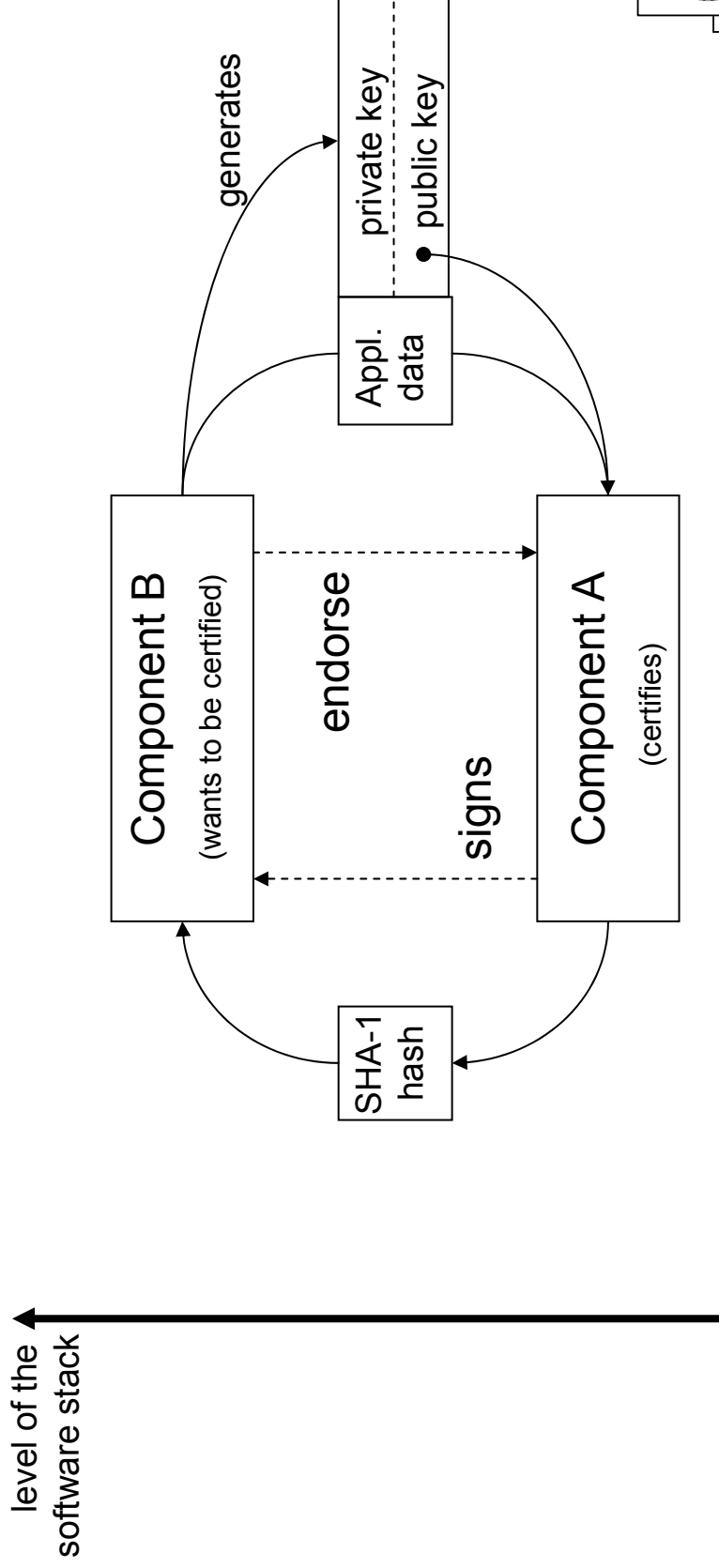


Tamper resistant chip (e.g. TCPA)



Attestation (II)

