

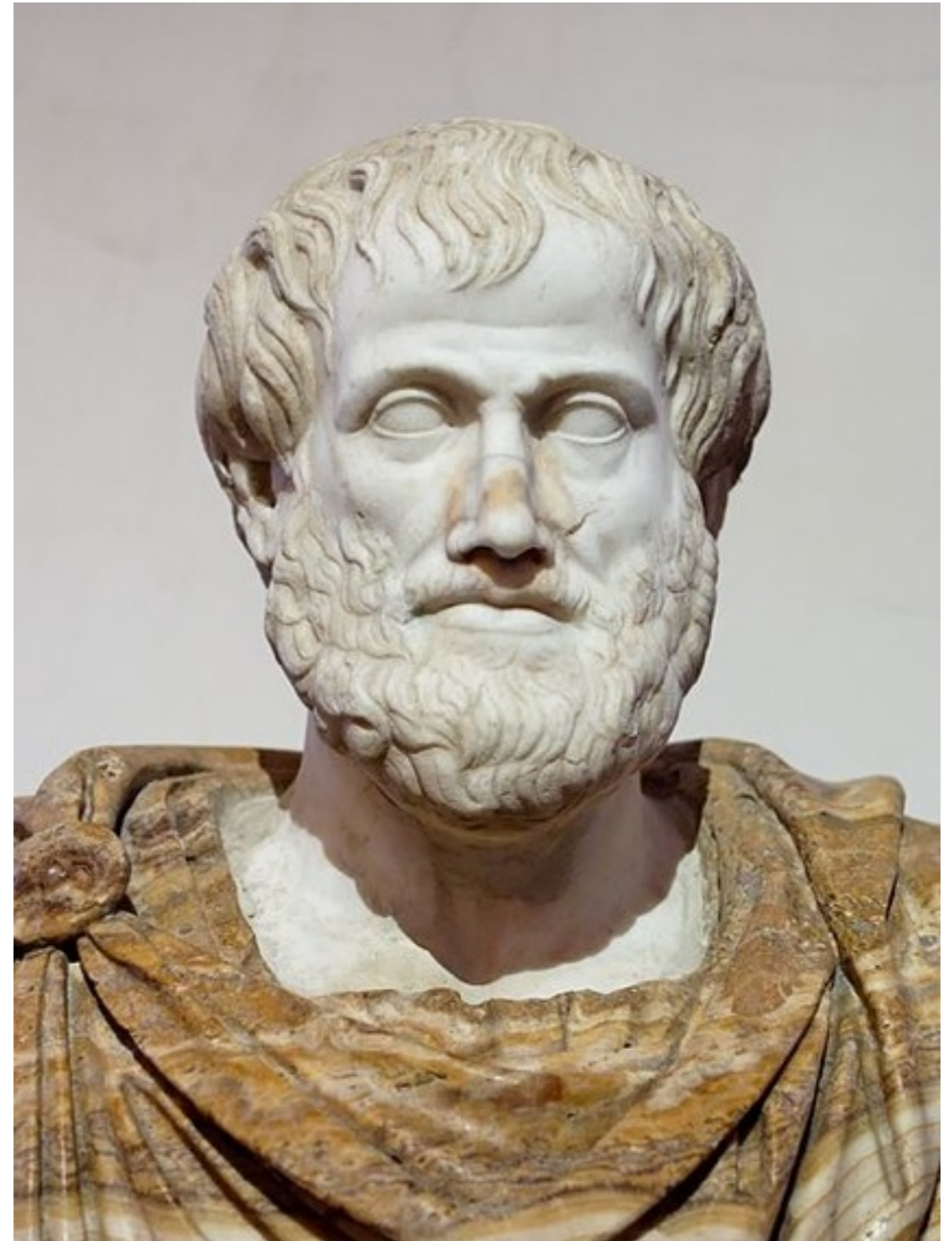
Logic

In the beginning

Aristotle +/- 350 B.C.

