

Oplossingen Tentamen Procesalgebra 28 november 2002

1. i.

$$\begin{aligned}
 &(a.\epsilon + a.b.\epsilon) \parallel c.\epsilon = \\
 &(a.\epsilon + a.b.\epsilon) \parallel c.\epsilon + c.\epsilon \parallel (a.\epsilon + a.b.\epsilon) + (a.\epsilon + a.b.\epsilon) | c.\epsilon = \\
 &a.\epsilon \parallel c.\epsilon + a.b.\epsilon \parallel c.\epsilon + c.(a.\epsilon + a.b.\epsilon) + a.\epsilon | c.\epsilon + a.b.\epsilon | c.\epsilon = \\
 &a.c.\epsilon + a.(b.\epsilon \parallel c.\epsilon) + c.(a.\epsilon + a.b.\epsilon) + \delta + \delta = \\
 &a.c.\epsilon + a.(b.c.\epsilon + c.b.\epsilon) + c.(a.\epsilon + a.b.\epsilon)
 \end{aligned}$$

1. ii.

$$\begin{aligned}
 &\partial_{\{a,b\}}(a.\epsilon \parallel b.\epsilon \parallel b.\epsilon) = \\
 &\partial_{\{a,b\}}(a.\epsilon \parallel (b.\epsilon \parallel b.\epsilon) + b.\epsilon \parallel (a.\epsilon \parallel b.\epsilon) + b.\epsilon \parallel (a.\epsilon \parallel b.\epsilon) + (a.\epsilon | b.\epsilon) \parallel b.\epsilon + \\
 &(a.\epsilon | b.\epsilon) \parallel b.\epsilon + (b.\epsilon | b.\epsilon) \parallel a.\epsilon) = \\
 &\partial_{\{a,b\}}(a.b.b.\epsilon + b.(a.b.\epsilon + b.a.\epsilon) + c.b.\epsilon) = \\
 &\delta + \delta + c.\delta = \\
 &c.\delta
 \end{aligned}$$

1. iii.

$$\begin{aligned}
 &\tau_a(a.b.\epsilon \parallel a.b.\epsilon) = \\
 &\tau_a(a.b.\epsilon \parallel a.b.\epsilon + a.b.\epsilon \parallel a.b.\epsilon + a.b.\epsilon | a.b.\epsilon) = \\
 &\tau_a(a.(b.\epsilon \parallel a.b.\epsilon) + \delta) = \\
 &\tau_a(a.(b.\epsilon \parallel a.b.\epsilon + a.b.\epsilon \parallel b.\epsilon + b.\epsilon | a.b.\epsilon)) = \\
 &\tau_a(a.(b.a.b.\epsilon + a.(b.\epsilon \parallel b.\epsilon) + c.(b.\epsilon))) = \\
 &\tau_a(a.(b.a.b.\epsilon + a.b.b.\epsilon + c.b.\epsilon)) = \\
 &\tau.(b.\tau.b.\epsilon + \tau.b.b.\epsilon + c.b.\epsilon) = \\
 &\tau.(b.b.\epsilon + \tau.b.b.\epsilon + c.b.\epsilon)
 \end{aligned}$$

1. iv.

$$\begin{aligned}
 &\pi_2(a.b.\delta \parallel \tau.\delta) = \\
 &\pi_2(a.b.\delta \parallel \tau.\delta + \tau.\delta \parallel a.b.\delta + a.b.\delta | \tau.\delta) = \\
 &\pi_2(a.(b.\delta \parallel \tau.\delta) + \tau.a.b.\delta + \delta) = \\
 &a.\pi_1(b.\delta \parallel \tau.\delta) + \tau.\pi_2(a.b.\delta) = \\
 &a.\pi_1(b.\tau.\delta + \tau.b.\delta) + \tau.a.b.\delta = \\
 &a.(b.\pi_0(\tau.\delta) + \tau.\pi_1(b.\delta)) + \tau.a.b.\delta = \\
 &a.(b.\tau.\delta + \tau.b.\delta) + \tau.a.b.\delta = \\
 &a.(\tau.(b.\delta + \delta) + b.\delta) + \tau.a.b.\delta = \\
 &a.b.\delta + \tau.a.b.\delta
 \end{aligned}$$

2.i.

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2.ii

u

2.iii.

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3.i.

Omdat er voor alle termen x, y basistermen x', y' zijn met $x = x', y = y'$, is het voldoende om de eigenschap aan te tonen voor basistermen. We bewijzen de eigenschap met inductie naar de structuur van x .

Basis

$x \equiv \epsilon$. Dan $\epsilon|\tau.y = \delta$.

$x \equiv \delta$. Dan $\delta|\tau.y = \delta$.

IH $x'|\tau.y = \delta, x''|\tau.y = \delta$ voor alle y .

$x = a.x'$: $a.x'|\tau.y = \delta$

$x = x' + x''$: $(x' + x'')|\tau.y = x'|\tau.y + x''|\tau.y = \delta + \delta = \delta$. □

3.ii.

Tegenvoorbeeld: $x = \epsilon, y = \epsilon$. Dan $x \ll \tau.y = \epsilon \ll \tau.\epsilon = \delta$, maar $x \ll y = \epsilon \ll \epsilon = \epsilon$.

