

Coalgebra

Love and Beauty in Science

Ana Sokolova



Do you know any
coalgebra?

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coalgebra?

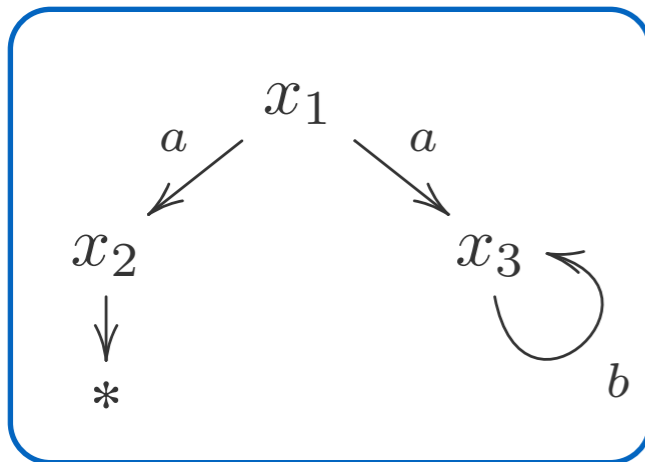
Yes, you know many coalgebras !

Some coalgebras

Some coalgebras

NFA

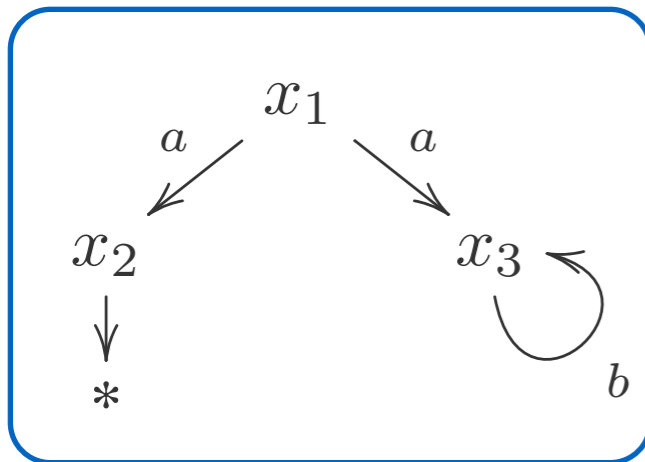
$$X \rightarrow 2 \times (\mathcal{P}X)^A$$



Some coalgebras

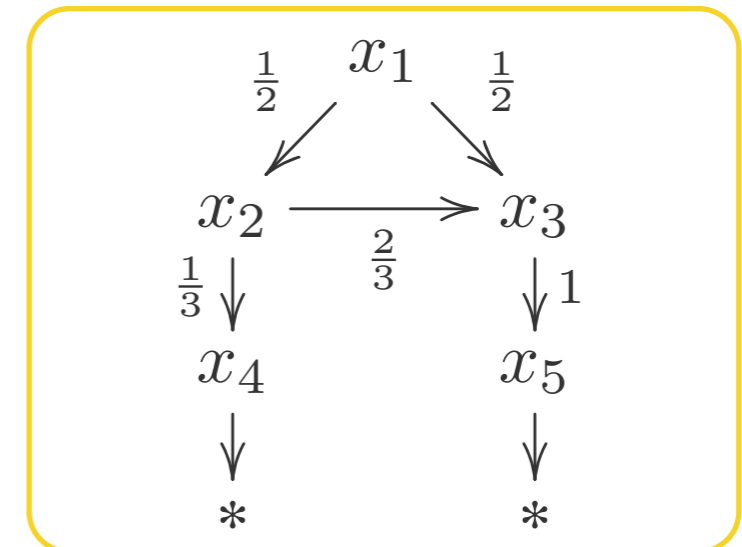
NFA

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MC

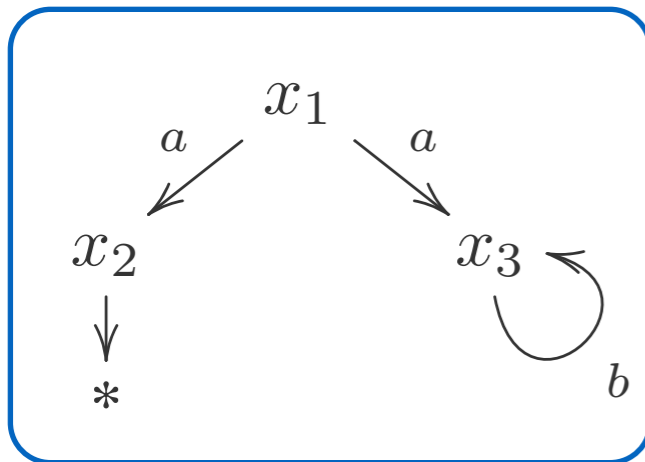
$$X \rightarrow \mathcal{D}X + I$$



Some coalgebras

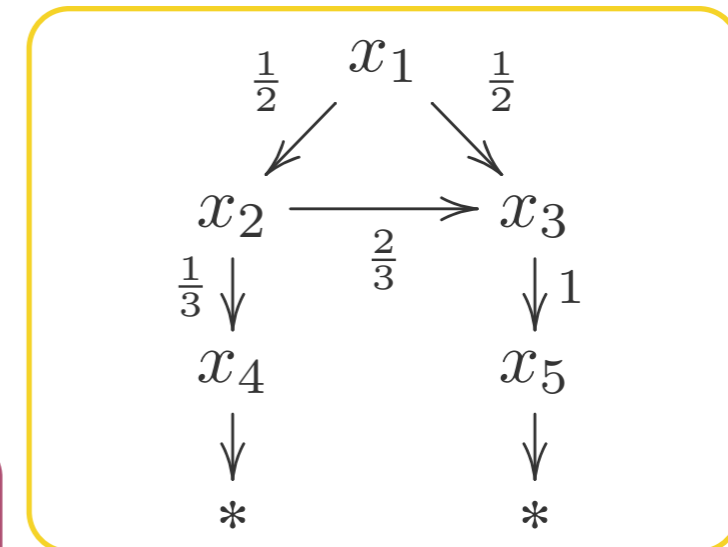
NFA

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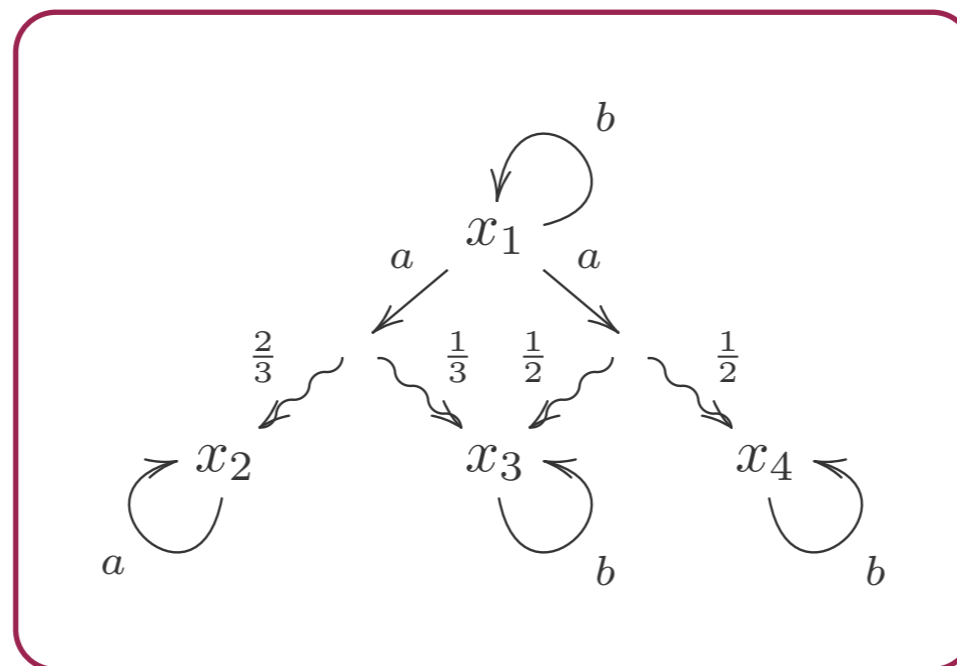
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PA

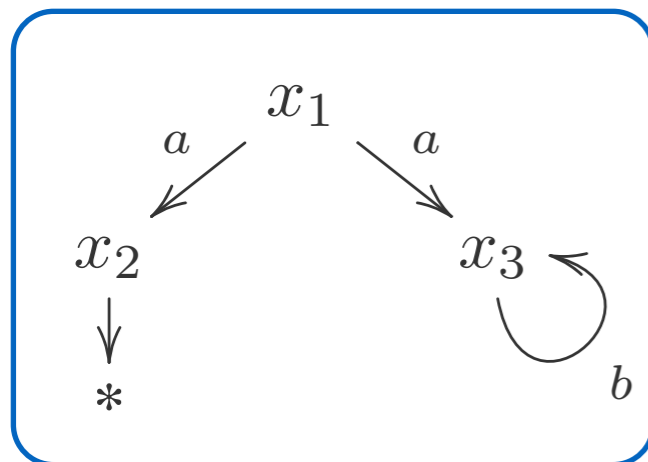
$$X \rightarrow (\mathcal{P}\mathcal{D}X)^A$$