Organon

19 syllogisms



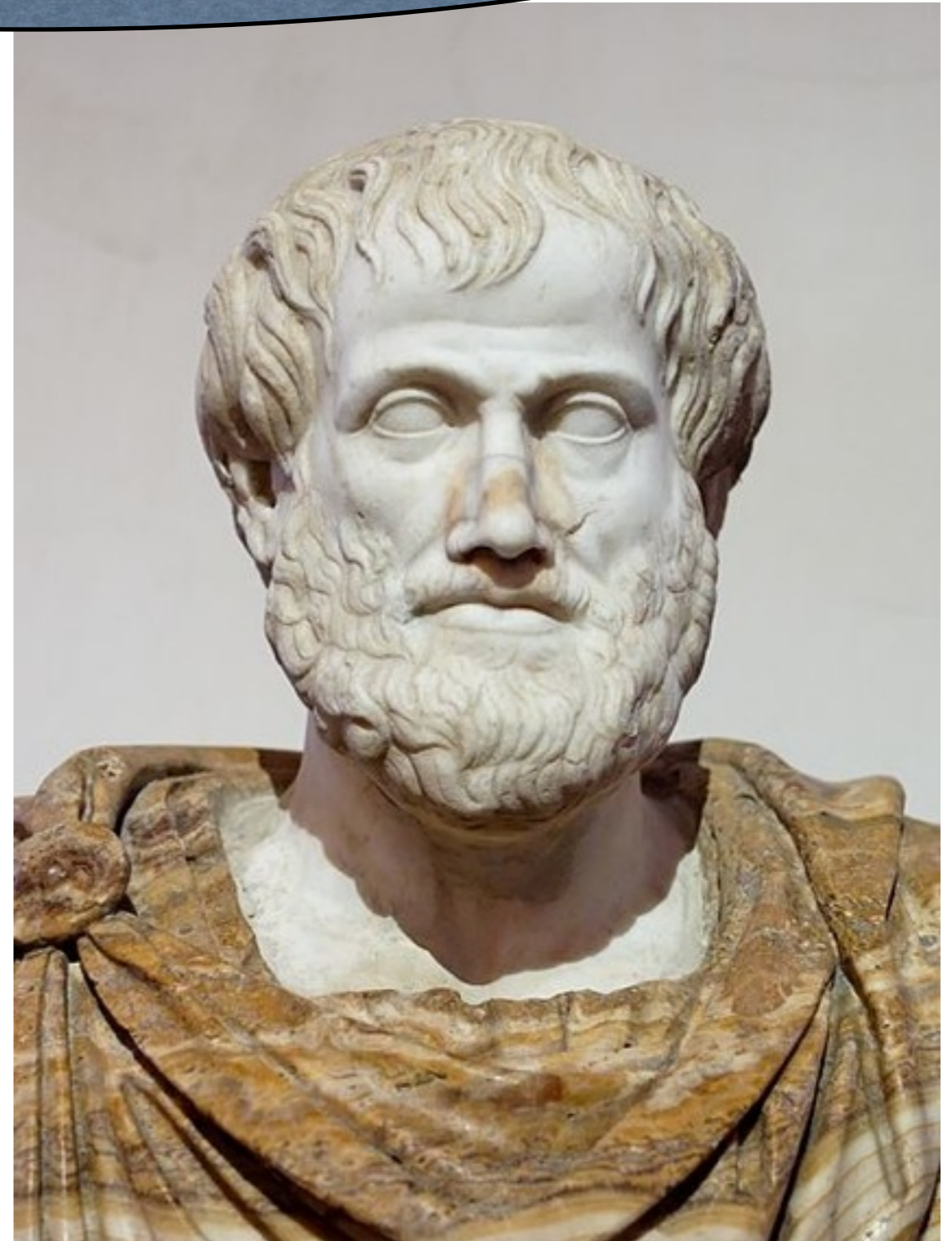
Logic = study of correct reasoning

In the beginning

Aristotle +/- 350 B.C.

Organon

19 syllogisms



Barbara syllogism

All K's are L's
All L's are M's

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Barbara syllogism

only later called so,
in the Middle Ages

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Logic (Logos, Greek for word, understanding, reason) deals with general reasoning laws in the form of formulas with parameters.

Propositions

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Def. A proposition (**Aussage**) is a grammatically correct sentence that is either true or false.

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logic deals with patterns!
what matters are not particular
propositions but the way in
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from them

Propositions

Def. A proposition (**Aussage**) is a grammatically correct sentence that is either true or false.

Connectives

- \wedge for “and”
- \vee for “or”
- \neg for “not”
- \Rightarrow for “if .. then” or “implies”
- \Leftrightarrow for “if and only if”

logic deals with patterns!
what matters are not particular propositions but the way in which (abstract) propositions are combined and what follows from them

Abstract propositions

Abstract propositions

Definition

- Basis** Propositional variables are abstract propositions.
- Step (Case 1)** If P is an abstract proposition, then so is $(\neg P)$.
- Step (Case 2)** If P and Q are abstract propositions, then so are $(P \wedge Q)$, $(P \vee Q)$, $(P \Rightarrow Q)$, and $(P \Leftrightarrow Q)$.

