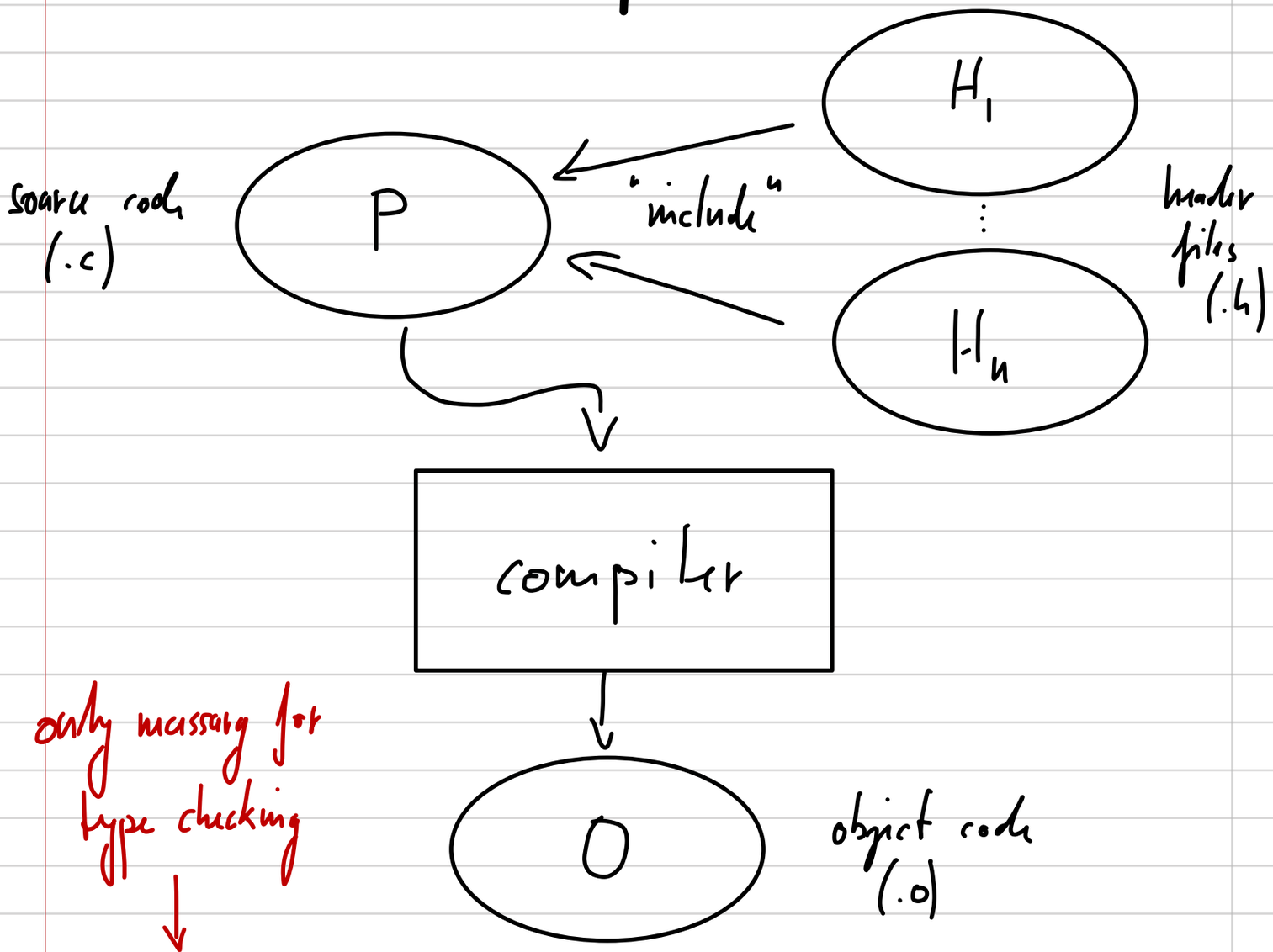
Compiler



only necessary for type checking

header file: syntactically equivalent to source code except that procedure bodies are missing → declarations only

compiler first parses all included header files → constructs symbol table, no code generation!