Some coalgebras

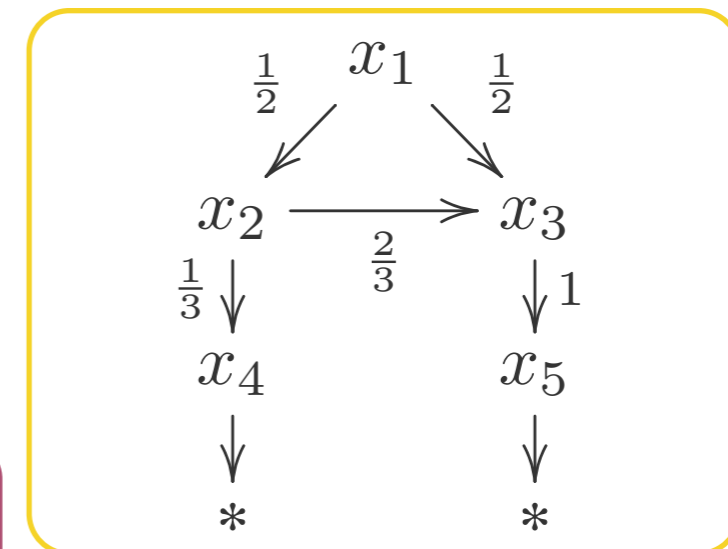
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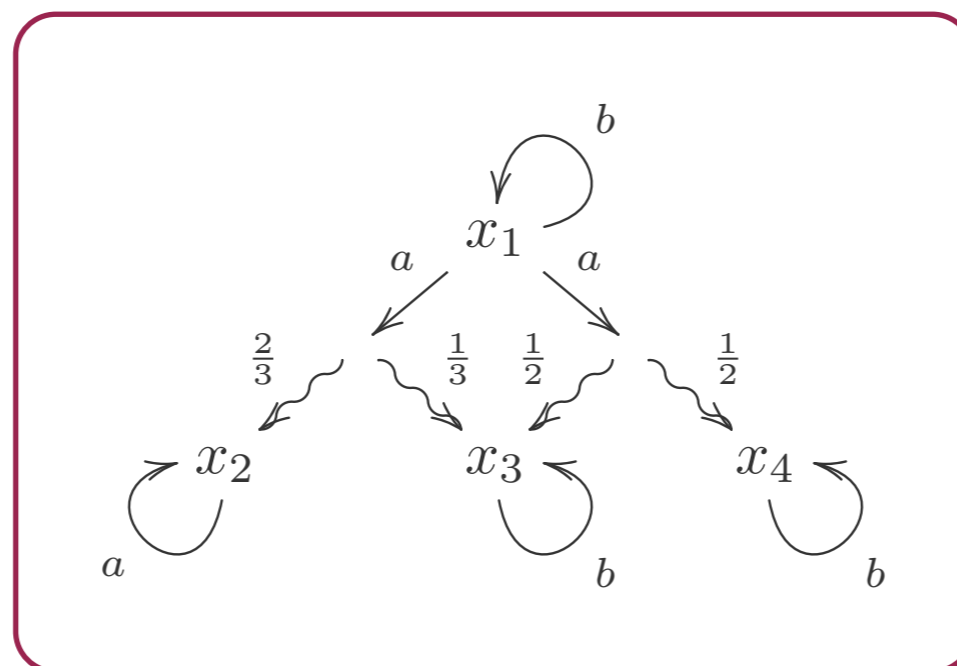
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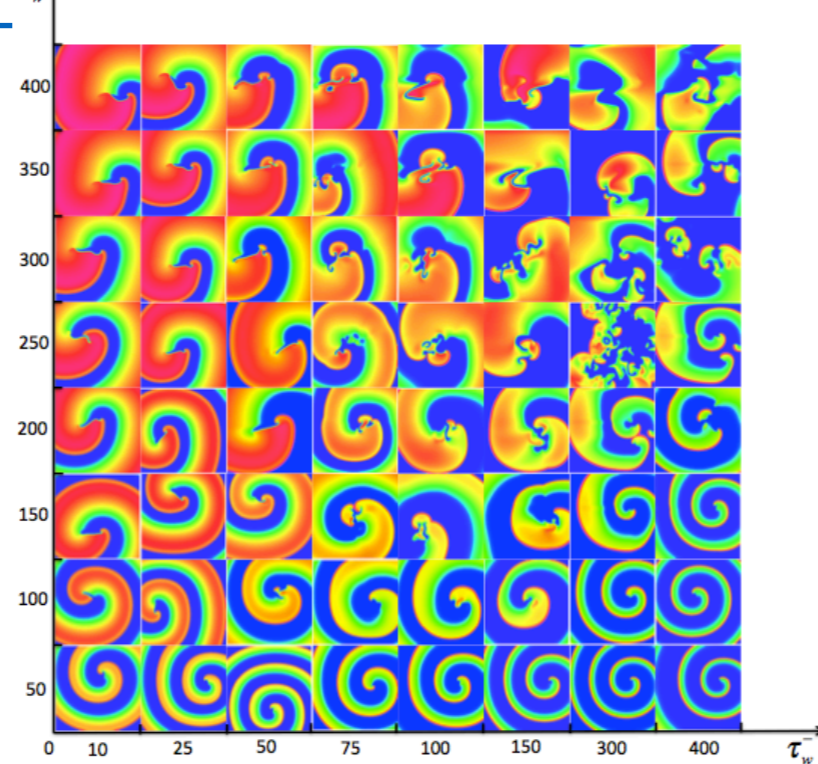
Various transitions systems / automata are coalgebras

Where do they live ?

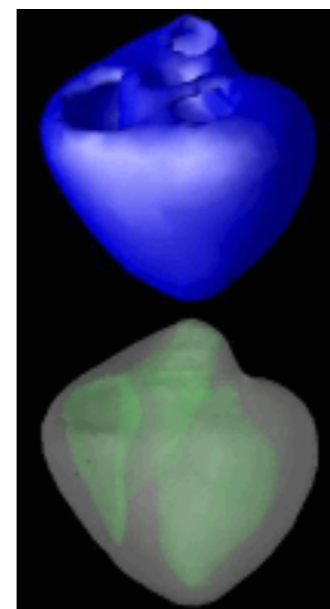
Where do they live ?



τ_w^+ 2D Tissue Model



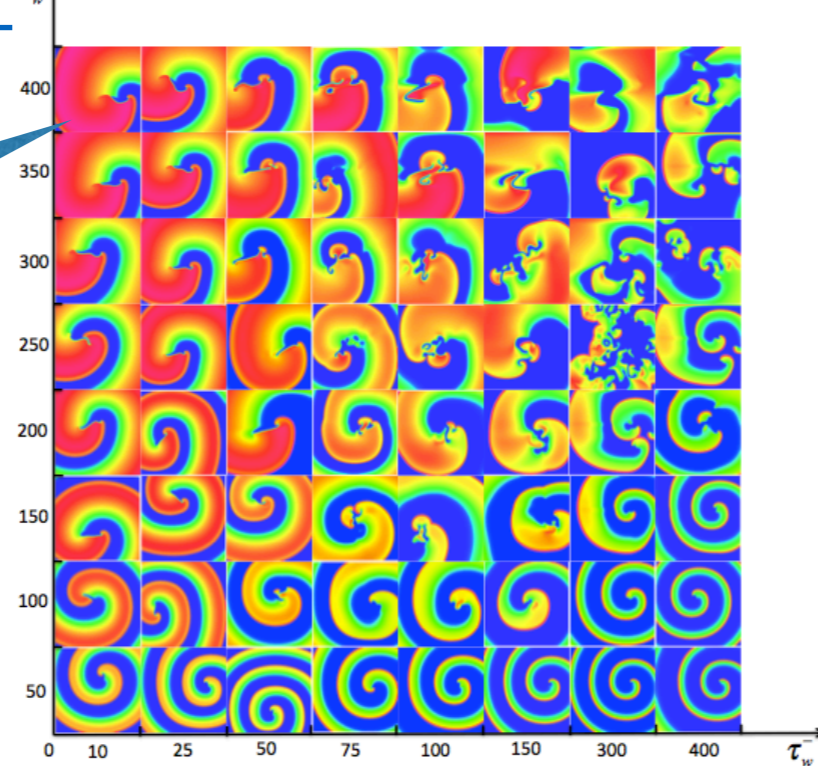
3D Organ Model



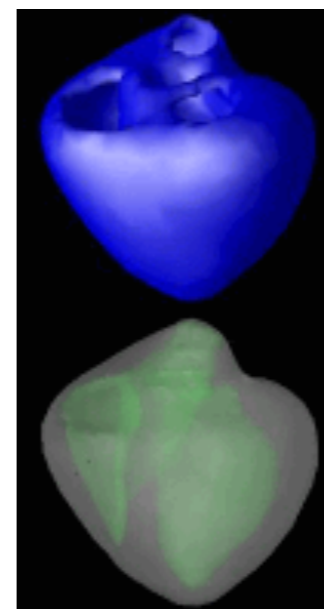
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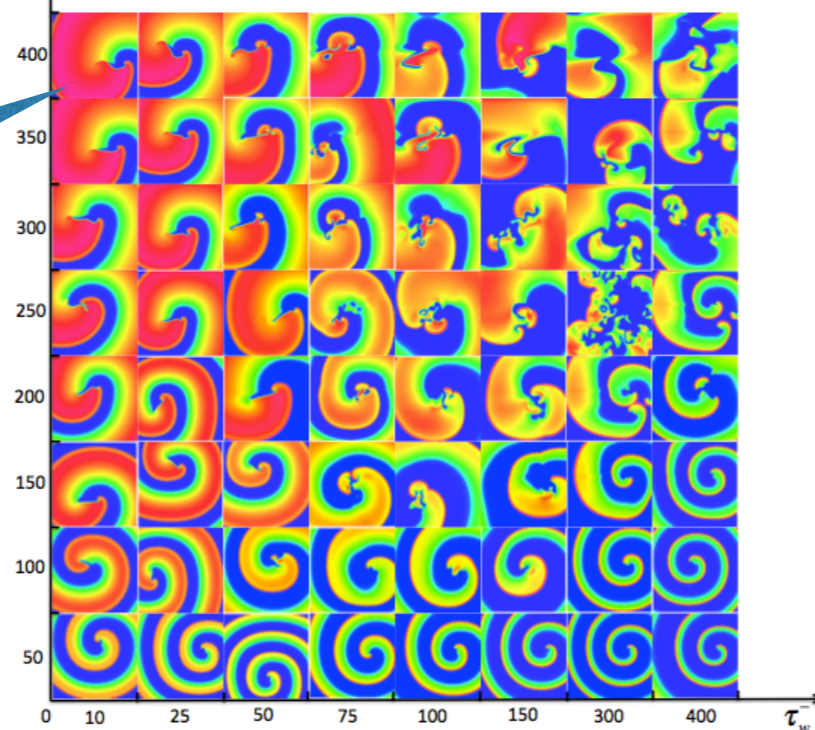