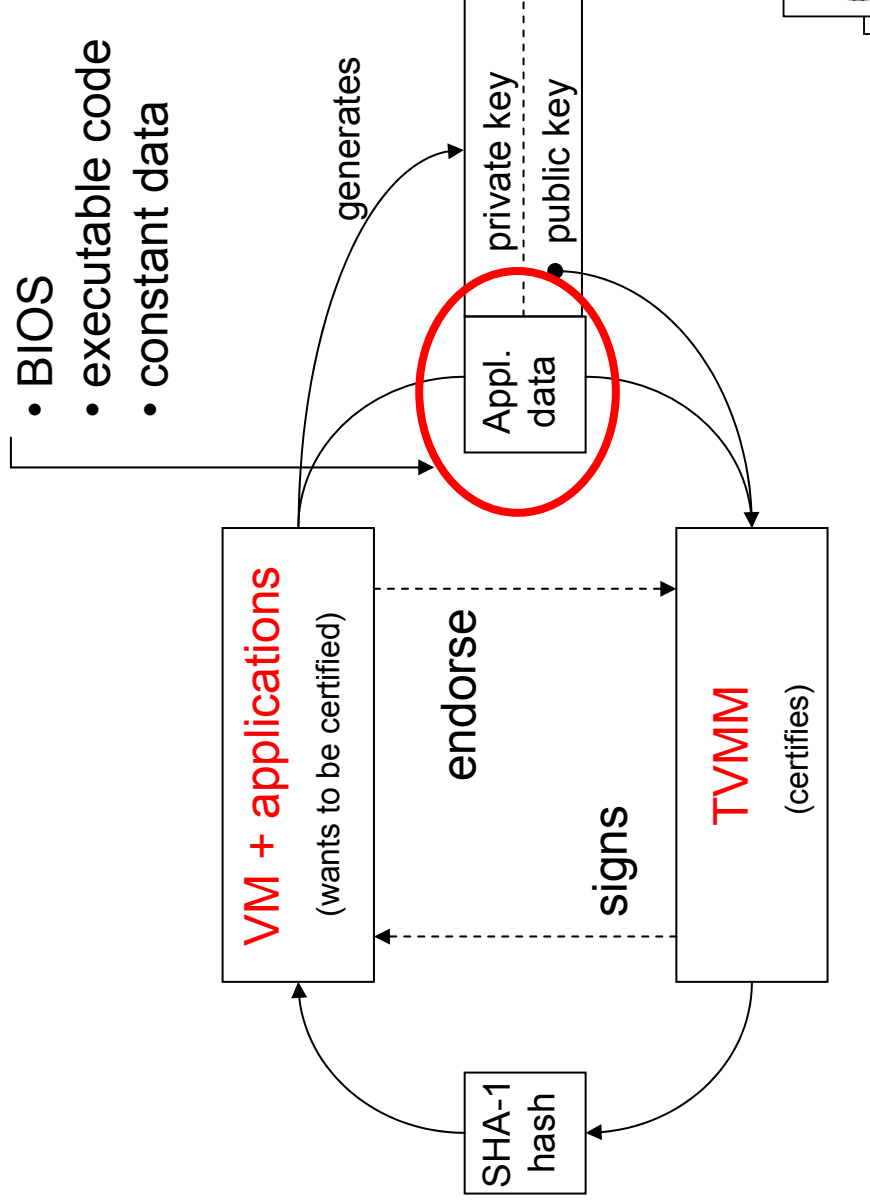
- Signing



Attestation (III)

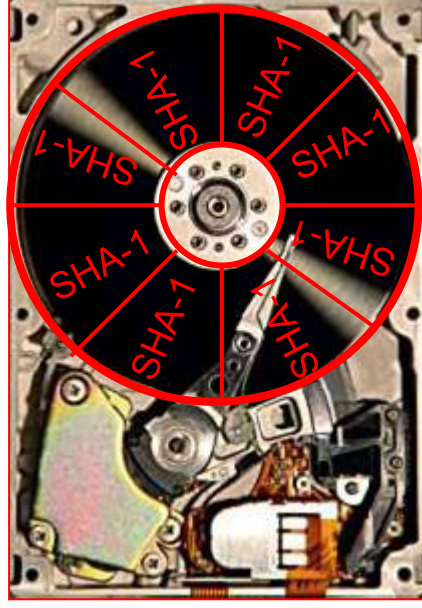
- Signing

level of the software stack



Implementing Attestation (I)