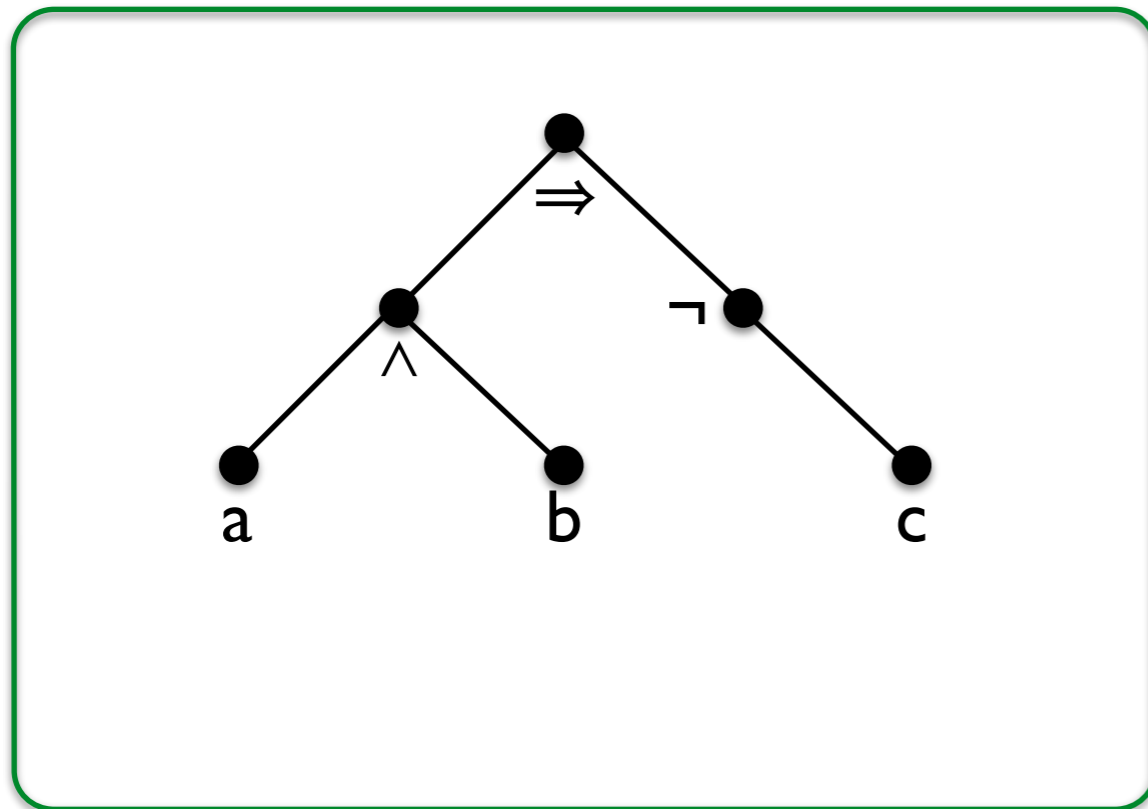
Abstract propositions

Definition

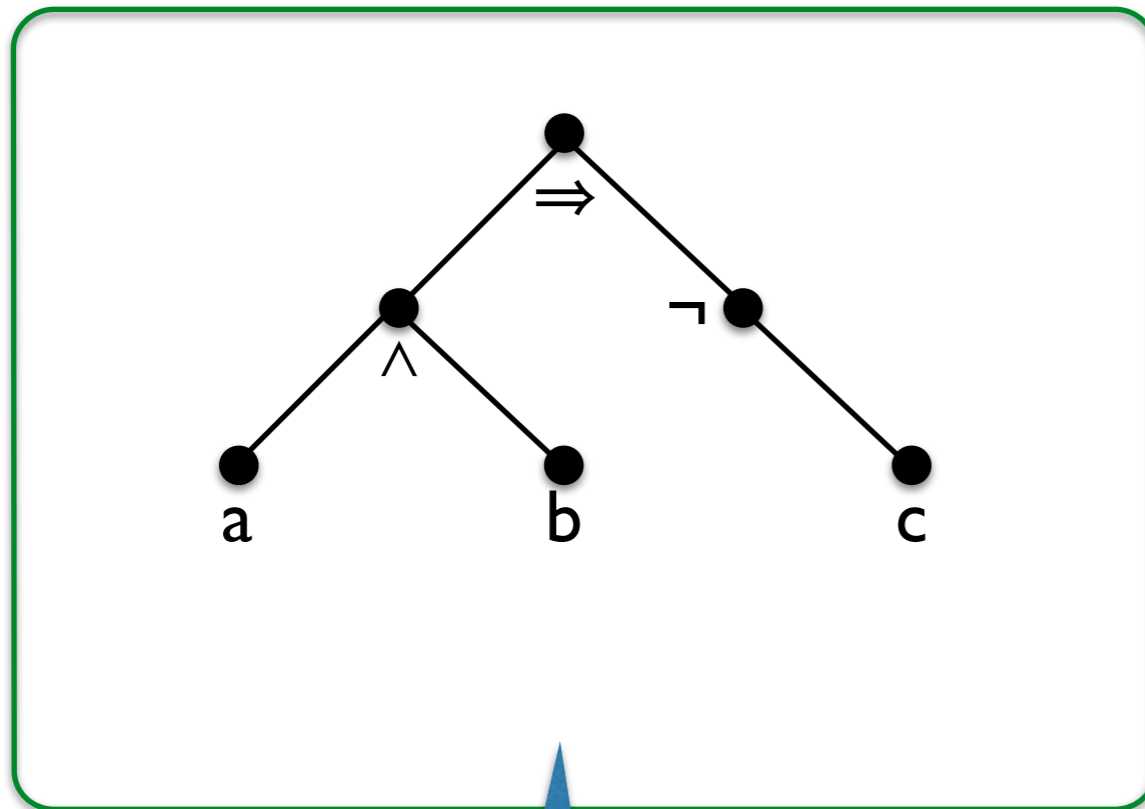
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a recursive/inductive
definition

...and their structure

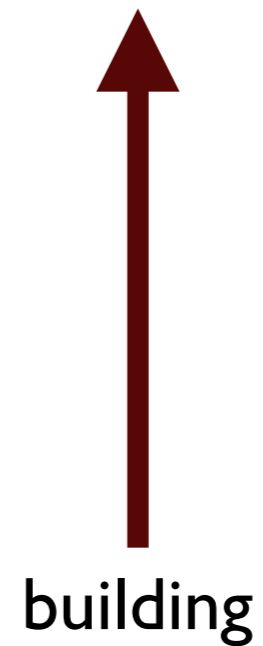
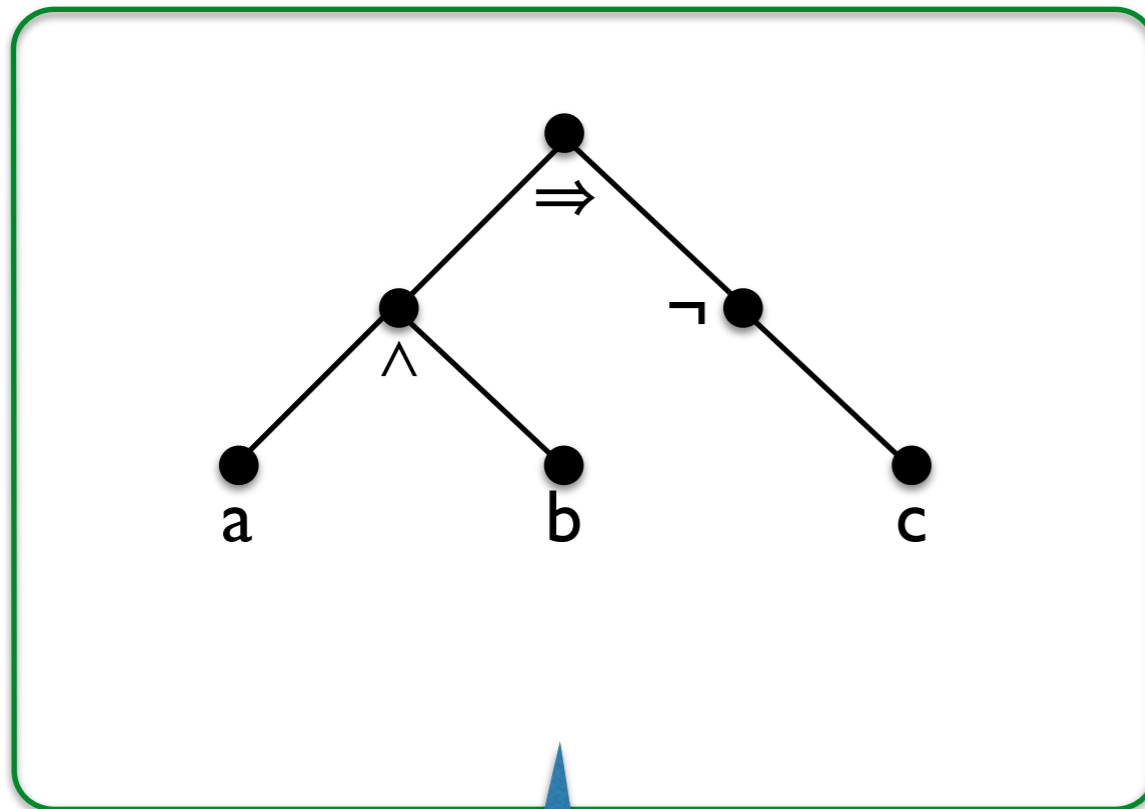