Bartocci et al.
TCS09, CAV11

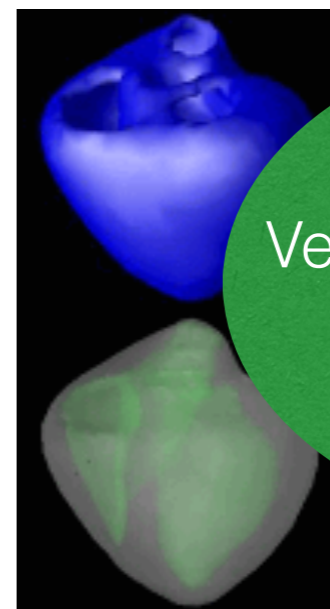
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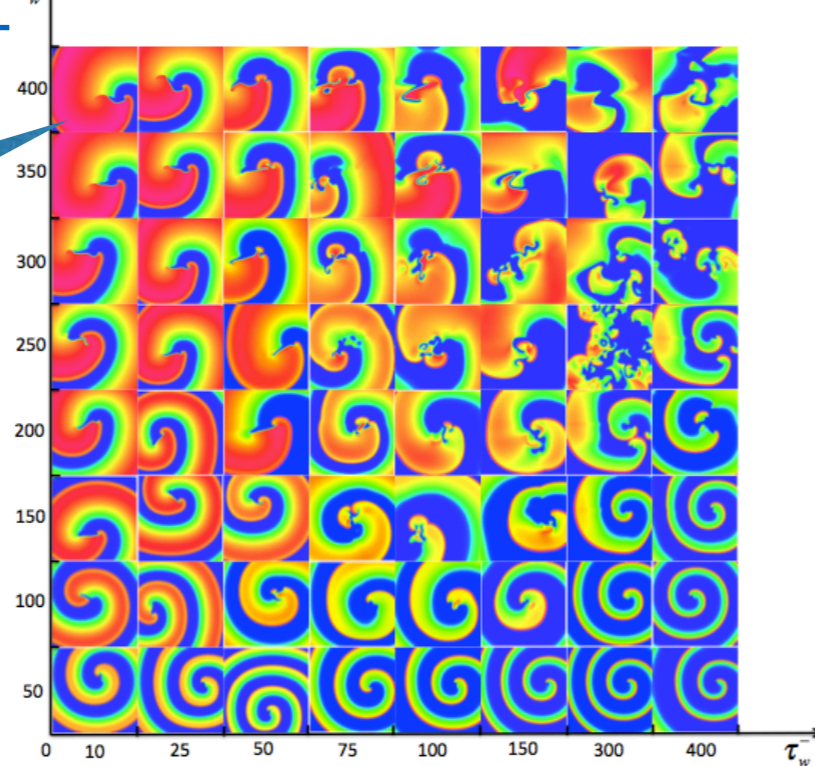
Verification requires clear semantics

Bartocci et al.
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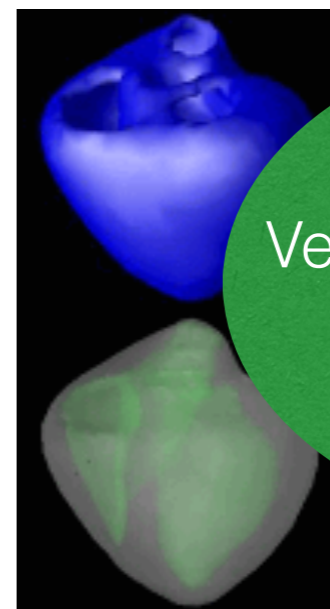
Where do they live ?



2D Tissue Model



3D Organ Model



Verification requires clear semantics

and suffers from state-space explosion

Bartocci et al.
TCS09, CAV11

Behavioural Equivalences

NFA

$$X \rightarrow 2^x (\mathcal{P}X)^A$$

Behavioural Equivalences

NFA

$$X \rightarrow 2^x (\mathcal{P}X)^A$$

language
equivalence

Behavioural Equivalences

NFA

$$X \rightarrow 2 \times (\mathcal{P}X)^A$$

Two states are equivalent iff the languages recognised from these two states are the same.

language
equivalence

Behavioural Equivalences

NFA

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language
equivalence

bisimilarity

Behavioural Equivalences

NFA

$$X \rightarrow 2 \times (\mathcal{P}X)^A$$

Two states are equivalent iff the languages recognised from these two states are the same.

language
equivalence

An equivalence relation $R \subseteq X \times X$ is a bisimulation of the NFA $(o, n): X \rightarrow 2 \times (\mathcal{P}X)^A$ iff whenever $(x, y) \in R$, we have $o(x) = o(y)$ and for all $a \in A$

$$x \xrightarrow{a} x' \quad \Rightarrow \quad \exists y'. y \xrightarrow{a} y' \wedge (x', y') \in R.$$

Bisimilarity, denoted by \sim , is the largest bisimulation.

bisimilarity

Behavioural Equivalences

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language
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bisimulation

R

bisimilarity

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language
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bisimulation



bisimilarity

Behavioural Equivalences

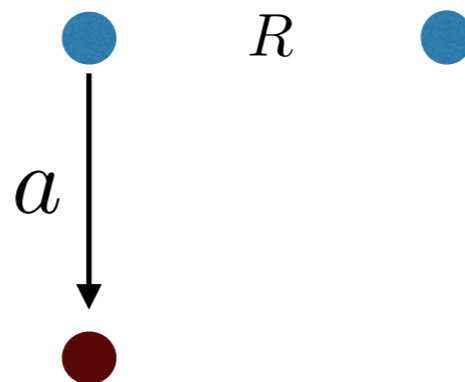
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Two states are equivalent iff the languages recognised from these two states are the same.

language equivalence

bisimulation



bisimilarity

Behavioural Equivalences

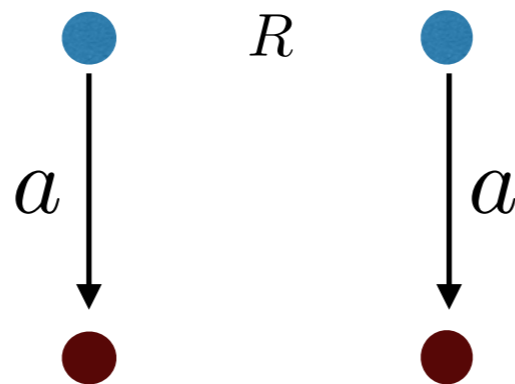
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Behavioural Equivalences

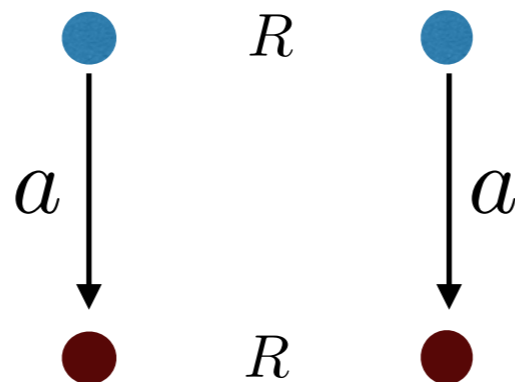
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bisimilarity

Behavioural Equivalences

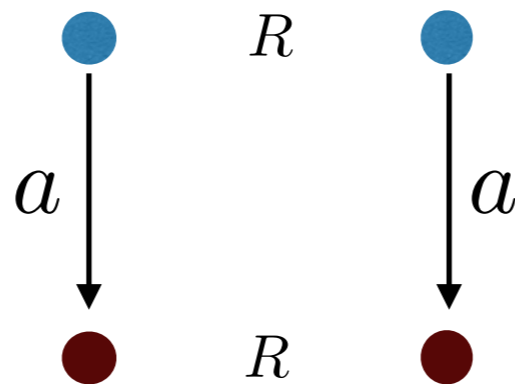
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largest
bisimulation

bisimilarity

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimilarity

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + 1$$

An equivalence relation $R \subseteq X \times X$ is a bisimulation of the MC $c: X \rightarrow \mathcal{D}X + 1$ iff whenever $(x, y) \in R$, then either $c(x) = c(y) = *$ or for all R -equivalence classes C we have

$$\sum_{z \in C} c(x)(z) = \sum_{z \in C} c(y)(z).$$

Bisimilarity, denoted by \sim , is the largest bisimulation.

bisimilarity

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + 1$$

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bisimilarity

Why are they both called bisimilarity ?

Behavioural Equivalences

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bisimilarity

What do they have in common ?

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimilarity

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimulation

R

bisimilarity

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimulation



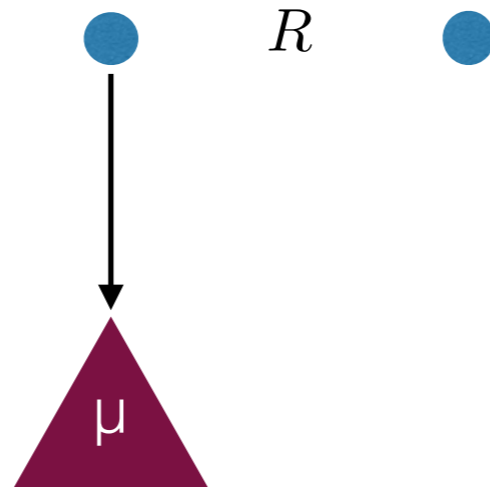
bisimilarity

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimulation



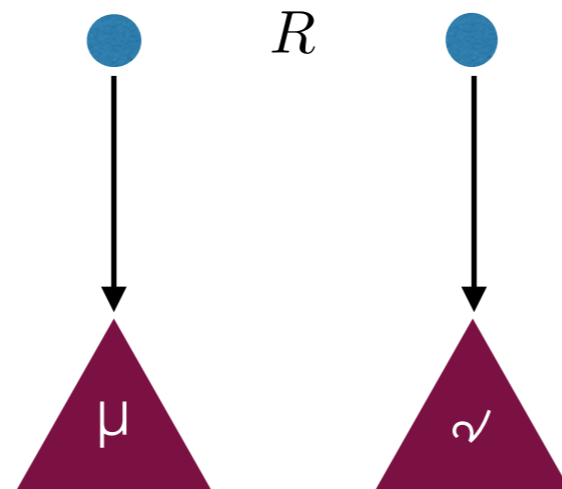
bisimilarity

Behavioural Equivalences

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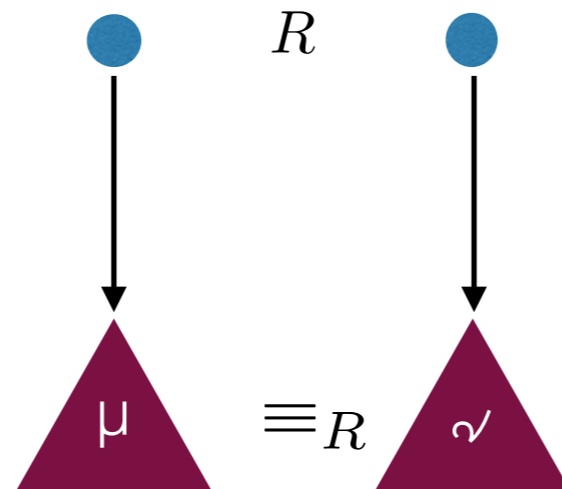
bisimilarity