- A VM image consists of several parts (entities)
 - Mutable (NVRAM, cache, user data)
 - Immutable (BIOS, executables, constant data)
- List of hashes for attestable parts
- e.g. an attestable (virtual) disk



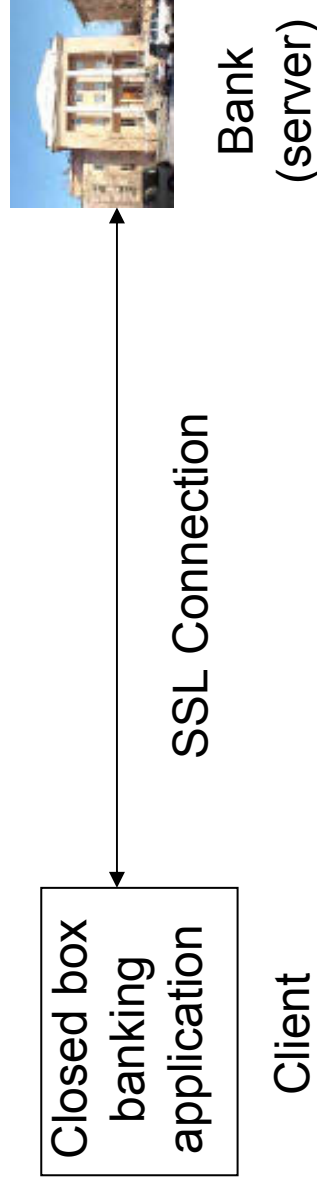
Implementing Attestation (II)

- Example:
 - Hashing a 4 GB entity into 20-byte SHA-1 with a 4 kB block size*
 - *20 MB of hashes*
 - *have to be verified against the VM descriptor hash*
 - *problem regarding memory and time*

Implementing Attestation (III)

- Ahead-of-Time Attestation
 - All components are proved during the boot process
 - Used only for small, high-assurance VMs
- Optimistic Attestation
 - If Ahead-of-Time is impractical
 - Individual blocks are „lazy“ checked

Example attestation

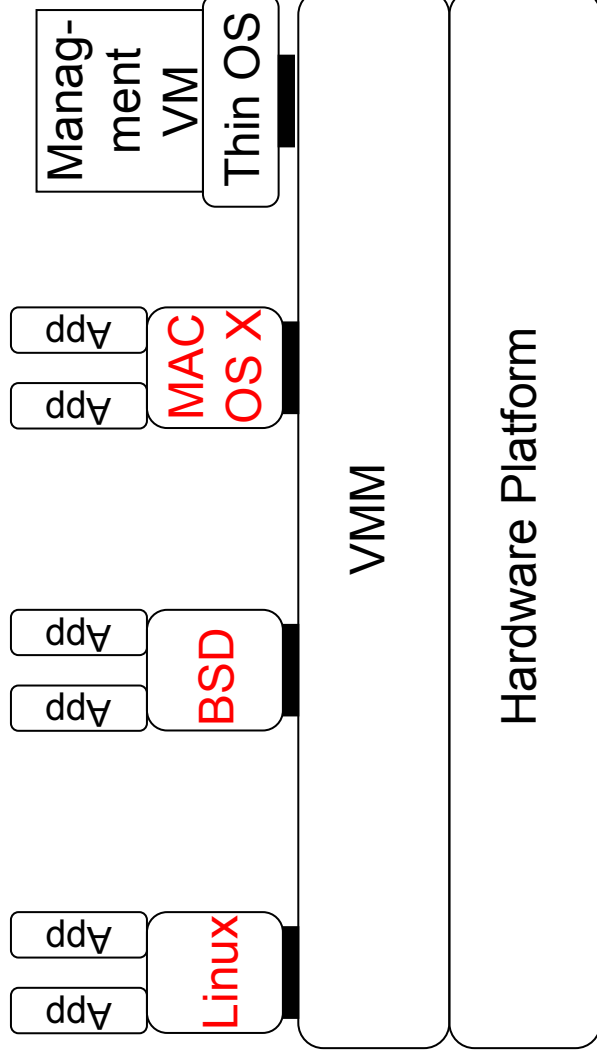


- Certificate is exchanged during SSL handshake
- The server verifies
 - lowest certification (hardware)
 - middle certification chain (BIOS, bootloader, TVMM)
 - highest certification (application, e.g. Quicken)