...and their structure



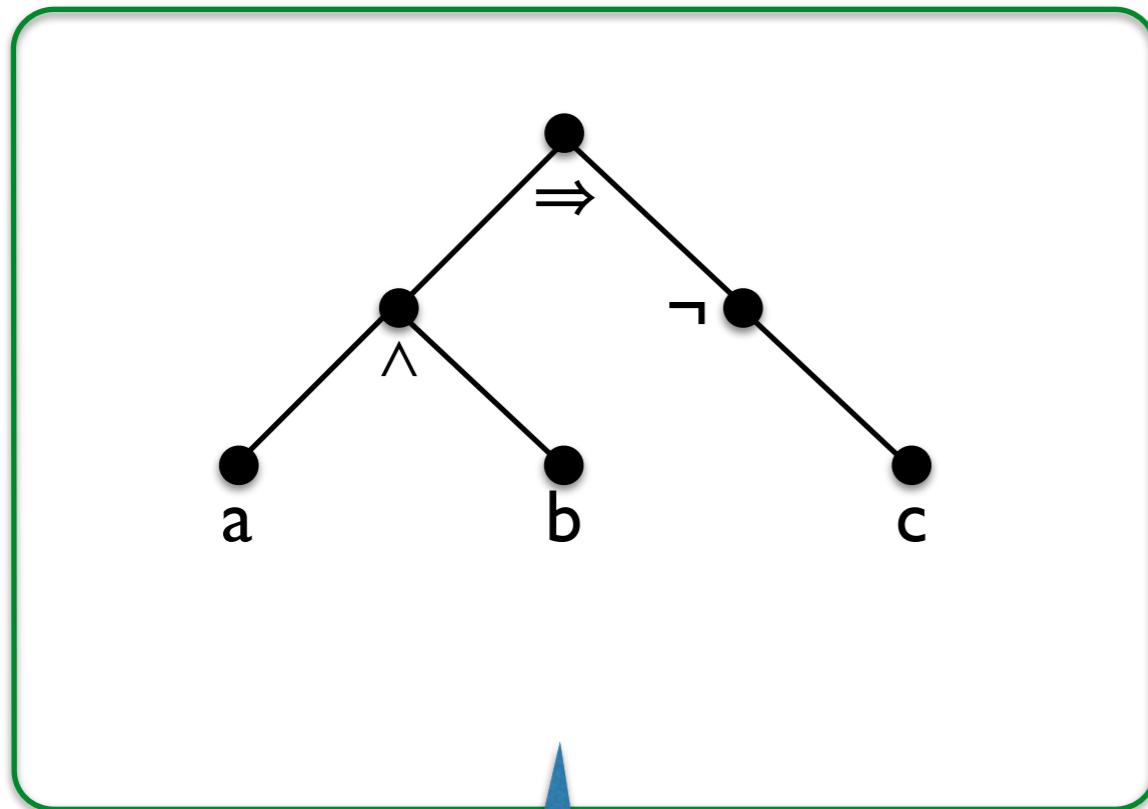
the tree of
 $((a \wedge b) \Rightarrow (\neg c))$