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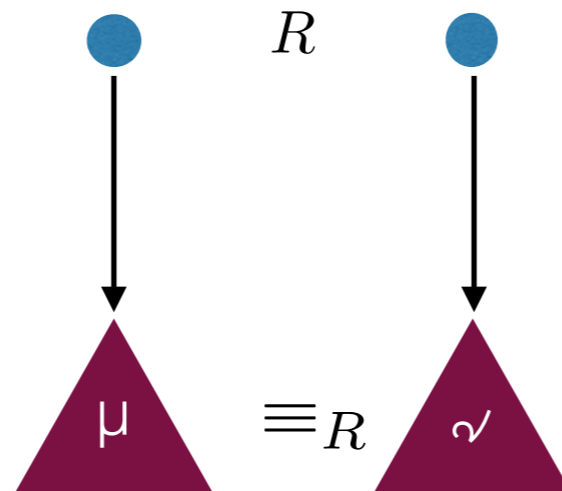
bisimilarity

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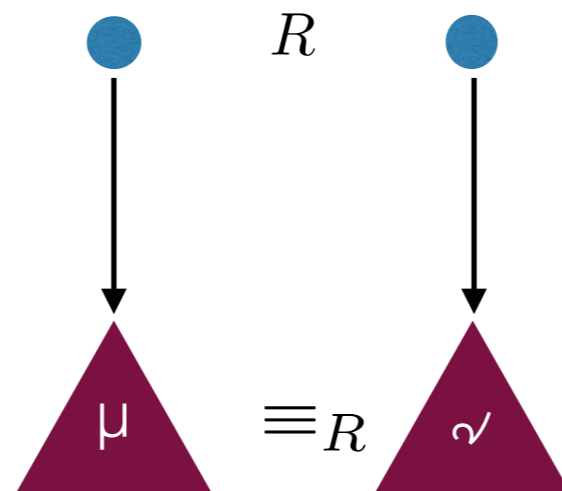
lifting of R to distributions

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimulation



bisimilarity

lifting of R to distributions

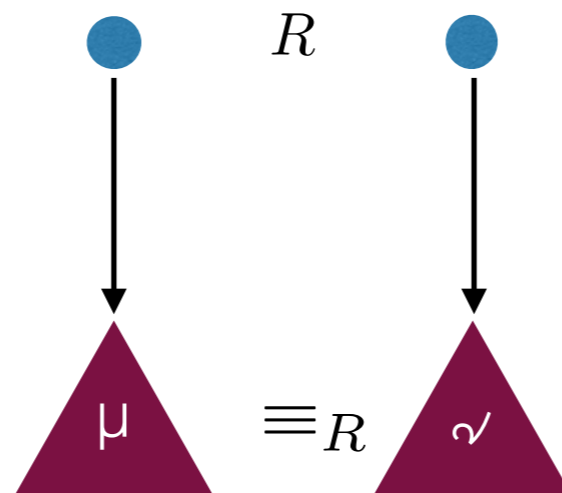
assign the same probability to "R-classes"

Behavioural Equivalences

MC

$$X \rightarrow \mathcal{D}X + I$$

bisimulation



largest bisimulation

bisimilarity

lifting of R to distributions

assign the same probability to "R-classes"



Coalgebra

Uniform framework for dynamic transition systems, based on category theory.

A coalgebra is generic transition system:



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states



Coalgebra

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$$X \xrightarrow{\mathcal{C}} FX$$

states

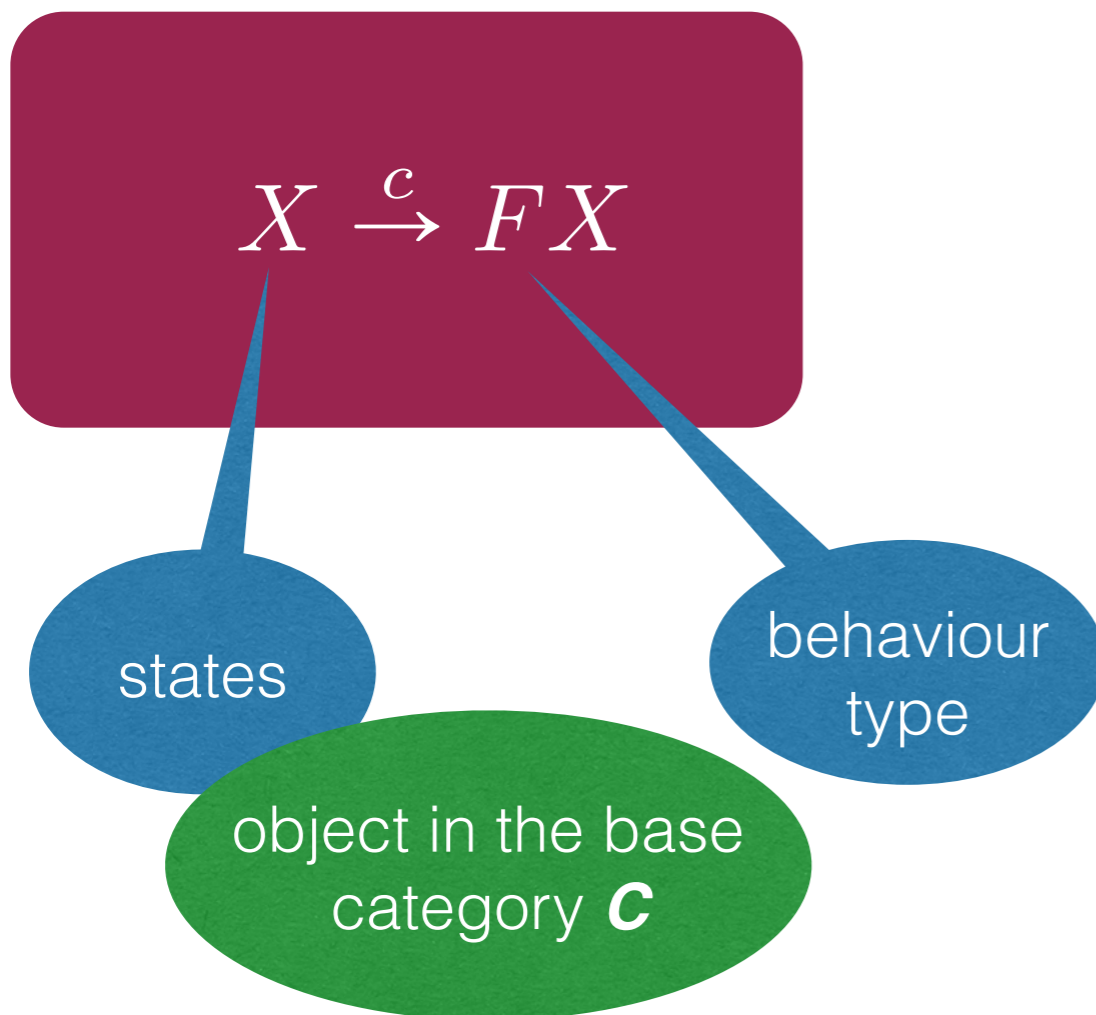
object in the base
category \mathcal{C}



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Uniform framework for dynamic transition systems, based on category theory.

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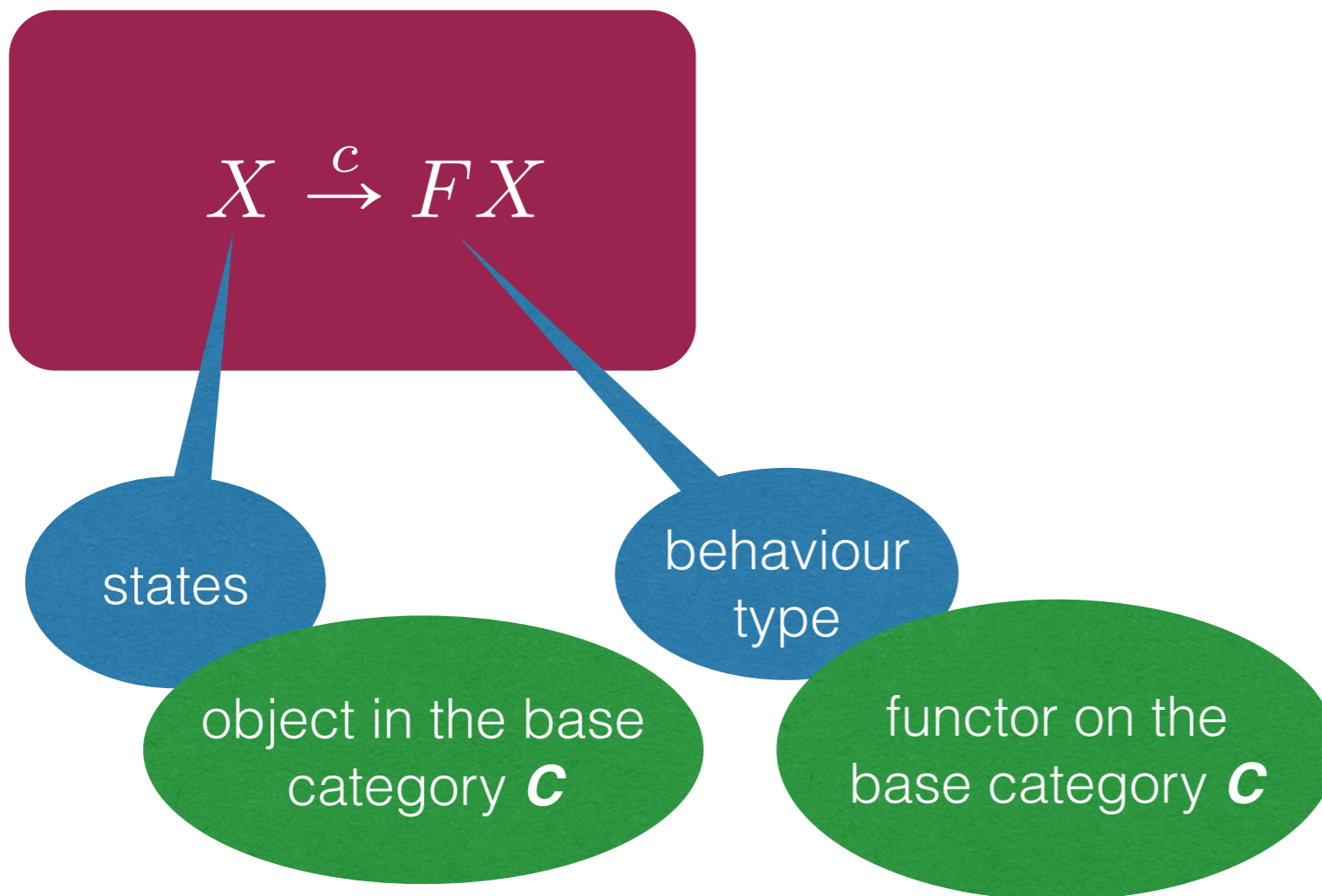