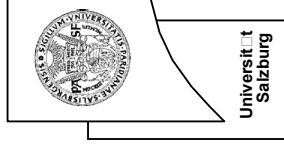
Management VM (I)

- High-Level configuration with the Management VM



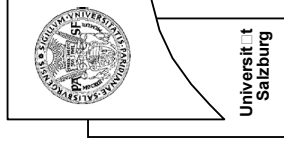
Management VM (II)

- API (i)
 - device-id ← CREATE-DEVICE(type, params)
 - CONNECT(device-id-1, device-id-2)
 - DISCONNECT(...)
 - vm-id ← CREATE-VM(config)



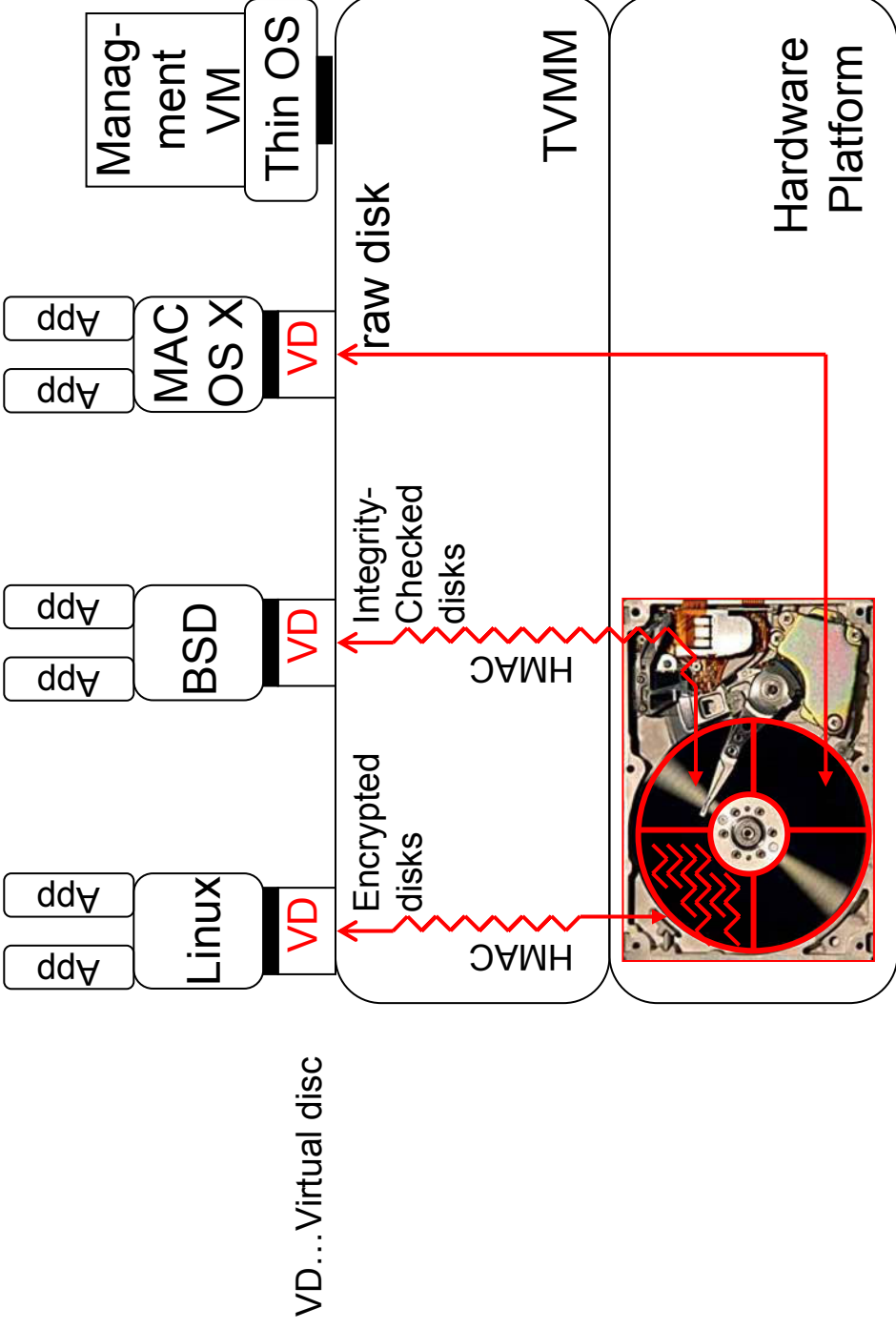
Management VM (III)