3.iii.

$\tau.(\tau.x||y) =$

$\tau.\tau.x \ll y =$

$\tau.(\tau.(x + \delta) + \delta) \ll y =$

$\tau.(x + \delta) \ll y =$

$\tau.(x||y)$.

3.iv.

$\tau.x||\tau.y =$

$\tau.x \ll \tau.y + \tau.y \ll \tau.x + \tau.x|\tau.y =$

$\tau.(x||\tau.y) + \tau.(y||\tau.x) + \delta =$

$\tau.(\tau.y||x) + \tau.(\tau.x||y) =$

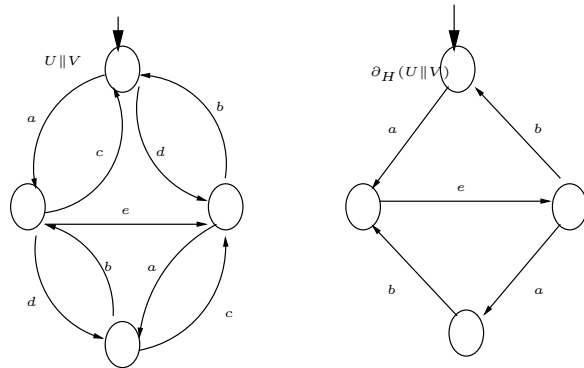
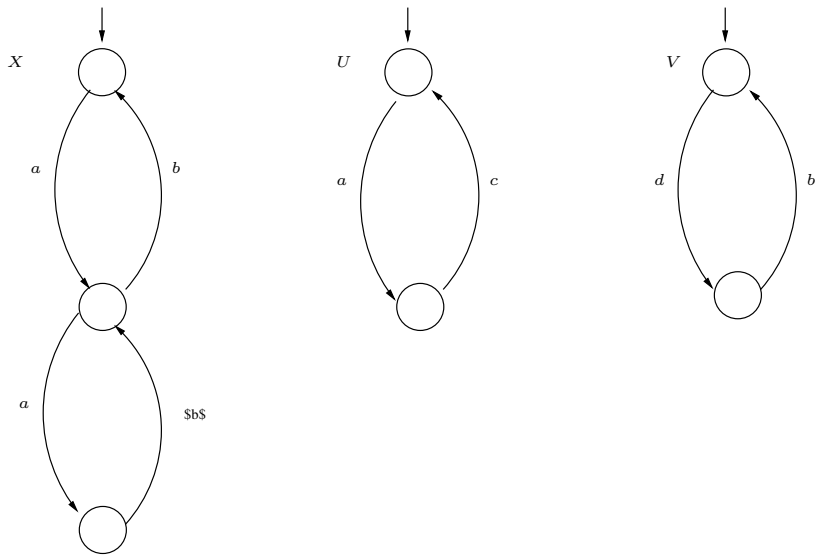
$$\begin{aligned} \tau.(y\|x) + \tau.(x\|y) &= \\ \tau.(x\|y) & \end{aligned}$$

4.i.

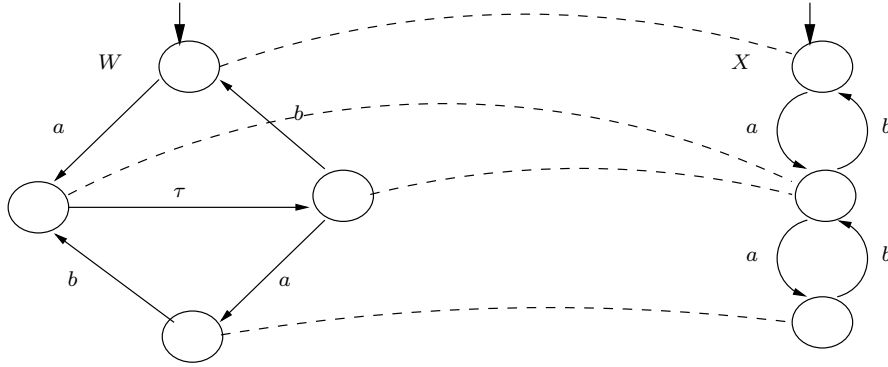
$$\begin{aligned} U\|V &= a.(U'\|V) + d.(U\|V') \\ U'\|V &= c.(U\|V) + d.(U'\|V') + e.(U\|V') \\ U\|V' &= a.(U'\|V') + b.(U\|V) \\ U'\|V' &= c.(U\|V') + b.(U'\|V) \end{aligned}$$

$$\begin{aligned} \partial_H(U\|V) &= a.\partial_H(U'\|V) \\ \partial_H(U'\|V) &= e.\partial_H(U\|V') \\ \partial_H(U\|V') &= a.\partial_H(U'\|V') + b.\partial_H(U\|V) \\ \partial_H(U'\|V') &= b.\partial_H(U\|V') \end{aligned}$$

Dus de procesgrafen voor $X, U, V, U\|V, \partial_H(U\|V)$ zijn:



4.ii.
De procesgrafieën voor X en W zijn rb-bisimilaar.