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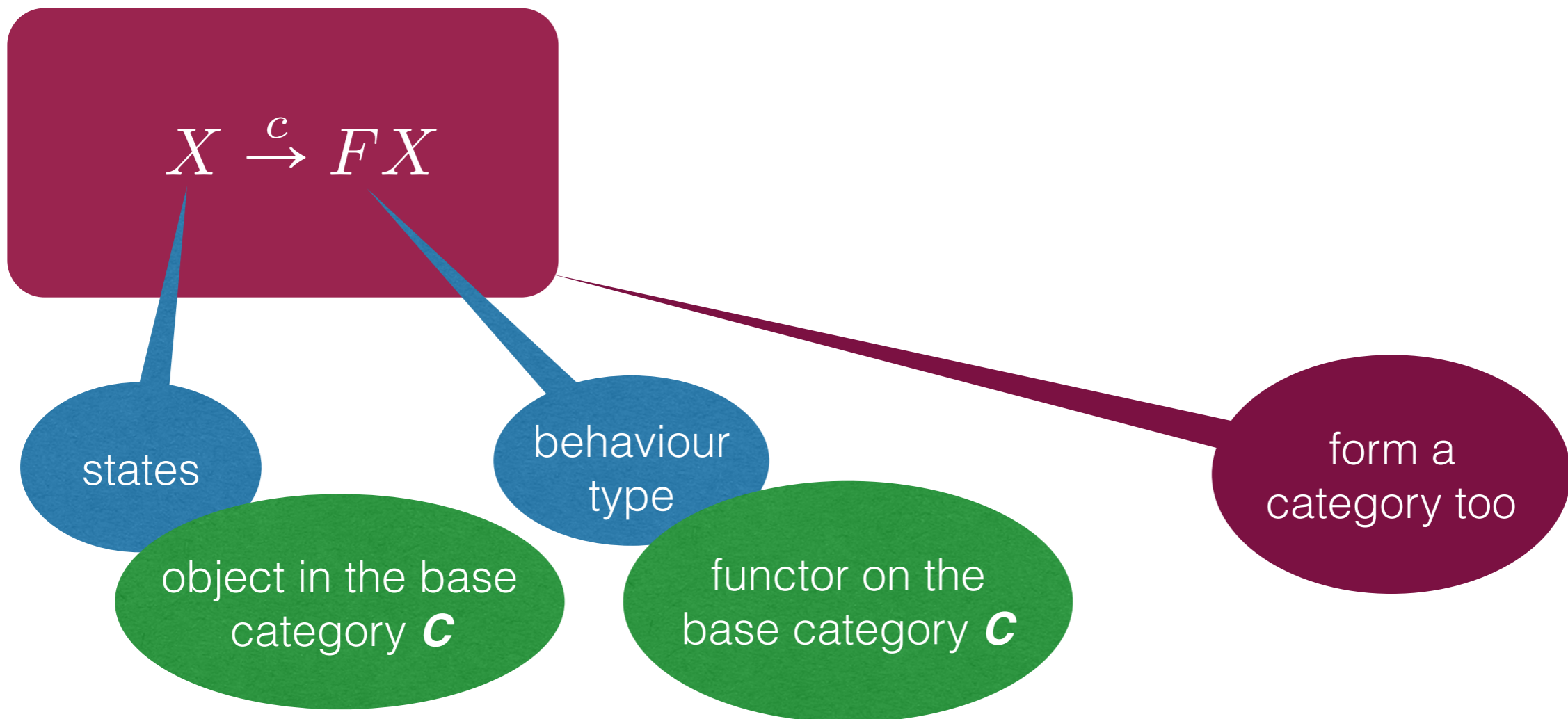




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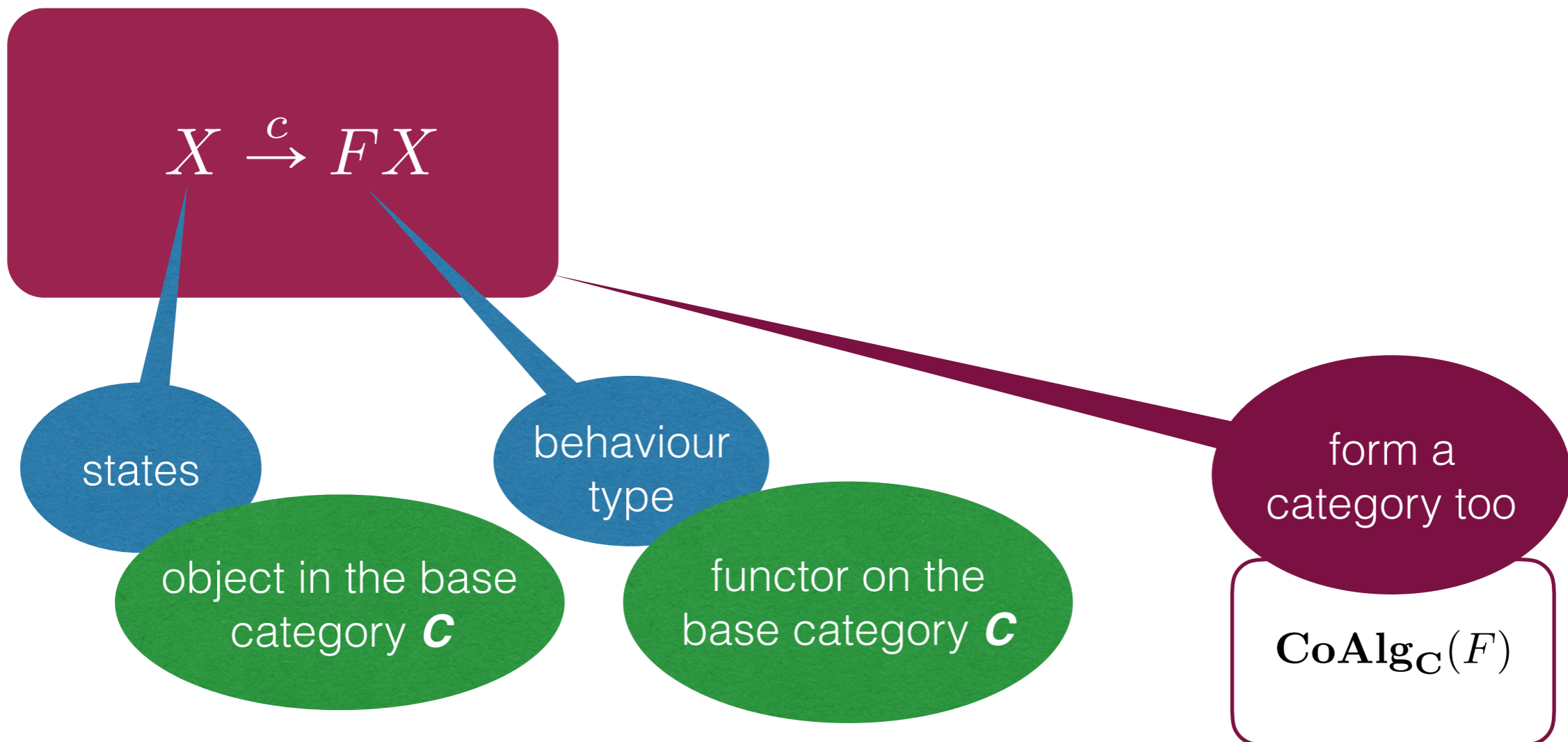