- API (ii)
 - ATTACH(vm-id, device-id)
 - DETACH(...)
 - ON(vm-id)
 - OFF(vm-id)
 - SUSPEND(vm-id)
 - RESUME(vm-id)



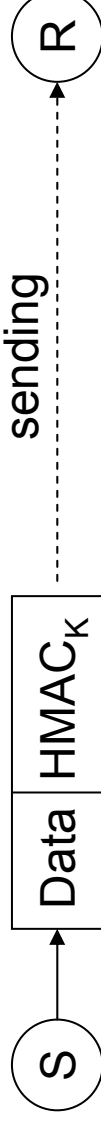
Storage Interface

- Three different types of disk access control

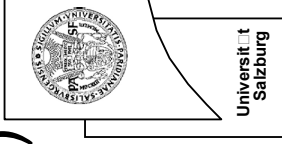


The HMAC Construction

- MAC – Messaging Authentication Code
- HMAC – MAC using hashes
- Using an authentication tag HMAC_K

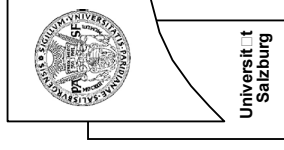


- $\text{HMAC}_K = H(K, H(K, \text{Data}))$ (H...SHA-1, K...Key)



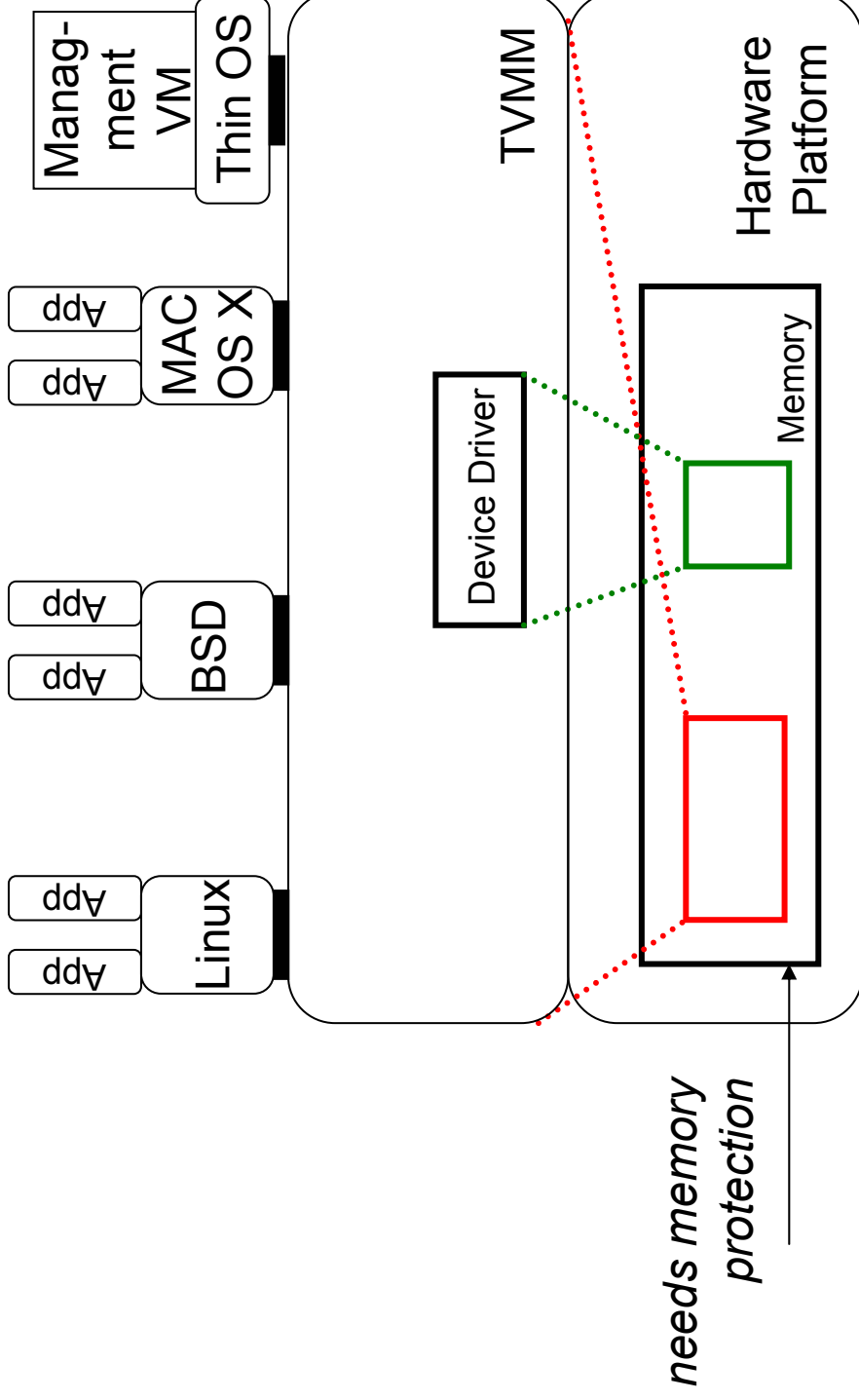
Device Driver Security (I)