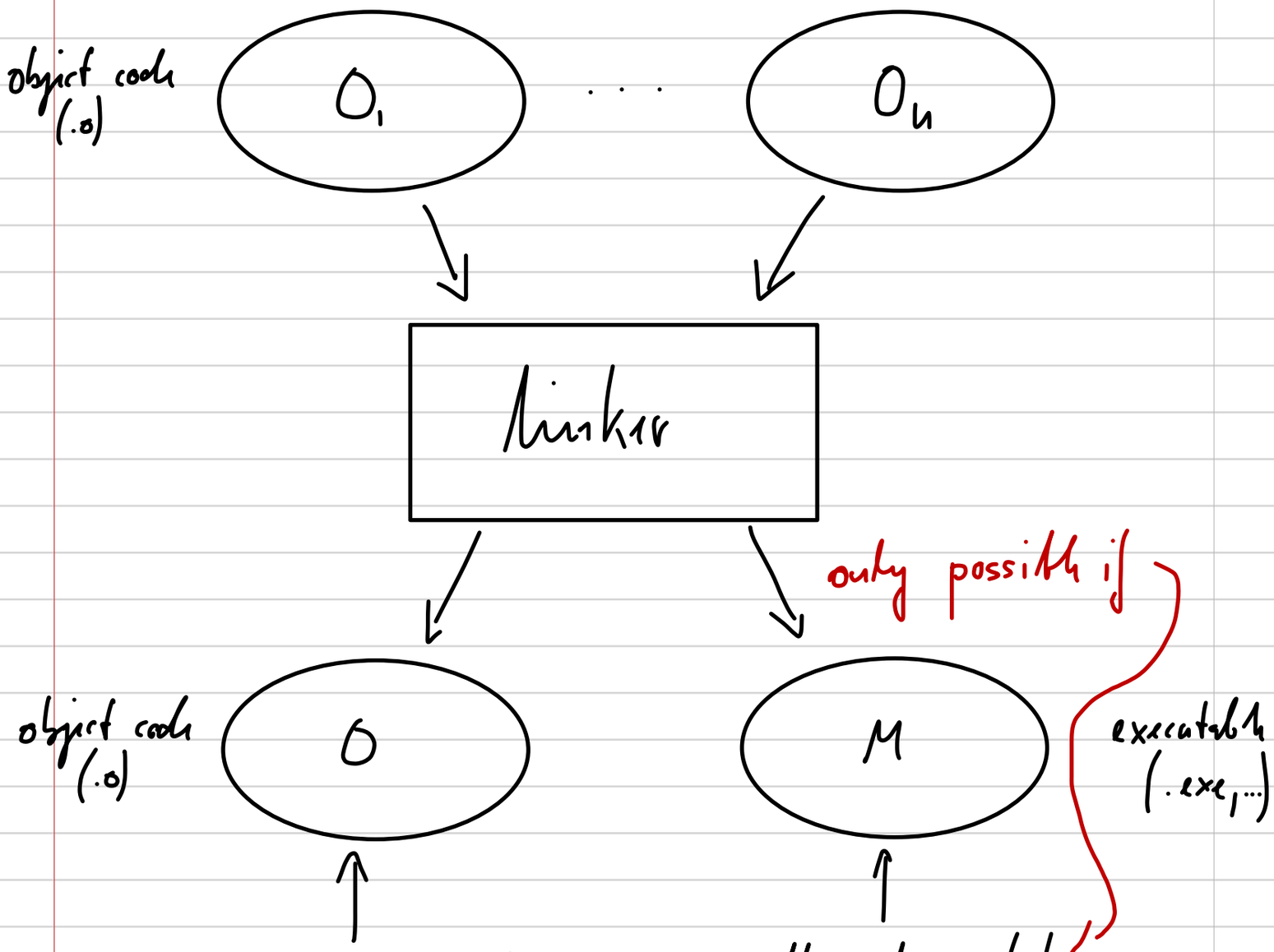
source code: uses global variables and procedures declared in imported header files

object code: contains relocatable code generated for source code, relocatable data (global vars, strings) for source code (not header files), and "symbols" (definitions, uses)

compiler then parses source code and generates object code

names of global variables, strings, and procedures

Linker



Java: .class files correspond to .o files but are only loaded and linked dynamically at runtime on demand } no executables!