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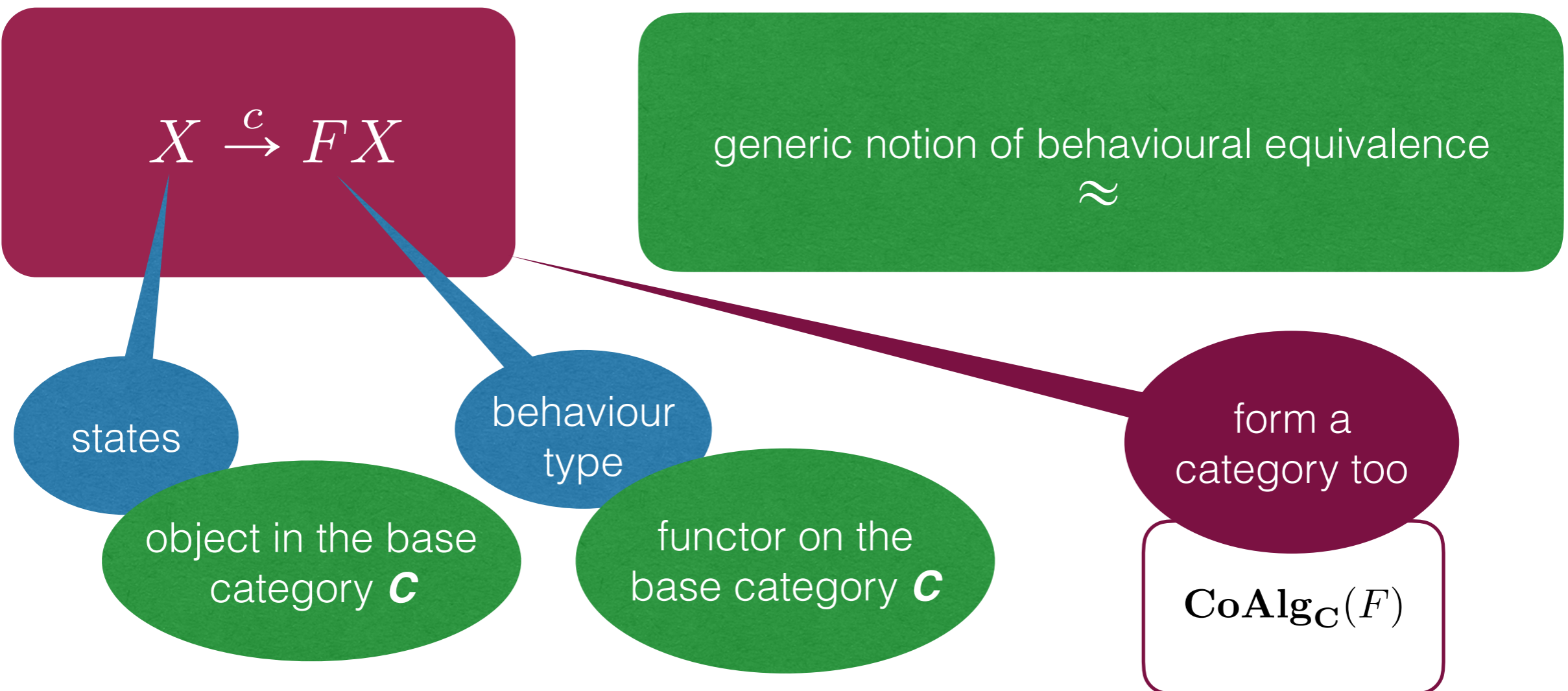




Coalgebra

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A coalgebra is generic transition system:





The category of F -coalgebras

$\mathbf{CoAlg}_{\mathbf{C}}(F)$

Objects = coalgebras

Arrows = coalgebra homomorphisms

Two states $x, y \in X$ are behaviourally equivalent, notation $x \approx y$ iff there exists a coalgebra homomorphism $h: X \rightarrow Y$ from $c: X \rightarrow FX$ to some coalgebra $d: Y \rightarrow FY$ such that $h(x) = h(y)$.



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behaviour-
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$$X \xrightarrow{c} FX$$

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$$\begin{array}{ccc} X & \xrightarrow{h} & Y \\ c_X \downarrow & & \downarrow c_Y \\ FX & \xrightarrow{Fh} & FY \end{array}$$

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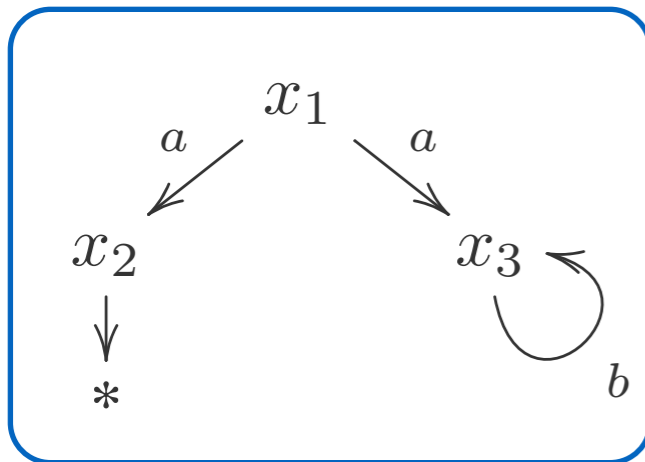
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Behavioural equivalence is bisimilarity

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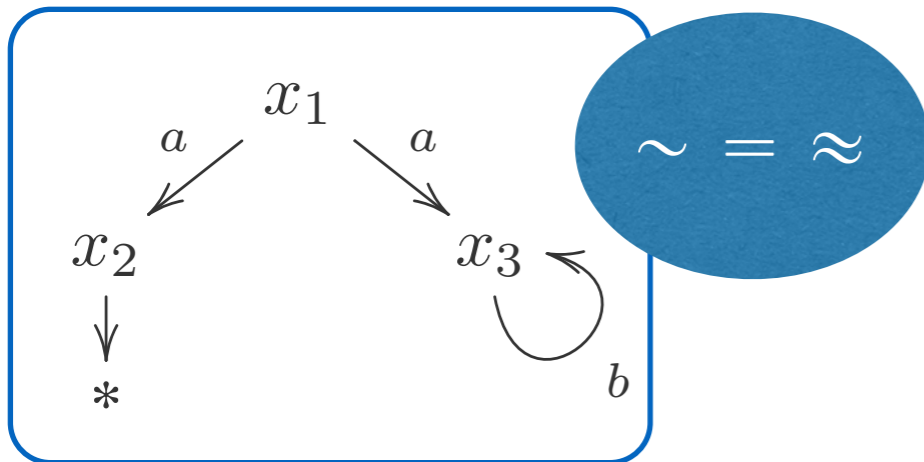
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Behavioural equivalence is bisimilarity

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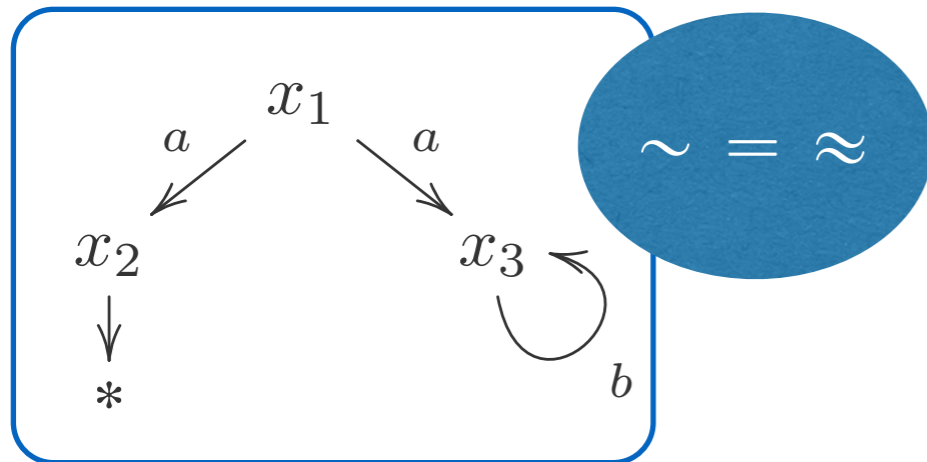
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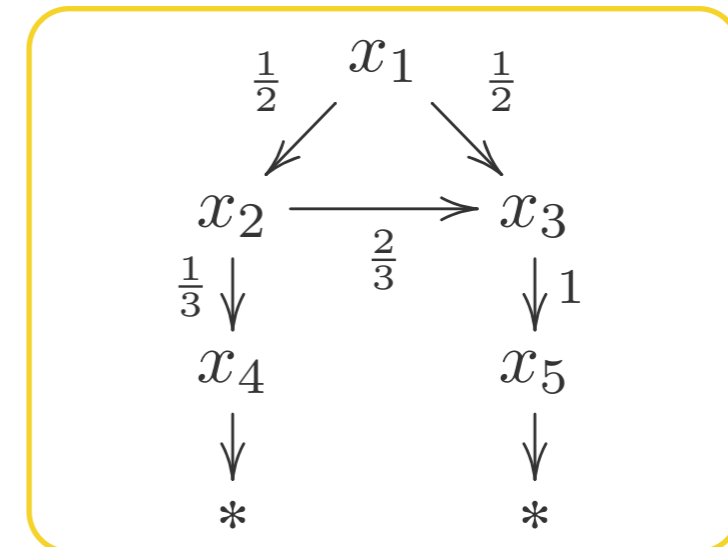
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MC