4.iii.

Van 4.i. hebben we dat

$$W = a.W'$$

$$W' = \tau.W''$$

$$W'' = a.W''' + b.W$$

$$W''' = b.W'$$

of

$$W = a.\tau.W'' = a.W''$$

$$W'' = a.W''' + b.W$$

$$W''' = b.W'$$

en het is nu duidelijk dat $X = W$.

4.iv.

We bewijzen dat $\pi_n(X) = X_n, \pi_n(X') = X'_n$ en $\pi_n(X'') = X''_n$ voor alle $n \geq 1$ met inductie naar n .

$$\mathbf{Basis} \quad \pi_1(X) = a.\delta = X_1, \pi_1(X') = a.\delta + b.\delta = X'_1, \pi_1(X'') = b.\delta = X''_1$$

$$\mathbf{IH} \quad \pi_n(X) = X_n, \pi_n(X') = X'_n, \pi_n(X'') = X''_n.$$

$$\text{Dan } \pi_{n+1}(X) = a.\pi_n(X'') + b.\pi_n(X) = a.X''_n + b.X_n = X_{n+1}.$$

$$\pi_{n+1}(X') = a.\pi_n(X'') + b.\pi_n(X) = a.X''_n + b.X_n = X'_{n+1}.$$

$$\pi_{n+1}(X'') = b.\pi_n(X') = b.X'_n = X''_{n+1}. \quad \square$$

5.i.

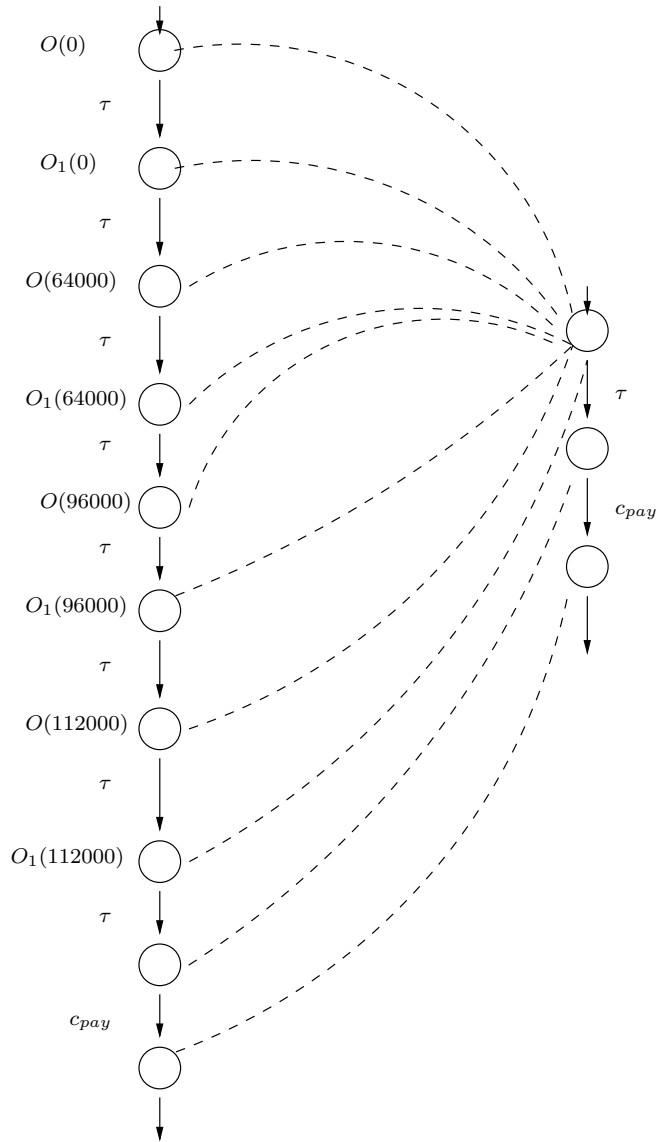
$$X(b) = \partial_H(P_1(b) \| P_2) = c_{offer}.\partial_H(Q_1(b) \| Q_2(b)) = c_{offer}.X_1(b)$$

$$X_1(b) = c_{reject}.\partial_H(P_1((b+T+1)div2) \| P_2) = X_1(b) =$$

$$= \begin{cases} c_{reject}.X((b+T+1)div2), & \text{als } 0 \leq b < B \\ c_{accept}.c_{pay}.\epsilon, & \text{als } b \geq B \end{cases}$$

5.ii.

$O(b) = \tau.\tau_I(\partial_H(Q_1(b)\|Q_2(b))) = \tau.O_1(b)$
 $O_1(b) = \tau.\tau_I(\partial_H(P_1((b+T+1)div2)\|P_2)) = \tau.O((b+T+1)div2), 0 \leq b < B$
 $O_1(b) = \tau.c_{pay}.\epsilon, b \geq B$
 Voor $T = 128.000, B = 112.000$ is de procesgraaf van $O(0)$ de volgende:



Het is rb-bisimilaair met de graaf voor $\tau.c_{pay}.\epsilon$.

5.iii. Als $T < B$ vindt er geen verkoop plaats.