...and their structure

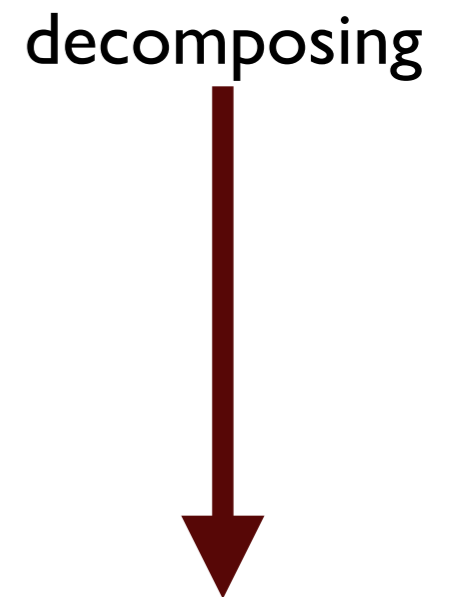


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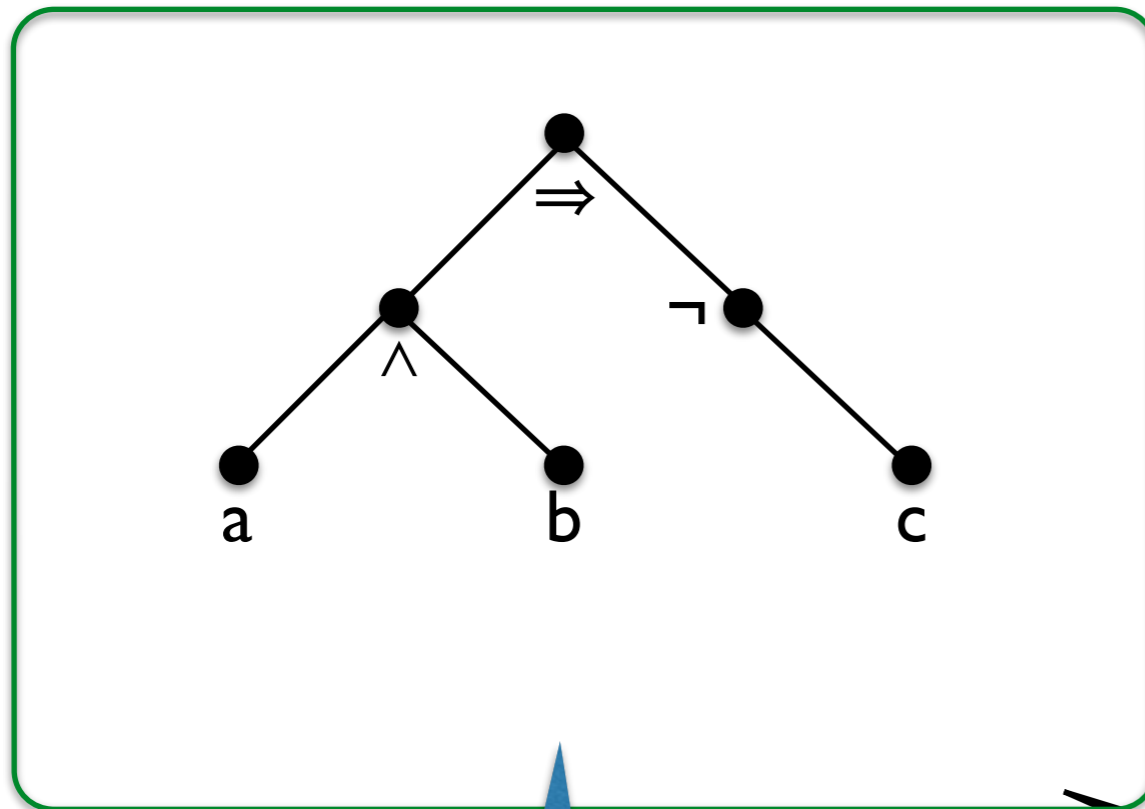
...and their structure



the tree of
 $((a \wedge b) \Rightarrow (\neg c))$



...and their structure



↑
building

decomposing
↓

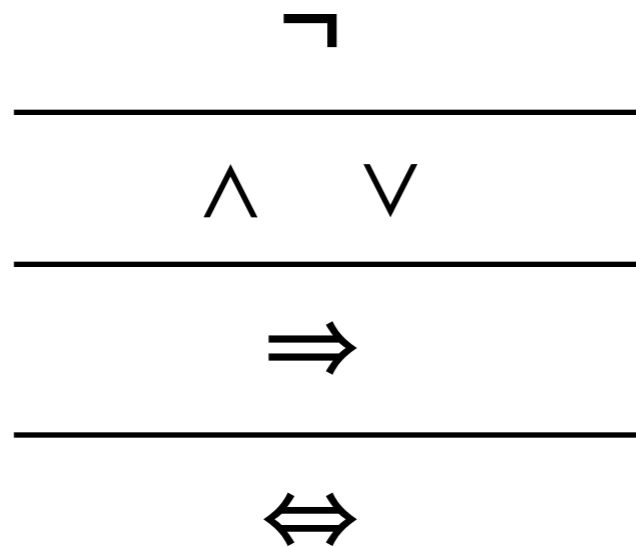
the tree of
 $((a \wedge b) \Rightarrow (\neg c))$

tree representation
(no need of
parenthesis)

Dropping parenthesis

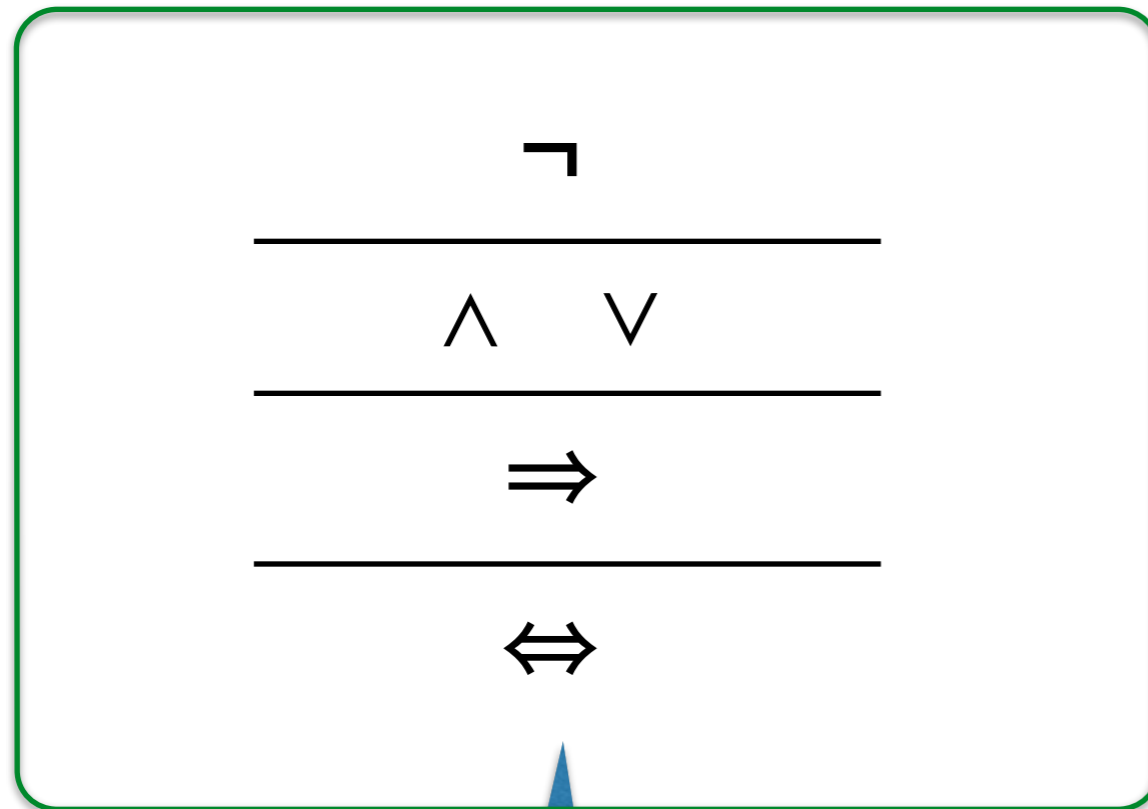
$$\frac{\neg}{\wedge \vee} \Rightarrow \Leftrightarrow$$