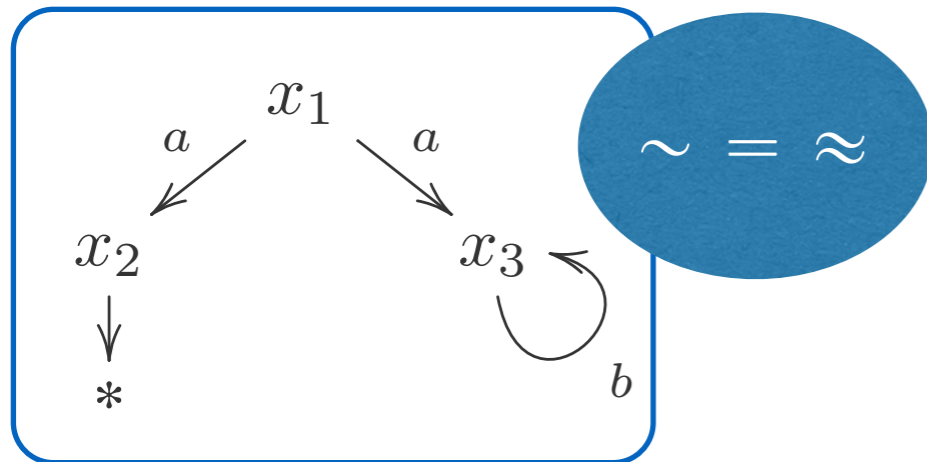
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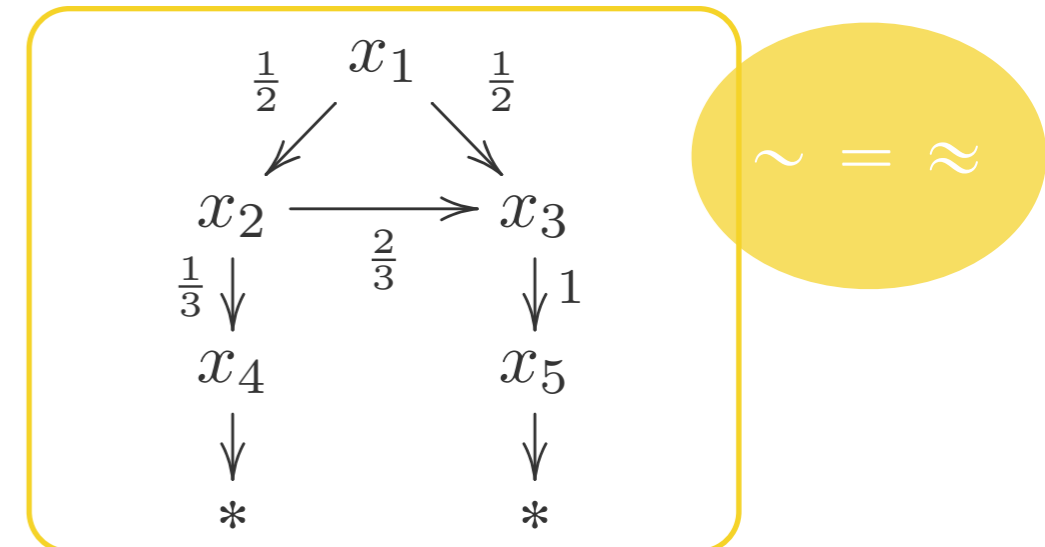
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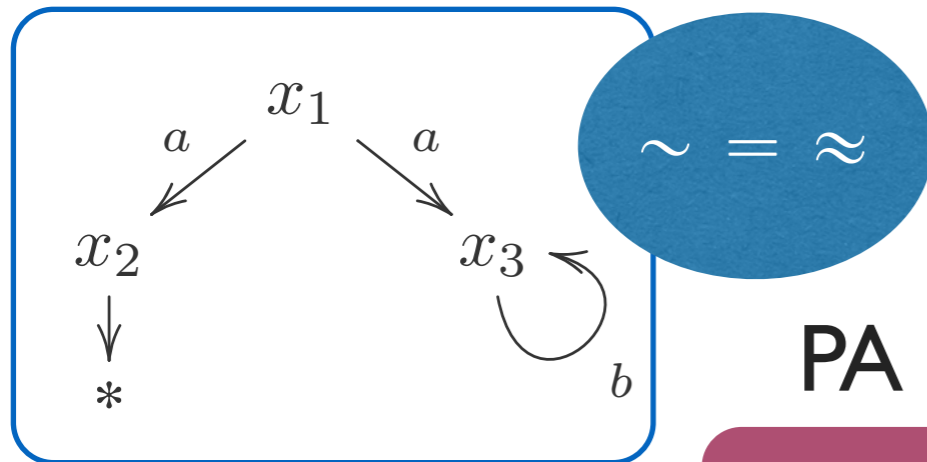
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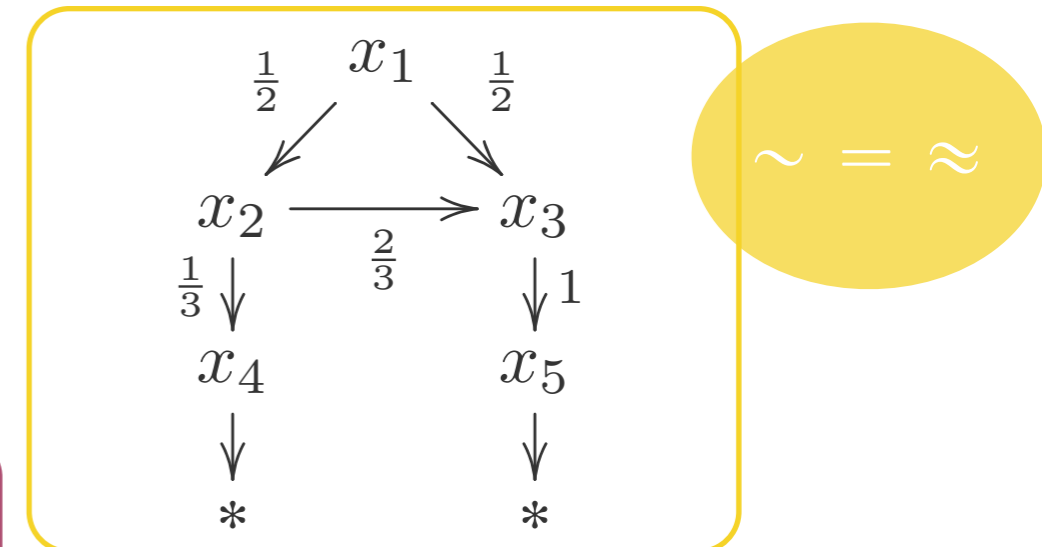
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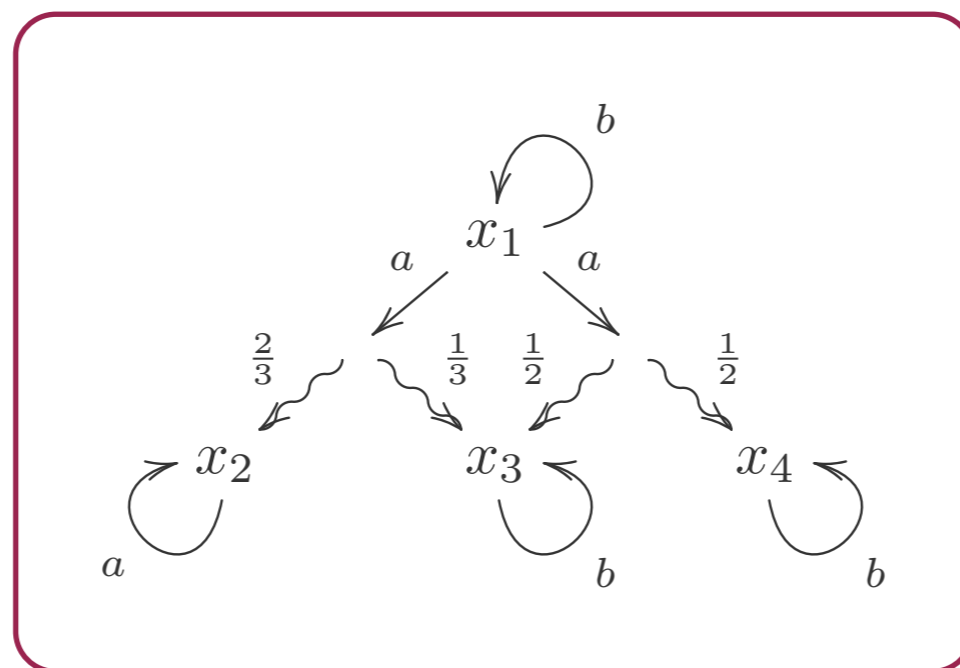
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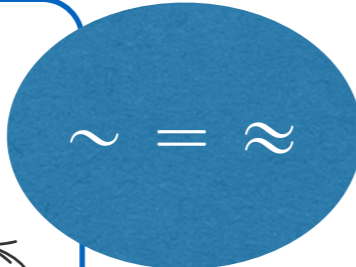
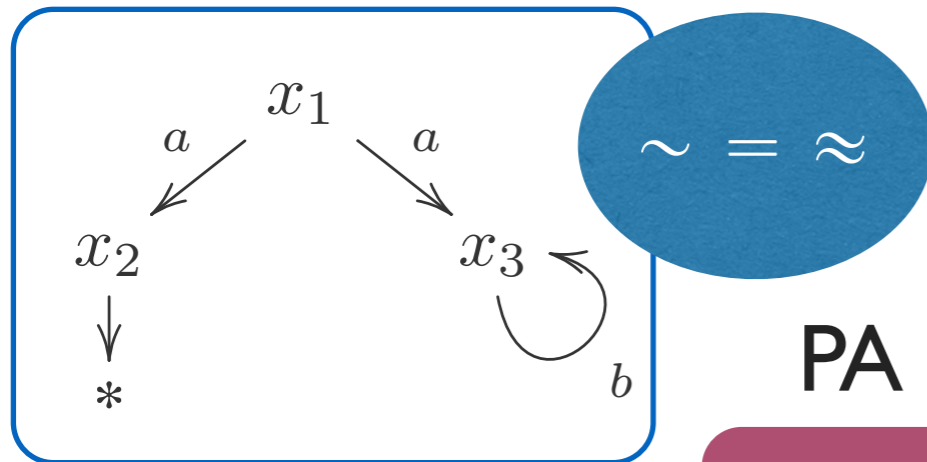
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Behavioural equivalence is bisimilarity