[DLX Tools (wirth, Knight, Ashenden)] → DLX Object File Format

```
objectFile = relocatableSegment { relocatableSegment }
  { symdef } { linkup } { fixup } .

relocatableSegment = ( ".codeseg" | ".dataseg" ) segmentName
  segmentSize { memoryData } .
segmentName = "$" identifier .
segmentSize = integer .
memoryData = word .

.codeseg $module1_code 80 ...
.dataseg $module1_variables 20
.dataseg $module1_strings 160 ...
.codeseg $module2_code 120 ...
.dataseg $module2_variables 8
.dataseg $module2_strings 100 ...
...

symdef = ".symdef" symbolName relocatableAddress .
symbolName = "#" identifier .
relocatableAddress = segmentName ("+" | "-") offset .
offset = integer .

.symdef #x $module1_variables-4
.symdef #y $module1_variables-8
...
.symdef #module1_string1 $module1_strings-32
...
.symdef #test $module1_code+0
.symdef #main $module1_code+42
...
```

later

distinction simplifies relocation!

or any other data representation

no memory data means all zeros

size is 80 even if there is more memory data (less means fill rest with zeros)

→ code segments will be located before data segments!

offset relative to: beginning ; end

global variables

string constants

procedures

linkup and Fixup

objectFile = relocatableSegment { relocatableSegment }
{ symdef } { linkup } { fixup } .

linkup = ".linkup" relocatableAddress symbolicReference .
symbolicReference = symbolName .

.linkup $module1_code+20 #x
.linkup $module1_code+24 #y
.linkup $module1_code+110 #module1_string1
.linkup $module2_code+40 #test
...

access of global variables
access of string constants
procedure calls

.symdef #x $module1_variables-4
.symdef #y $module1_variables-8
.symdef #module1_string1 $module1_strings-32
.symdef #test $module1_code+0

linking

fixup = ".fixup" relocatableAddress relocatableAddress

.fixup $module1_code+20 $module1_variables-4
.fixup $module1_code+24 $module1_variables-8
.fixup $module1_code+110 $module1_strings-32
.fixup $module2_code+40 $module1_code+0
...

relocation type omitted for simplicity

Executable

```
executableFile = absoluteSegment start .

absoluteSegment =
 ".code" { memoryData } ".data" { memoryData } .

start = ".start" offset .

.code
  memoryData<$module1_code, 80>
  memoryData<$module2_code, 120>
  ...
.data
  memoryData<$module1_variables, 20>
  memoryData<$module1_strings, 160>
  memoryData<$module2_variables, 8>
  memoryData<$module2_strings, 100>
  ...
.start 42
```

distinction not massary !

$$address(\$i+o) = o + \sum_{1 \le j < i} size(\$j)$$

$$offset(\$i-o) = -o - \sum_{i < j \le \#\text{of data segments}} size(\$j)$$

.symdef #main $module1_code+42

.fixup $module1_code+20 $module1_variables-4
---> set parameter c @ address 20 in code segment to -272
(-272 = -4-160-8-100)

.fixup $module1_code+24 $module1_variables-8
---> set parameter c @ address 24 in code segment to -276
(-276 = -8-160-8-100)

.fixup $module1_code+110 $module1_strings-32
---> set parameter c @ address 110 in code segment to -140
(-140 = -32-8-100)

.fixup $module2_code+40 $module1_code+0
---> set parameter c @ address 120 in code segment to 0
(120 = 40+80)

Loading and go

1. load code segment into memory starting @ address 0
2. load data segment into memory right after code segment
3. set PC to start offset — runtime PC!
4. set GP to size of code segment plus size of data segment
5. set SP to size of memory
6. set HP to GP (heap "bump" pointer) — use HP: reg[26] in malloc implementation
7. push console parameters onto stack
8. set LINK to 0 — termination condition
9. start emulator
10. terminate on RET to 0

Bootstrapping

↝ implement system procedures in object file:

library code

(standard approach when linker is available)

↝ or implement as part of emulator:

virtual machine

(requires special instructions)

↝ or have compiler generate code

(problem: code duplication)