Dropping parenthesis

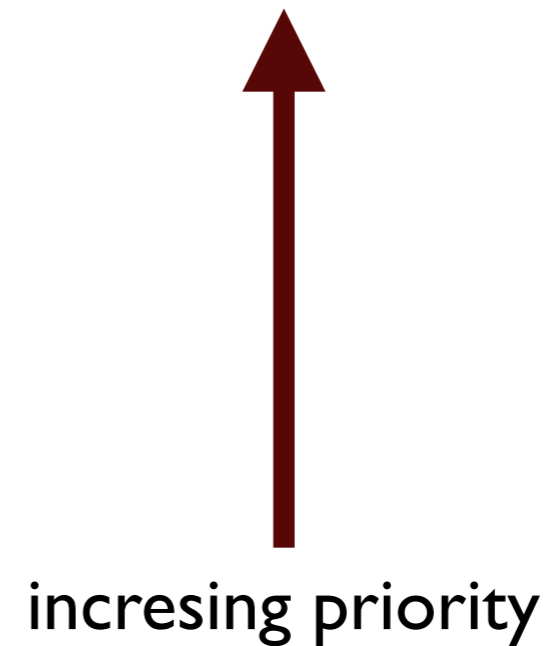


priority schema
(top binds the most)

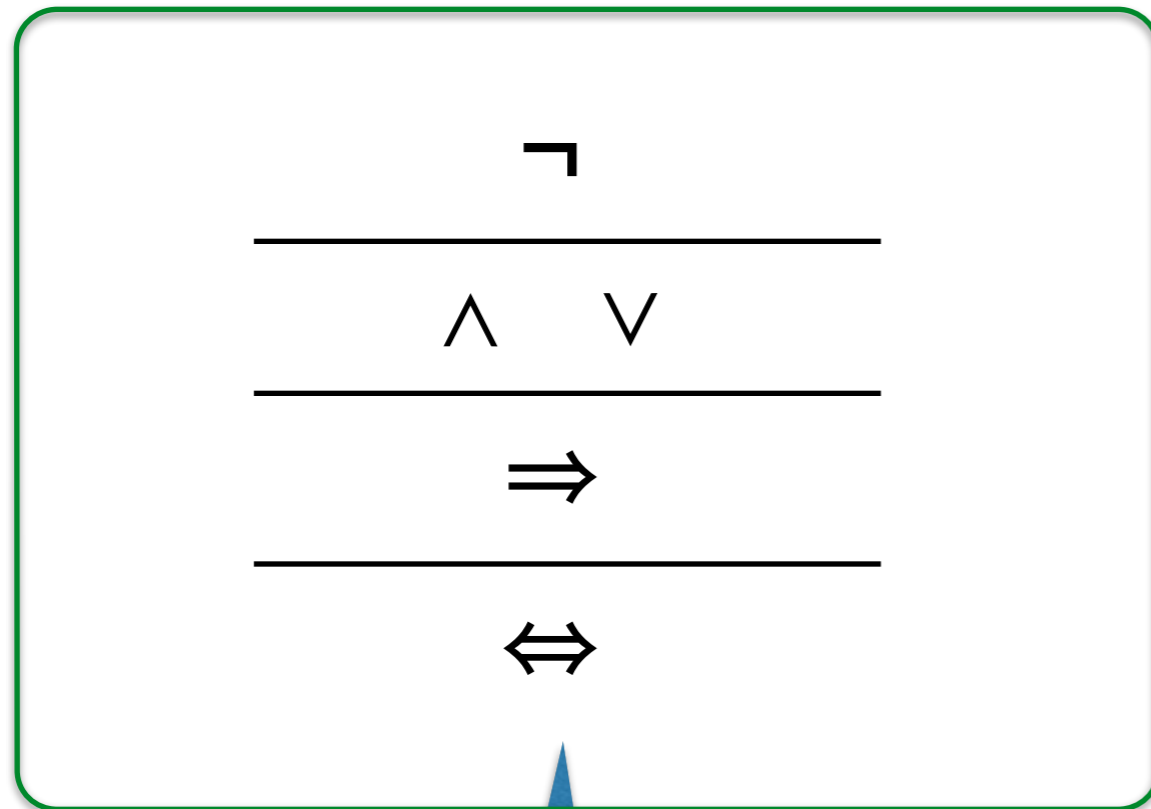
Dropping parenthesis



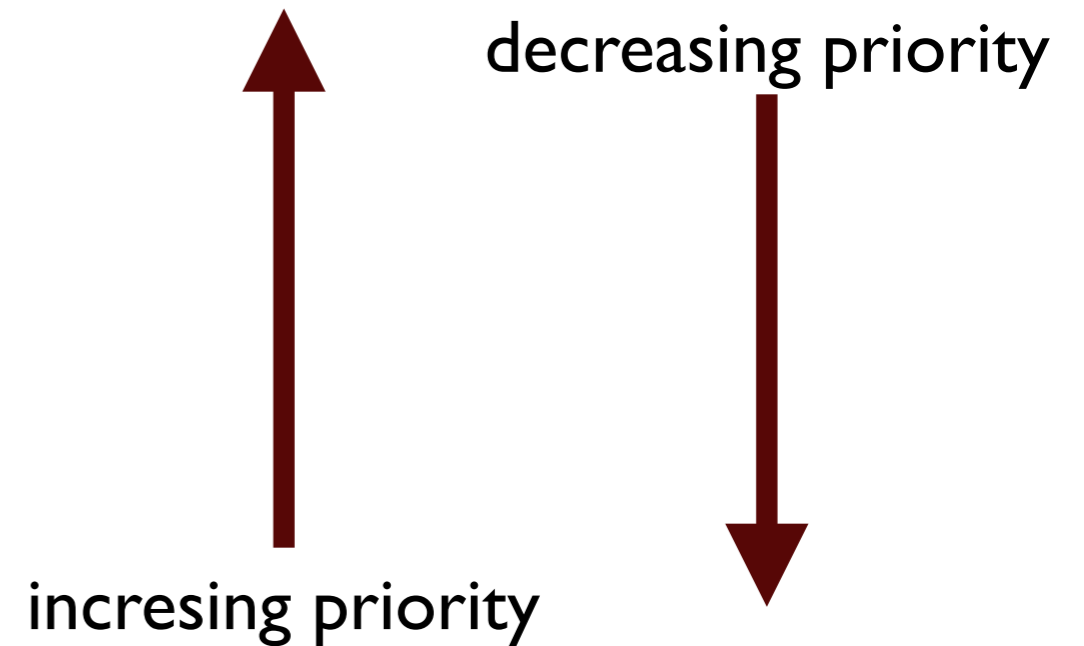