- Today's driver can be very „large“
(e.g. video, modem, wireless network driver)
- Worst quality code in many cases
- → Device driver as part of the *TVM*'s trusted computing base?



Device Driver Security (II)

- High level configuration management

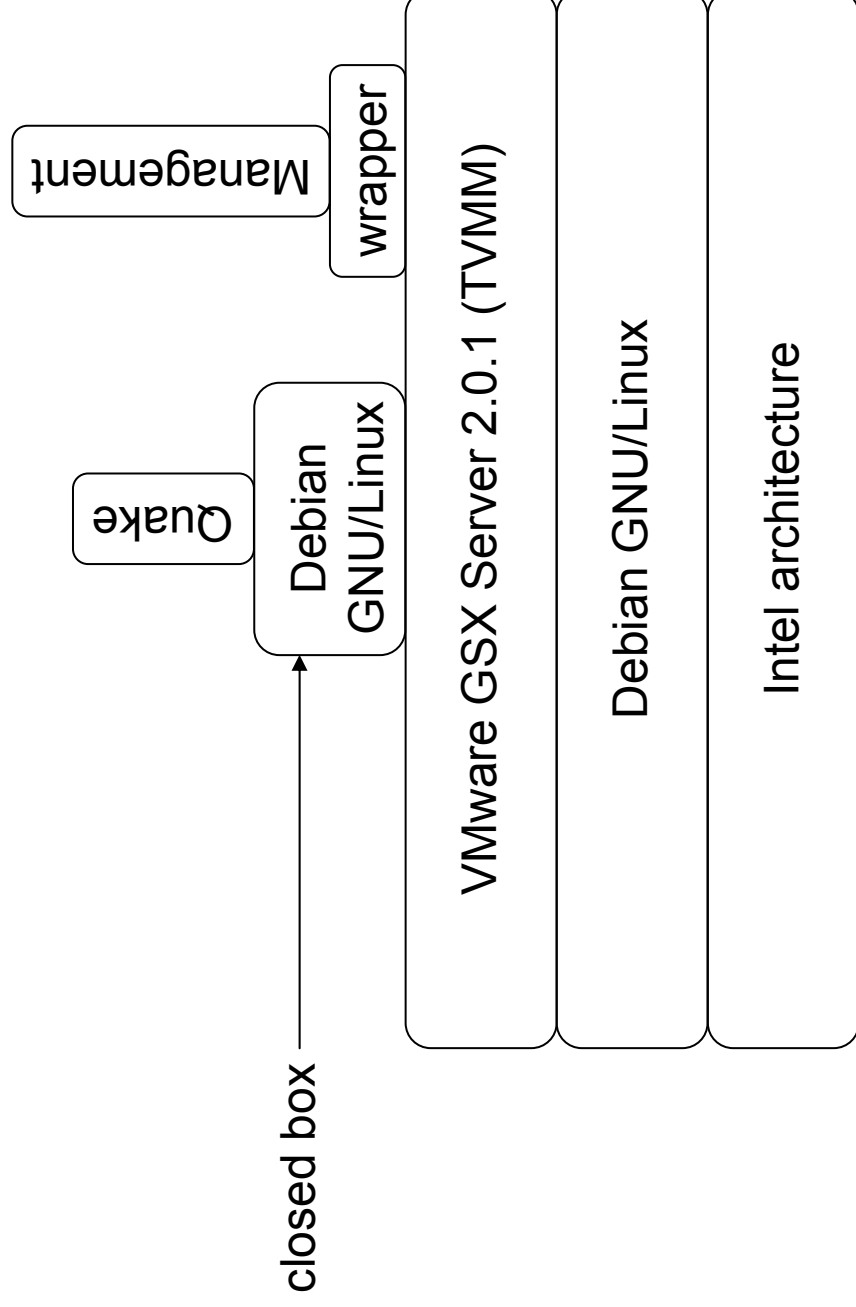


needs memory protection

Prototype implementation

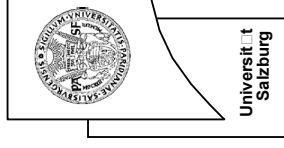
Prototype implementation (I)

- Used Architecture