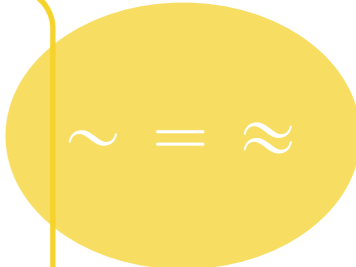
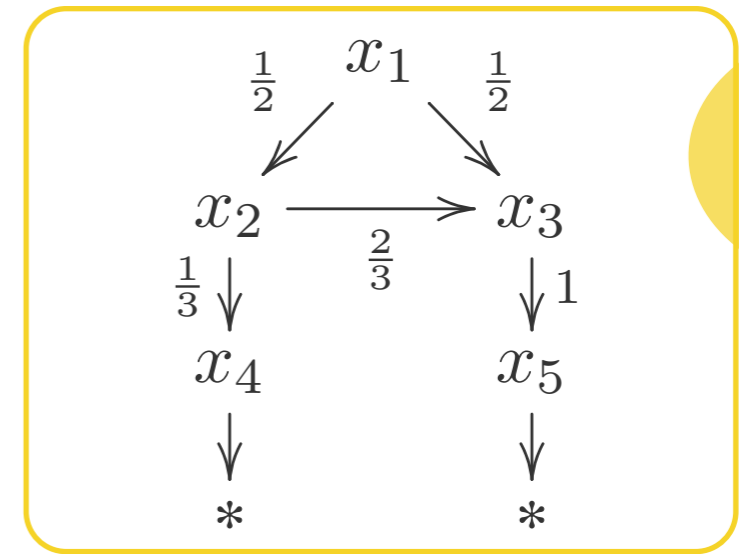
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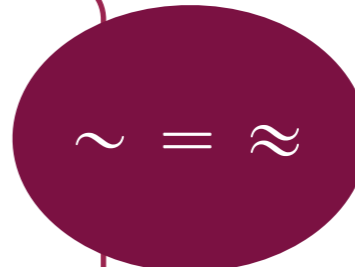
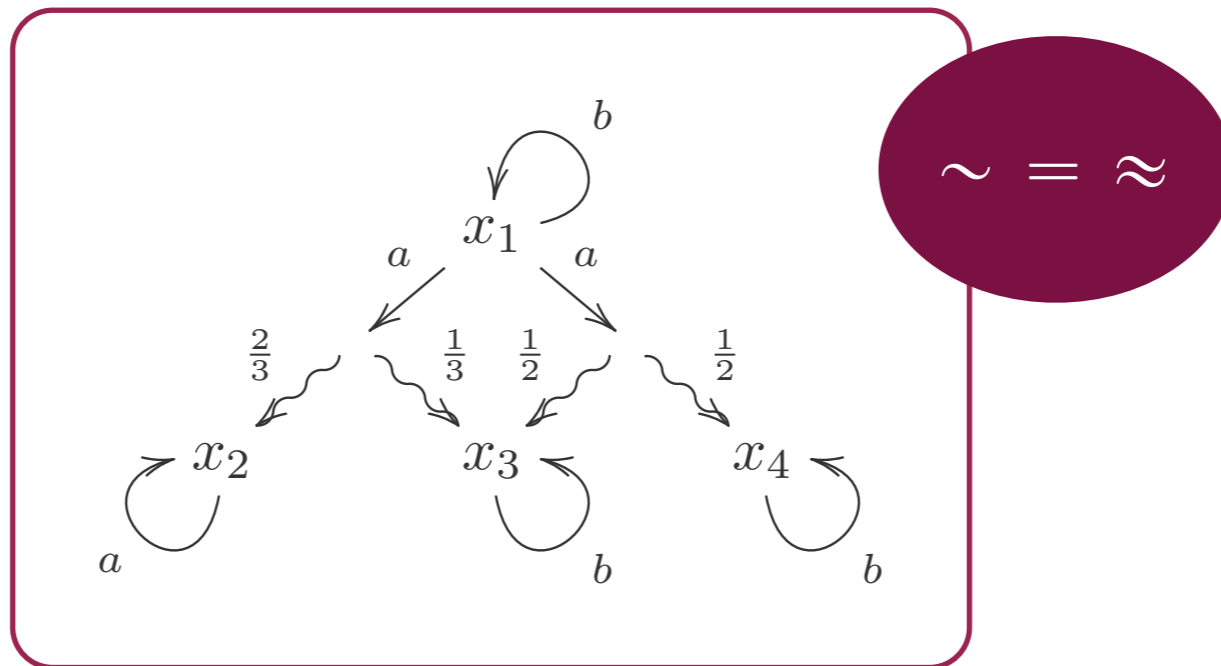
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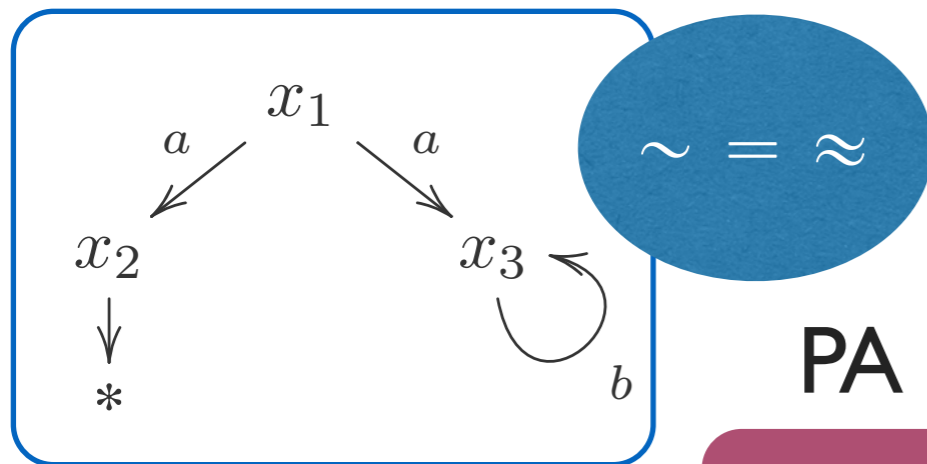
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Behavioural equivalence is bisimilarity

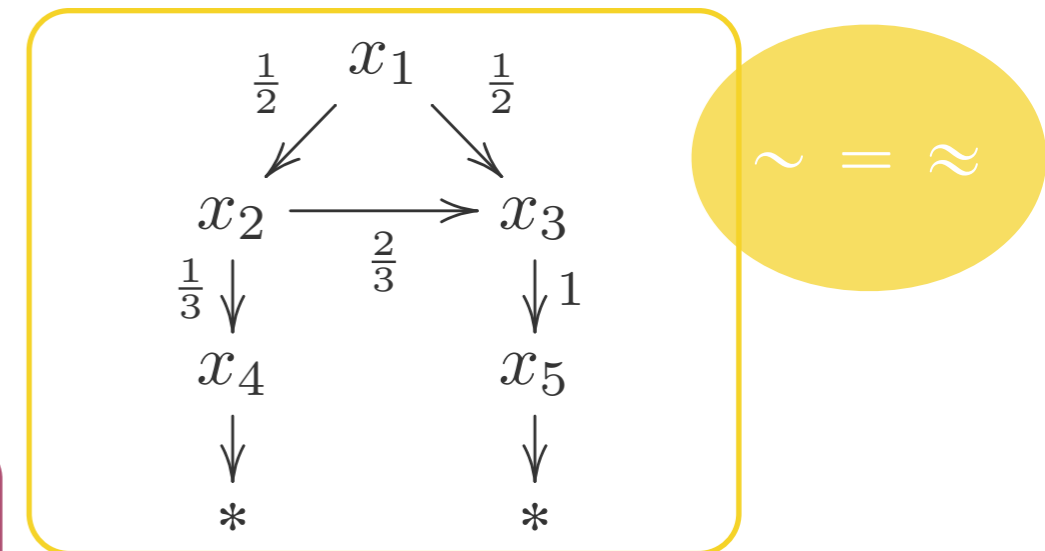
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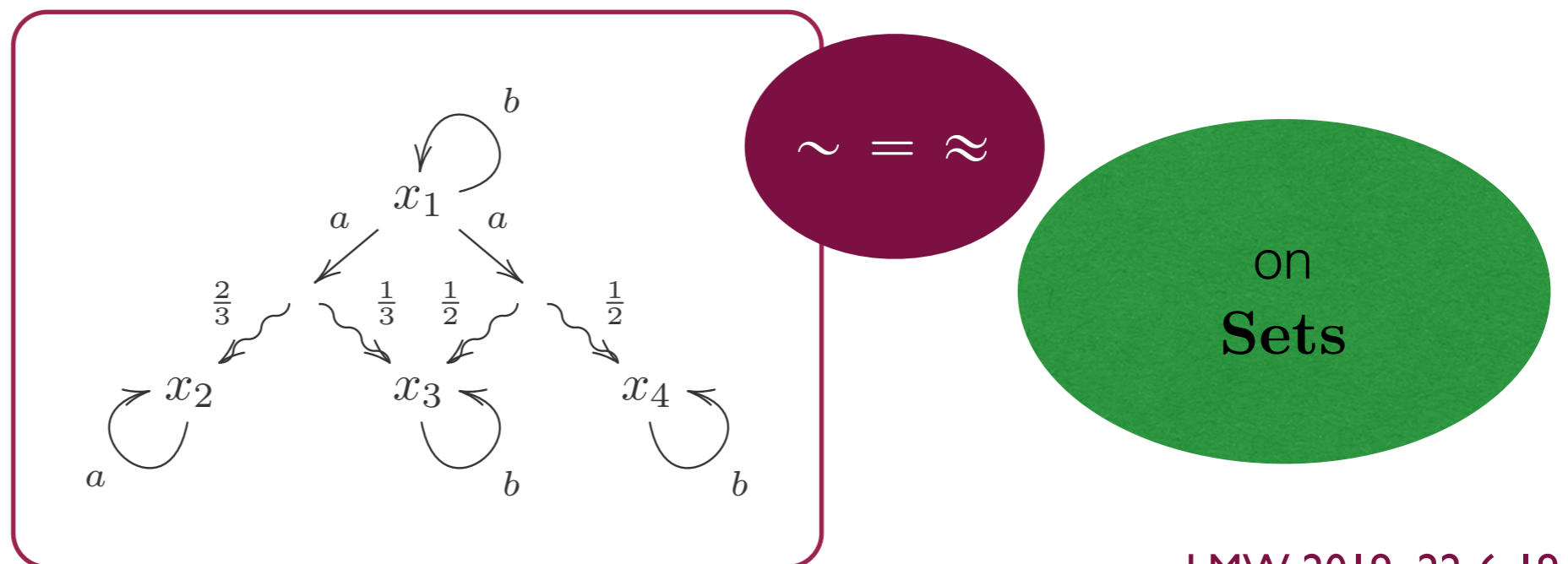
MC

$$X \rightarrow \mathcal{D}X + I$$



PA

$$X \rightarrow (\mathcal{P}\mathcal{D}X)^A$$



on Sets

Isn't that beautiful ?

Isn't that beautiful ?



if yes, read Rutten
and Jacobs!

Isn't that beautiful ?

if yes, read Rutten
and Jacobs!

and come to my talk
tomorrow at WiL

Isn't that beautiful ?

if yes, read Rutten
and Jacobs!

and come to my talk
tomorrow at Wil

and to the talk of
our LICS paper on
Wednesday





Beyond coalgebra

Beyond coalgebra

What is the best about
doing science ?



it's rewarding

relevance

elegance

What I love about
doing science

communicating

striving for
perfection

novelty

creativity

community

beauty

open

challenging

meaningful

work with people

explaining

work alone

freedom

discovering

integrity

joy

Let's bring some order here

communicating			challenging		open
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- Do what you do best and love
- Choose relevant topics
- Mix topics
- Learn from masters
- Dare to be independent
- Leave a trace
- There is plenty to learn: Exchange roles
- Do not worry much
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