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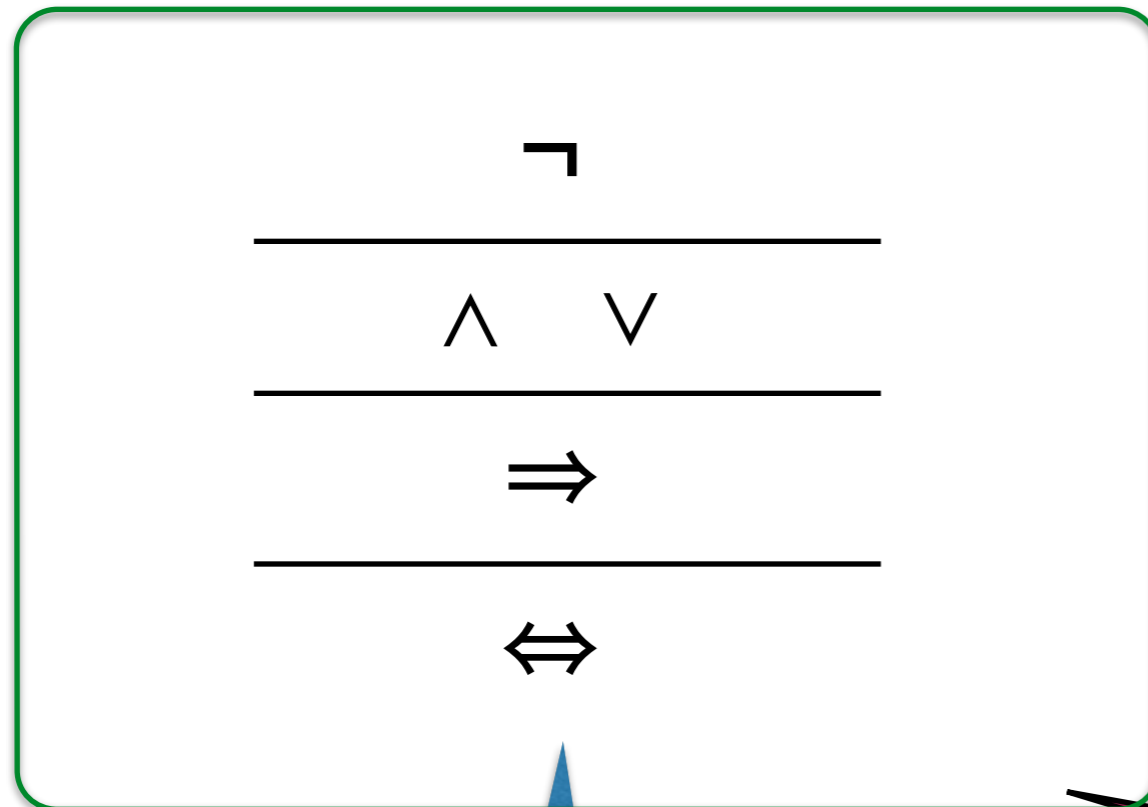
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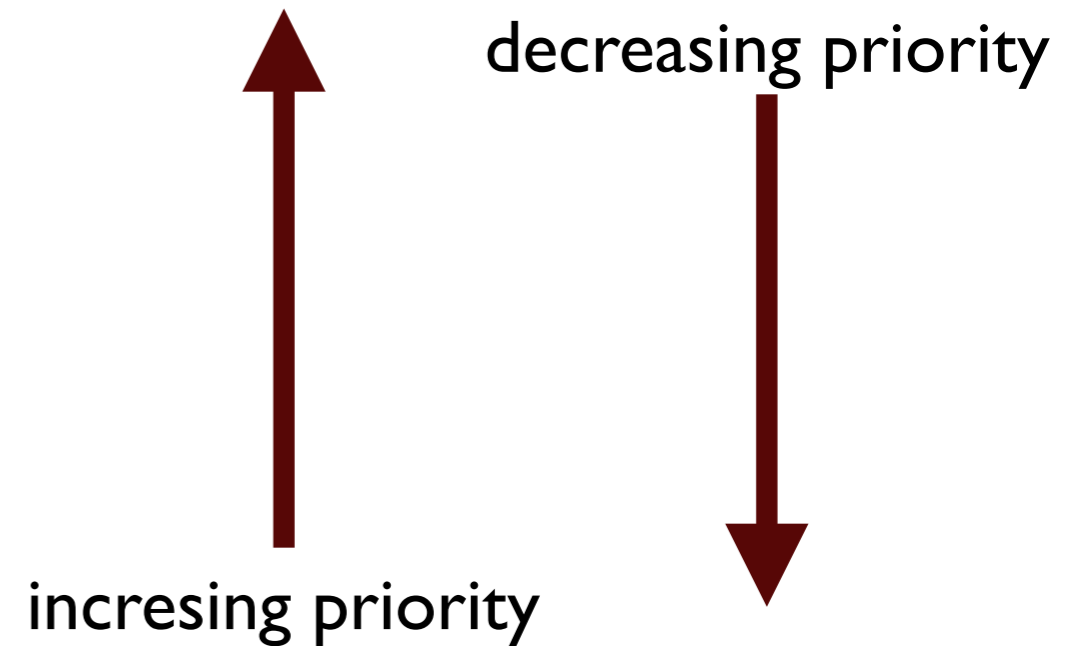
priority schema
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Dropping parenthesis