Prototype implementation (II)

- Secure storage
 - Using a dynamic preloaded library
 - Library implements
 - Ahead-of-time attestation
 - Optimistic attestation
 - Integrity-checked storage
 - Encrypted storage



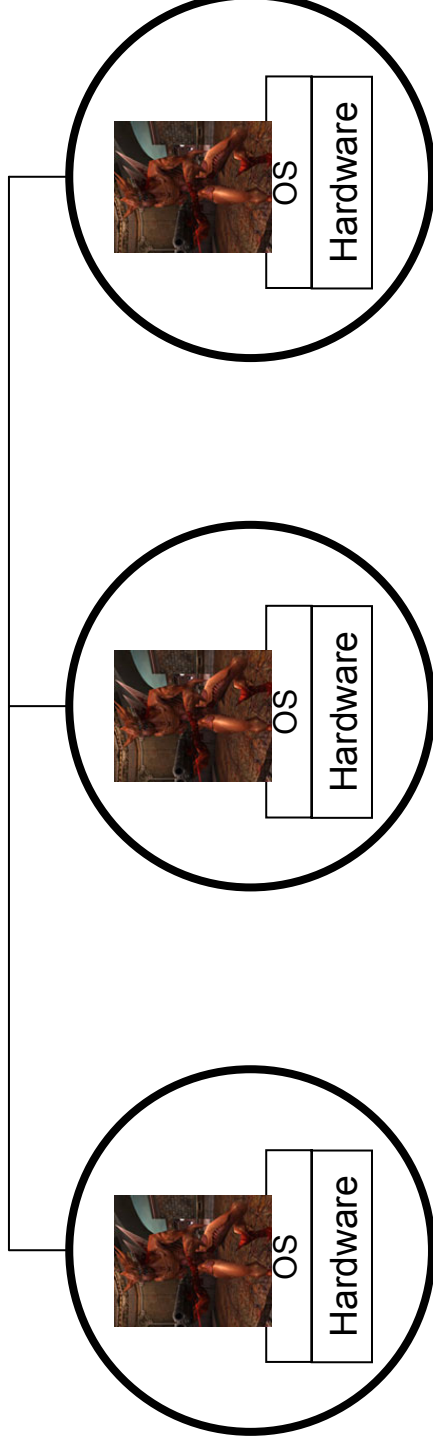
Prototype implementation (III)

- System management
 - VMware GSX Server provides configuration interface
 - Python wrapper implement the management API
 - OpenSSL library for certification management

