Formale Systeme Proseminar

Tasks for Week 15, 25.1.2018

 ${\bf Task} \ {\bf 1} \ {\bf Construct} \ {\bf a} \ {\bf FA} \ {\bf for} \ {\bf the} \ {\bf language}$

 $L = \{ w \in \{a, b\}^* \mid w \text{ has at least three } a \text{'s or at least two } b \text{'s} \}.$

Note that this language is a union of two languages.

Task 2 Construct a FA for the language

$$L = \{w_1 w_2 \in \{0,1\}^* \mid w_1 = 0^{2n}, w_2 = (100)^m, \text{ for some } n, m \in \mathbb{N}\}.$$

Note that the regular expression for L is $(00)^* \cdot (100)^*$.

Task 3 Let L be the language of all strings over $\{0, 1\}$ that do not contain a pair of 1's that are separated by an odd number of symbols. Give the state diagram of a DFA with 5 states that recognises L.

Task 4 Construct a DFA for the language L^* where

 $L = 01 \cup (00)^* 11.$

- **Task 5** Construct a DFA for the language $L_1 \cdot L_2$ where $L_1 = \{a, b\}^*$ and $L_2 = \{aabab\}.$
- **Task 6** Let L be a regular language, $L \subseteq \Sigma^*$. Show that the reversed language of L defined as

 $L^R = \{ w \in \Sigma^* \mid w^R \in L \}$

where reversed words are defined inductively by

$$\varepsilon^R = \varepsilon, (ua)^R = au^R \text{ for } a \in \Sigma, u \in \Sigma^*$$

is regular as well.

Hint: From an automaton for L, construct an automaton for L^R .

- **Task 7** Give state diagrams of NFAs with the specified number of states recognising each of the following languages. In all parts the alphabet is $\{0, 1\}$ and the language is given via its regular expression.
 - (a) The language 0 with two states.

- (b) The language 0^* with one state.
- (c) The language $(0 \cup 1)^* 00$ with three states,
- (d) The language $1^* \cdot (001^+)^*$ with three states.

All the best for the rest of your studies!