priority schema
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Example: $((a \wedge b) \Rightarrow (\neg c))$
becomes
 $a \wedge b \Rightarrow \neg c$

Truth tables

Conjunction

P	Q	$P \wedge Q$
0	0	0
0	1	0
1	0	0
1	1	1

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only true when both
P and Q are true

Truth tables

Disjunction

P	Q	$P \vee Q$
0	0	0
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1	0	1
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Truth tables

Disjunction

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true when either P
or Q or both are
true

Truth tables

Negation

Truth tables

Negation

unary connective

Truth tables

Negation

unary connective

P	$\neg P$
0	1
1	0

Truth tables

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true when P
is false

Truth tables

Implication

Truth tables

Implication



needs more attention

Truth tables

Implication

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P	Q	$P \Rightarrow Q$
0	0	1
0	1	1
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Truth tables

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needs more attention

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Implication

needs more attention

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1	0	0
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only false when P is true and Q is false

Truth tables

Bi-implication

Truth tables

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$$P \Leftrightarrow Q$$

is $(P \Rightarrow Q) \wedge (Q \Rightarrow P)$

Truth tables

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true when P and Q have the same truth value

Truth-functions

Def. A truth-function or Boolean function is a function

$$f: \{0, 1\}^n \longrightarrow \{0, 1\}$$

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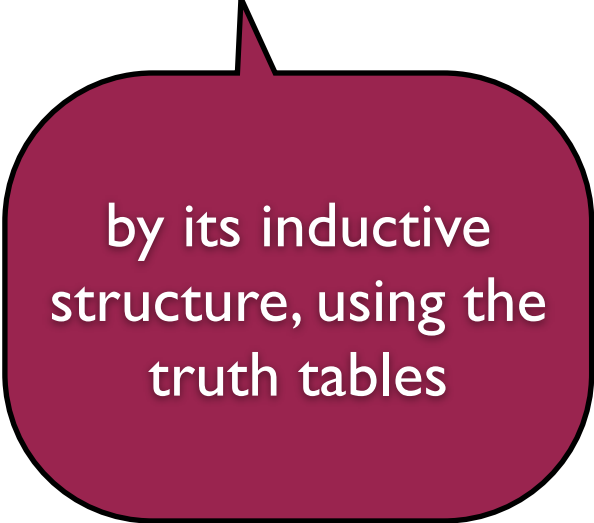
Property: Every abstract proposition $P(a_1, \dots, a_n)$ induces a truth-function.

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by its inductive structure, using the truth tables

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Notation in the book...

$$\left\{ \begin{array}{l} a, b \\ (0, 0) \longmapsto 0 \\ (0, 1) \longmapsto 1 \\ (1, 0) \longmapsto 0 \\ (1, 1) \longmapsto 1 \end{array} \right.$$

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a_1, \dots, a_n are the variables in P (and more) ordered in a sequence

Property: Every abstract proposition $P(a_1, \dots, a_n)$ induces a truth-function.

by its inductive structure, using the truth tables

Notation in the book...

{	a, b
	$(0, 0) \longmapsto 0$
	$(0, 1) \longmapsto 1$
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	$(1, 1) \longmapsto 1$

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Truth-functions

a_1, \dots, a_n are the variables in P (and more) ordered in a sequence

Property: Every abstract proposition $P(a_1, \dots, a_n)$ with ordered and specified variables induces a truth-function.

Note:

The sequence of specified variables matters!

$P(a,b,c): (a \wedge b) \vee b$

induces

a, b, c

$(0,0,0)$	\mapsto	0
$(0,0,1)$	\mapsto	0
$(0,1,0)$	\mapsto	1
$(0,1,1)$	\mapsto	1
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Equivalence of propositions

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on any sequence containing their common variables

Property: The relation $\stackrel{\text{val}}{=}$ is an equivalence on the set of all abstract propositions.

Example

Are the following equivalent? $b \wedge \neg b$ and $c \wedge \neg c$

b	c	$\neg b$	$\neg c$	$b \wedge \neg b$	$c \wedge \neg c$
0	0				
0	1				
1	0				
1	1				

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1	0	0	1	0	0
1	1	0	0	0	0

Their truth values are the same, so they are equivalent

$$b \wedge \neg b \stackrel{val}{=} c \wedge \neg c$$