Prototype implementation (IV)

- Trusted Quake

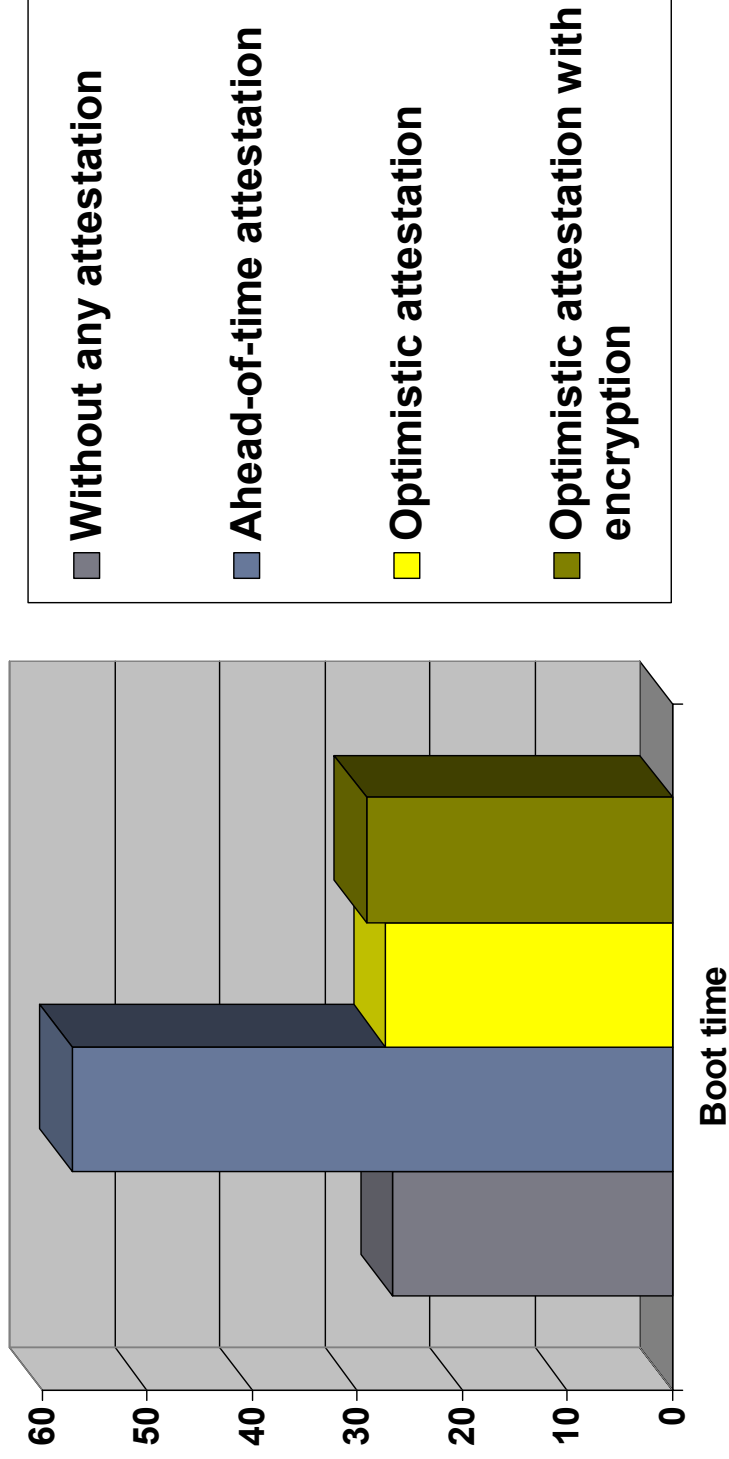
authenticate each other via attestation



- Linux boots directly into quake
- Attestation via dynamic preloaded library
- Using 160-bit SHA-1 HMAC and 56-bit DES keys

Prototype implementation (V)

- Measurement results



Conclusion

- Flexible architecture for trusted computing
- Allowing open box and closed box VMs
- Attestation and Isolation as basis for application authentication
- Application security can be tailored to their needs
- Still some „new architectural environment“ is needed

